# Discrete matrix delayed exponential for two delays and its property 

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#### Abstract

In recent papers, a discrete matrix delayed exponential for a single delay was defined and its main property connected with the solution of linear discrete systems with a single delay was proved. In the present paper, a generalization of the concept of discrete matrix delayed exponential is designed for two delays and its main property is proved as well.


## Introduction

Throughout the paper, we use the following notation. For integers $s, t$, $s \leq t$, we define a set $\mathbb{Z}_{s}^{t}:=\{s, s+1, \ldots, t-1, t\}$. Similarly, we define sets $\mathbb{Z}_{-\infty}^{t}:=\{\ldots, t-1, t\}$ and $\mathbb{Z}_{s}^{\infty}:=\{s, s+$ $1, \ldots\}$. The function $\lfloor\cdot\rfloor$ used below is the floor integer function. We employ the following property of the floor integer function:

$$
\begin{equation*}
x-1<\lfloor x\rfloor \leq x, \tag{1}
\end{equation*}
$$

where $x \in \mathbb{R}$.
Define binomial coefficients as customary, i.e., for $n \in \mathbb{Z}$ and $k \in \mathbb{Z}$,

$$
\binom{n}{k}:= \begin{cases}\frac{n!}{k!(n-k)!} & \text { if } n \geq k \geq 0  \tag{2}\\ 0 & \text { otherwise }\end{cases}
$$

We recall that for a well-defined discrete function $f(k)$, the forward difference operator $\Delta$ is defined as $\Delta f(k)=f(k+1)-f(k)$. In the paper, we also adopt the customary notation $\sum_{i=i_{1}}^{i_{2}} g_{i}=0$ if $i_{2}<i_{1}$. In the case of double sums, we set

$$
\begin{equation*}
\sum_{i=i_{1}, j=j_{1}}^{i_{2}, j_{2}} g_{i j}=0 \tag{3}
\end{equation*}
$$

if at least one of the inequalities $i_{2}<i_{1}, j_{2}<j_{1}$ holds.
In $[1,2]$, a discrete matrix delayed exponential for a single delay $m \in \mathbb{N}$ was defined as follows.

[^0]Definition 1 For an $r \times r$ constant matrix $B, k \in \mathbb{Z}$, and fixed $m \in \mathbb{N}$, we define the discrete matrix delayed exponential $\mathrm{e}_{m}^{B k}$ as follows:

$$
\mathrm{e}_{m}^{B k}:= \begin{cases}\Theta & \text { if } k \in \mathbb{Z}_{-\infty}^{-m-1}, \\ I+\sum_{j=1}^{\ell} B^{j} \cdot\binom{k-m(j-1)}{j} & \text { if } \ell=0,1,2, \ldots, k \in \mathbb{Z}_{(\ell-1)(m+1)+1}^{\ell(m+1)},\end{cases}
$$

where $\Theta$ is an $r \times r$ null matrix and $I$ is an $r \times r$ unit matrix.

Next, the main property (Theorem 1 below) of discrete matrix delayed exponential for a single delay $m \in \mathbb{N}$ is proved in [1].

Theorem 1 Let B be a constant $r \times r$ matrix. Then, for $k \in \mathbb{Z}_{-m}^{\infty}$,

$$
\begin{equation*}
\Delta \mathrm{e}_{m}^{B k}=B \mathrm{e}_{m}^{B(k-m)} . \tag{4}
\end{equation*}
$$

The paper is concerned with a generalization of the notion of discrete matrix delayed exponential for two delays and a proof of one of its properties, similar to the main property (4) of discrete matrix delayed exponential for a single delay.

## Discrete matrix delayed exponential for two delays and its main property

We define a discrete $r \times r$ matrix function $\mathrm{e}_{m n}^{B C k}$ called the discrete matrix delayed exponential for two delays $m, n \in \mathbb{N}, m \neq n$ and for two $r \times r$ commuting constant matrices $B$, $C$ as follows.

Definition 2 Let $B, C$ be constant $r \times r$ matrices with the property $B C=C B$ and let $m, n \in \mathbb{N}, m \neq n$ be fixed integers. We define a discrete $r \times r$ matrix function $\mathrm{e}_{m n}^{B C k}$ called the discrete matrix delayed exponential for two delays $m, n$ and for two $r \times r$ constant matrices $B, C$ :

$$
\mathrm{e}_{m n}^{B C k}:= \begin{cases}\Theta & \text { if } k \in \mathbb{Z}_{-\infty}^{-\max \{m, n\}-1}, \\ I & \text { if } k \in \mathbb{Z}_{-\max \{m, n\}}^{0}, \\ I+(B+C) \sum_{i=0, j=0}^{p_{(k)}^{-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1}} & \text { if } k \in \mathbb{Z}_{1}^{\infty},\end{cases}
$$

where

$$
\begin{equation*}
p_{(k)}:=\left\lfloor\frac{k+m}{m+1}\right\rfloor, \quad q_{(k)}:=\left\lfloor\frac{k+n}{n+1}\right\rfloor . \tag{5}
\end{equation*}
$$

The main property of $\mathrm{e}_{m n}^{B C k}$ is given by the following theorem.

Theorem 2 Let $B, C$ be constant $r \times r$ matrices with the property $B C=C B$ and let $m, n \in \mathbb{N}$, $m \neq n$ be fixed integers. Then

$$
\begin{equation*}
\Delta \mathrm{e}_{m n}^{B C k}=B \mathrm{e}_{m n}^{B C(k-m)}+C \mathrm{e}_{m n}^{B C(k-n)} \tag{6}
\end{equation*}
$$

holds for $k \geq 0$.

Proof Let $k \geq 1$. From (1) and (5), we can see easily that, for an integer $k \geq 0$ satisfying

$$
\left(p_{(k)}-1\right)(m+1)+1 \leq k \leq p_{(k)}(m+1) \wedge\left(q_{(k)}-1\right)(n+1)+1 \leq k \leq q_{(k)}(n+1),
$$

the relation

$$
\Delta \mathrm{e}_{m n}^{B C k}=\Delta\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}^{-1, q_{(k)}-1}} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1}\right]
$$

holds in accordance with Definition 2 of $\mathrm{e}_{m n}^{B C k}$. Since $\Delta I=\Theta$, we have

$$
\begin{equation*}
\Delta \mathrm{e}_{m n}^{B C k}=\Delta\left[(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1}\right] . \tag{7}
\end{equation*}
$$

Considering the increment by its definition, i.e.,

$$
\begin{equation*}
\Delta \mathrm{e}_{m n}^{B C k}=\mathrm{e}_{m n}^{B C(k+1)}-\mathrm{e}_{m n}^{B C k}, \tag{8}
\end{equation*}
$$

we conclude that it is reasonable to divide the proof into four parts with respect to the value of integer $k$. In case one, $k$ is such that

$$
\left(p_{(k)}-1\right)(m+1)+1 \leq k<p_{(k)}(m+1) \wedge\left(q_{(k)}-1\right)(n+1)+1 \leq k<q_{(k)}(n+1)
$$

in case two

$$
k=p_{(k)}(m+1) \wedge\left(q_{(k)}-1\right)(n+1)+1 \leq k<q_{(k)}(n+1),
$$

in case three

$$
\left(p_{(k)}-1\right)(m+1)+1 \leq k<p_{(k)}(m+1) \wedge k=q_{(k)}(n+1)
$$

and in case four

$$
k=p_{(k)}(m+1) \wedge k=q_{(k)}(n+1) .
$$

We see that the above cases cover all the possible relations between $k, p_{(k)}$ and $q_{(k)}$.
In the proof, we use the identities

$$
\begin{equation*}
\binom{n+1}{k}=\binom{n}{k}+\binom{n}{k-1}, \tag{9}
\end{equation*}
$$

where $n, k \in \mathbb{N}$ and

$$
\begin{equation*}
\binom{i}{i}=\binom{i-1}{i-1}, \quad\binom{j}{0}=\binom{j-1}{0}, \quad\binom{i+j}{i}=\binom{i+j-1}{i-1}+\binom{i+j-1}{i}, \tag{10}
\end{equation*}
$$

where $i, j \in \mathbb{N}$, which are derived from (2) and (9).
I. $\left(p_{(k)}-1\right)(m+1)+1 \leq k<p_{(k)}(m+1) \wedge\left(q_{(k)}-1\right)(n+1)+1 \leq k<q_{(k)}(n+1)$

From (1) and (5), we get

$$
\begin{aligned}
& p_{(k-m)}=\left\lfloor\frac{k-m+m}{m+1}\right\rfloor \leq \frac{k}{m+1}<p_{(k)}, \\
& p_{(k-m)}=\left\lfloor\frac{k-m+m}{m+1}\right\rfloor>\frac{k}{m+1}-1=\frac{k-m-1}{m+1}>p_{(k)}-2 .
\end{aligned}
$$

Therefore, $p_{(k-m)}=p_{(k)}-1$ and, by Definition 2,

Similarly, omitting details, we get (using (1), and (5)) $q_{(k-n)}=q_{(k)}-1$ and

$$
\begin{equation*}
\mathrm{e}_{m n}^{B C(k-n)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k-n)}-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-n-m i-n j}{i+j+1} . \tag{12}
\end{equation*}
$$

Let $q_{(k-m)} \geq 1$. We show that

$$
\begin{equation*}
\binom{k-m-m i-n j}{i+j+1}=0 \quad \text { if } i \geq 0, j \geq q_{(k-m)} . \tag{13}
\end{equation*}
$$

In accordance with (1),

$$
q_{(k-m)}=\left\lfloor\frac{k-m+n}{n+1}\right\rfloor>\frac{k-m+n}{n+1}-1=\frac{k-m-1}{n+1}
$$

or

$$
k-m<(n+1) q_{(k-m)}+1 \leq(m+1) i+(n+1) j+1 \quad \text { if } i \geq 0, j \geq q_{(k-m)} .
$$

From the last inequality, we get

$$
k-m-m i-n j<i+j+1 \quad \text { if } i \geq 0, j \geq q_{(k-m)}
$$

and (13) holds by (2). For that reason and since $q_{(k-m)} \leq q_{(k)}$, we can replace $q_{(k-m)}$ by $q_{(k)}$ in (11). Thus, we have

$$
\begin{equation*}
\mathrm{e}_{m n}^{B C(k-m)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1} . \tag{14}
\end{equation*}
$$

It is easy to see that, due to (3), formula (14) can be used instead of (11) if $q_{(k-m)}<1$ also.
Let $p_{(k-n)} \geq 1$. Similarly, we can show that

$$
\binom{k-n-m i-n j}{i+j+1}=0 \quad \text { if } i \geq p_{(k-n), j \geq 0}
$$

and, since $p_{(k-n)} \leq p_{(k)}$, we can replace $p_{(k-n)}$ by $p_{(k)}$ in (12). Thus, we have

$$
\begin{equation*}
\mathrm{e}_{m n}^{B C(k-n)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1} . \tag{15}
\end{equation*}
$$

It is easy to see that, due to (3), formula (15) can be used instead of (12) if $p_{(k-n)}<1$, too.
Due to (1), we also conclude that

$$
\begin{equation*}
p_{(k+1)}=p_{(k)}, \quad q_{(k+1)}=q_{(k)} \tag{16}
\end{equation*}
$$

because

$$
p_{(k+1)}=\left\lfloor\frac{k+1+m}{m+1}\right\rfloor \leq \frac{k}{m+1}+1<p_{(k)}+1
$$

and

$$
p_{(k+1)}=\left\lfloor\frac{k+1+m}{m+1}\right\rfloor>\frac{k+1+m}{m+1}-1=\frac{k}{m+1} \geq p_{(k)}-1+\frac{1}{m+1} .
$$

The second formula can be proved similarly.
Now we are able to prove that

$$
\begin{align*}
\Delta \mathrm{e}_{m n}^{B C k}= & B \mathrm{e}_{m n}^{B C(k-m)}+C \mathrm{e}_{m n}^{B C(k-n)} \\
= & B\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right] \\
& +C\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}\right] . \tag{17}
\end{align*}
$$

With the aid of (7), (8), (9) and (16), we get

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & \mathrm{e}_{m n}^{B C(k+1)}-\mathrm{e}_{m n}^{B C k} \\
= & I+(B+C) \sum_{i=0, j=0}^{p_{(k+1)}^{-1, q_{(k+1)}-1}} B^{i} C^{j}\binom{i+j}{i}\binom{k+1-m i-n j}{i+j+1} \\
& -I-(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1} \\
= & I+(B+C) \sum_{i=0, j=0}^{p_{(k)}^{-1, q_{(k)}-1}} B^{i} C^{j}\binom{i+j}{i}\binom{k+1-m i-n j}{i+j+1} \\
& -I-(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1} \\
= & (B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\left[\binom{k+1-m i-n j}{i+j+1}-\binom{k-m i-n j}{i+j+1}\right]
\end{aligned}
$$

$$
\begin{aligned}
= & (B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j} \\
= & (B+C)\left[I+\sum_{i=1}^{p_{(k)}-1} B^{i} C^{0}\binom{i}{i}\binom{k-m i}{i}+\sum_{j=1}^{q_{(k)}-1} B^{0} C^{j}\binom{j}{0}\binom{k-n j}{j}\right. \\
& \left.+\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j}\right] .
\end{aligned}
$$

By (10), we have

$$
\begin{aligned}
\Delta \mathbf{e}_{m n}^{B C k}= & (B+C)\left[I+\sum_{i=1}^{p_{(k)}-1} B^{i} C^{0}\binom{i-1}{i-1}\binom{k-m i}{i}\right. \\
& +\sum_{j=1}^{q_{(k)}-1} B^{0} C^{j}\binom{j-1}{0}\binom{k-n j}{j} \\
& +\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i-1}\binom{k-m i-n j}{i+j} \\
& \left.+\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i}\binom{k-m i-n j}{i+j}\right] \\
= & (B+C)\left[\begin{array}{c}
p_{(k)}-1, q_{(k)}-1 \\
I+B_{i=1, j=0}\binom{i+j-1}{i-1}\binom{k-m i-n j}{i+j} \\
\\
\end{array}+\sum_{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i}\binom{k-m i-n j}{i+j}\right] .
\end{aligned}
$$

Now in the first sum we replace the summation index $i$ by $i+1$ and in the second sum we replace the summation index $j$ by $j+1$. Then

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & (B+C)\left[I+\sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i+1} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right. \\
& \left.+\sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j+1}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}\right] \\
= & B+B(B+C) \sum_{i=0, j=0}^{p_{(k)}^{-2, q_{(k)}-1}} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1} \\
& +C+C(B+C) \sum_{i=0, j=0}^{p_{(k)}^{-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}} \\
= & B\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right]
\end{aligned}
$$

$$
\begin{aligned}
& +C\left[I+C(B+C) \sum_{i=0, j=0}^{p_{(k)}^{-1, q_{(k)}-2}} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}\right] \\
= & B \mathrm{e}_{m n}^{B C(k-m)}+C \mathrm{e}_{m n}^{B C(k-n)} .
\end{aligned}
$$

Due to (14) and (15), we conclude that formula (17) is valid.

## II. $k=p_{(k)}(m+1) \wedge\left(q_{(k)}-1\right)(n+1)+1 \leq k<q_{(k)}(n+1)$

In this case,

$$
\begin{aligned}
& p_{(k-m)}=\left\lfloor\frac{k-m+m}{m+1}\right\rfloor=\left\lfloor\frac{k}{m+1}\right\rfloor=p_{(k)}, \\
& p_{(k+1)}=\left\lfloor\frac{k+1+m}{m+1}\right\rfloor \leq \frac{k+1+m}{m+1}=\frac{k}{m+1}+1=p_{(k)}+1, \\
& p_{(k+1)}=\left\lfloor\frac{k+1+m}{m+1}\right\rfloor>\frac{k+1+m}{m+1}-1=\frac{k}{m+1}=p_{(k)}
\end{aligned}
$$

and $p_{(k+1)}=p_{(k)}+1$. In addition to this (see relevant computations performed in case I), we have $q_{(k-n)}=q_{(k)}-1$ and $q_{(k+1)}=q_{(k)}$.

Then

$$
\begin{aligned}
& \mathrm{e}_{m n}^{B C k}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1}, \\
& \mathrm{e}_{m n}^{B C(k+1)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k+1-m i-n j}{i+j+1}
\end{aligned}
$$

and

$$
\begin{align*}
& \mathrm{e}_{m n}^{B C(k-m)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k-m)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m-m i-n j}{i+j+1},  \tag{18}\\
& \mathrm{e}_{m n}^{B C(k-n)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k-n)}^{-1, q_{(k)}-2}} B^{i} C^{j}\binom{i+j}{i}\binom{k-n-m i-n j}{i+j+1} . \tag{19}
\end{align*}
$$

Like with the computations performed in the previous part of the proof, we get

$$
\binom{k-m-m i-n j}{i+j+1}=0 \quad \text { if } i \geq 0, j \geq q_{(k-m)}
$$

and

$$
\binom{k-n-m i-n j}{i+j+1}=0 \quad \text { if } i \geq p_{(k-n)}, j \geq 0 .
$$

So, we can substitute $q_{(k-m)}$ by $q_{(k)}$ in (18) and $p_{(k-n)}$ by $p_{(k)}$ in (19).
Accordingly, we have

$$
\begin{equation*}
\mathrm{e}_{m n}^{B C(k-m)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}, \tag{20}
\end{equation*}
$$

$$
\begin{equation*}
\mathrm{e}_{m n}^{B C(k-n)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1} . \tag{21}
\end{equation*}
$$

It is easy to see that, due to (3), formula (20) can also be used instead of (18) if $q_{(k-m)}<1$ and formula (21) can also be used instead of (19) if $p_{(k-n)}<1$.

We have to prove

$$
\begin{align*}
\Delta \mathrm{e}_{m n}^{B C k}= & B \mathrm{e}_{m n}^{B C(k-m)}+C \mathrm{e}_{m n}^{B C(k-n)} \\
= & B\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right] \\
& +C\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}\right] . \tag{22}
\end{align*}
$$

Therefore,

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & \mathrm{e}_{m n}^{B C(k+1)}-\mathrm{e}_{m n}^{B C k} \\
= & I+(B+C) \sum_{i=0, j=0}^{p_{(k)}, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k+1-m i-n j}{i+j+1} \\
& -I-(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1} \\
= & (B+C)\left[\begin{array}{c}
\sum_{i=0, j=0}^{p_{(k)}^{-1, q_{(k)}-1}} B^{i} C^{j}\binom{i+j}{i}\left[\binom{k+1-m i-n j}{i+j+1}-\binom{k-m i-n j}{i+j+1}\right] \\
\\
\end{array}+\sum_{j=0}^{q_{(k)}-1} B^{p_{(k)}} C^{j}\binom{p_{(k)}+j}{p_{(k)}}\binom{k+1-m p_{(k)}-n j}{p_{(k)}+j+1}\right] .
\end{aligned}
$$

With the aid of the equation $k=p_{(k)}(m+1)$, we get

$$
\binom{k+1-m p_{(k)}-n j}{p_{(k)}+j+1}=\binom{p_{(k)}+1-n j}{p_{(k)}+1+j}=0 \quad \text { if } j>0
$$

and, by (9), we have

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & (B+C)\left[\sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j}+B^{p_{(k)}}\right] \\
= & (B+C)\left[I+\sum_{i=1}^{p_{(k)}-1} B^{i} C^{0}\binom{i}{i}\binom{k-m i}{i}+\sum_{j=1}^{q_{(k)}-1} B^{0} C^{j}\binom{j}{0}\binom{k-n j}{j}\right. \\
& \left.+\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j}+B^{p_{(k)}}\right] .
\end{aligned}
$$

By (10), we have

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & (B+C)\left[I+\sum_{i=1}^{p_{(k)}-1} B^{i} C^{0}\binom{i-1}{i-1}\binom{k-m i}{i}+\sum_{j=1}^{q_{(k)}-1} B^{0} C^{j}\binom{j-1}{0}\binom{k-n j}{j}\right. \\
& +\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i-1}\binom{k-m i-n j}{i+j} \\
& \left.+\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i}\binom{k-m i-n j}{i+j}+B^{p_{(k)}}\right] \\
= & (B+C)\left[I+\sum_{i=1, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i-1}\binom{k-m i-n j}{i+j}\right. \\
& \left.+\sum_{i=0, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i}\binom{k-m i-n j}{i+j}+B^{p_{(k)}}\right] .
\end{aligned}
$$

Now we replace in the first sum the summation index $i$ by $i+1$ and in the second sum we replace the summation index $j$ by $j+1$. Then

$$
\begin{aligned}
& \Delta \mathrm{e}_{m n}^{B C k}=(B+C)\left[I+\sum_{i=0, j=0}^{p(k)-2, q(k)-1} B^{i+1} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right. \\
& \left.+\sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j+1}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}+B^{p_{(k)}}\right] \\
& =B+B(B+C) \sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}+B^{p(k)}(B+C) \\
& +C+C(B+C) \sum_{i=0, j=0}^{p(k)-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1} \\
& =B\left[I+(B+C)\left(\sum_{i=0, j=0}^{p(k)-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}+B^{p_{(k)}-1}\right)\right] \\
& +C\left[I+(B+C) \sum_{i=0, j=0}^{p(k)-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}\right] .
\end{aligned}
$$

For $k=p_{(k)}(m+1)$, we have

$$
B^{p_{(k)}-1}=\sum_{j=0}^{q_{(k)}-1} B^{p_{(k)}-1} C^{j}\binom{p_{(k)}-1+j}{p_{(k)}-1}\binom{k-m\left(p_{(k)}-1+1\right)-n j}{p_{(k)}-1+j+1},
$$

where

$$
\binom{k-m\left(p_{(k)}-1+1\right)-n j}{p_{(k)}-1+j+1}=\binom{k-m p_{(k)}-n j}{p_{(k)}+j}=0 \quad \text { if } j>0 .
$$

Thus,

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & B\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right] \\
& +C\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}\right] \\
= & B \mathrm{e}_{m n}^{B C(k-m)}+C \mathrm{e}_{m n}^{B C(k-n)}
\end{aligned}
$$

and formula (22) is proved.
III. $\left(p_{(k)}-1\right)(m+1)+1 \leq k<p_{(k)}(m+1) \wedge k=q_{(k)}(n+1)$

In this case, we have (see relevant computations in cases I and II)

$$
p_{(k-m)}=p_{(k)}-1, \quad p_{(k+1)}=p_{(k)}
$$

and

$$
q_{(k-n)}=q_{(k)}, \quad q_{(k+1)}=q_{(k)}+1 .
$$

Then

$$
\begin{aligned}
& \mathrm{e}_{m n}^{B C k}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1}, \\
& \mathrm{e}_{m n}^{B C(k+1)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}} B^{i} C^{j}\binom{i+j}{i}\binom{k+1-m i-n j}{i+j+1}
\end{aligned}
$$

and

$$
\begin{align*}
& \mathrm{e}_{m n}^{B C(k-m)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-2, q_{(k-m)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m-m i-n j}{i+j+1},  \tag{23}\\
& \mathrm{e}_{m n}^{B C(k-n)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k-n)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-n-m i-n j}{i+j+1} . \tag{24}
\end{align*}
$$

Like with the computations performed in case I, we can get

$$
\binom{k-m-m i-n j}{i+j+1}=0 \quad \text { if } i \geq 0, j \geq q_{(k-m)}
$$

and

$$
\binom{k-n-m i-n j}{i+j+1}=0 \quad \text { if } i \geq p_{(k-n)}, j \geq 0 .
$$

So, we can substitute $q_{(k)}$ for $q_{(k-m)}$ in (23) and $p_{(k)}$ for $p_{(k-n)}$ in (24).

Thus, we have

$$
\begin{align*}
& \mathrm{e}_{m n}^{B C(k-m)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1},  \tag{25}\\
& \mathrm{e}_{m n}^{B C(k-n)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}^{-1, q_{(k)}-1}} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1} . \tag{26}
\end{align*}
$$

It is easy to see that, due to (3), formula (25) can also be used instead of (23) if $q_{(k-m)}<1$ and formula (26) can also be used instead of (24) if $p_{(k-n)}<1$.

Now we have to prove

$$
\begin{align*}
\Delta \mathrm{e}_{m n}^{B C k}= & B \mathrm{e}_{m n}^{B C(k-m)}+C \mathrm{e}_{m n}^{B C(k-n)} \\
= & B\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right] \\
& +C\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}\right] . \tag{27}
\end{align*}
$$

Considering the difference by its definition, we get

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & \mathrm{e}_{m n}^{B C(k+1)}-\mathrm{e}_{m n}^{B C k} \\
= & I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}} B^{i} C^{j}\binom{i+j}{i}\binom{k+1-m i-n j}{i+j+1} \\
& -I-(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1} \\
= & (B+C)\left(\sum_{i=0, j=0}^{p_{(k)}^{-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\left[\binom{k+1-m i-n j}{i+j+1}-\binom{k-m i-n j}{i+j+1}\right]}\right. \\
& \left.+\sum_{i=0}^{p_{(k)}-1} B^{i} C^{q_{(k)}}\binom{i+q_{(k)}}{i}\binom{k+1-m i-n q_{(k)}}{i+q_{(k)}+1}\right] .
\end{aligned}
$$

With the aid of relation $k=q_{(k)}(n+1)$, we get

$$
\binom{k+1-m i-n q_{(k)}}{i+q_{(k)}+1}=\binom{q_{(k)}+1-m i}{q_{(k)}+1+i}=0 \quad \text { if } i>0
$$

and

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k} & =(B+C)\left[\sum_{i=0, j=0}^{p_{(k)}^{-1, q_{(k)}-1}} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j}+C^{q(k)}\right] \\
& =(B+C)\left[I+\sum_{i=1}^{p_{(k)}-1} B^{i} C^{0}\binom{i}{i}\binom{k-m i}{i}\right.
\end{aligned}
$$

$$
\begin{aligned}
& +\sum_{j=1}^{q_{(k)}-1} B^{0} C^{j}\binom{j}{0}\binom{k-n j}{j} \\
& \left.+\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j}+C^{q_{(k)}}\right] .
\end{aligned}
$$

By (10), we have

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & (B+C)\left[I+\sum_{i=1}^{p_{(k)}-1} B^{i} C^{0}\binom{i-1}{i-1}\binom{k-m i}{i}+\sum_{j=1}^{q_{(k)}-1} B^{0} C^{j}\binom{j-1}{0}\binom{k-n j}{j}\right. \\
& +\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i-1}\binom{k-m i-n j}{i+j} \\
& \left.+\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i}\binom{k-m i-n j}{i+j}+C^{q_{(k)}}\right] \\
= & (B+C)\left[I+\sum_{i=1, j=0}^{p_{(k)}^{-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i-1}\binom{k-m i-n j}{i+j}}\right. \\
& \left.+\sum_{i=0, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i}\binom{k-m i-n j}{i+j}+C^{q_{(k)}}\right] .
\end{aligned}
$$

Now we replace in the first sum the summation index $i$ by $i+1$ and in the second sum we replace the summation index $j$ by $j+1$. Then

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & (B+C)\left[I+\sum_{i=0, j=0}^{p_{(k)}^{-2, q_{(k)}-1}} B^{i+1} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right. \\
& \left.+\sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j+1}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}+C^{q_{(k)}}\right] \\
= & B+B(B+C) \sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1} \\
& +C+C(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}+C^{q_{(k)}(B+C)} \\
= & B\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right] \\
& +C\left[I+(B+C)\left(\sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}+C^{q_{(k)}-1}\right)\right] .
\end{aligned}
$$

For $k=q_{(k)}(n+1)$, we have

$$
C^{q_{(k)}-1}=\sum_{i=0}^{p_{(k)}-1} B^{i} C^{q_{(k)}-1}\binom{i+q_{(k)}-1}{i}\binom{k-m i-n\left(q_{(k)}-1+1\right)}{i+q_{(k)}-1+1},
$$

where

$$
\binom{k-m i-n\left(q_{(k)}-1+1\right)}{i+q_{(k)}-1+1}=\binom{k-m i-n q_{(k)}}{i+q_{(k)}}=0 \quad \text { if } i>0 .
$$

Thus,

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & B\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}^{-2, q_{(k)}-1}} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right] \\
& +C\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}\right] \\
= & B \mathrm{e}_{m n}^{B C(k-m)}+C \mathrm{e}_{m n}^{B C(k-n)}
\end{aligned}
$$

and formula (27) is proved.
IV. $k=p_{(k)}(m+1) \wedge k=q_{(k)}(n+1)$

In this case, we have (see similar combinations in cases II and III)

$$
p_{(k-m)}=p_{(k)}, \quad p_{(k+1)}=p_{(k)}+1
$$

and

$$
q_{(k-n)}=q_{(k)}, \quad q_{(k+1)}=q_{(k)}+1 .
$$

Then

$$
\begin{aligned}
& \mathrm{e}_{m n}^{B C k}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1}, \\
& \mathrm{e}_{m n}^{B C(k+1)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}, q_{(k)}} B^{i} C^{j}\binom{i+j}{i}\binom{k+1-m i-n j}{i+j+1}
\end{aligned}
$$

and

$$
\begin{align*}
& \mathrm{e}_{m n}^{B C(k-m)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k)}^{-1, q_{(k-m)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m-m i-n j}{i+j+1},}  \tag{28}\\
& \mathrm{e}_{m n}^{B C(k-n)}=I+(B+C) \sum_{i=0, j=0}^{p_{(k-n)}^{-1, q_{(k)}-1}} B^{i} C^{j}\binom{i+j}{i}\binom{k-n-m i-n j}{i+j+1} . \tag{29}
\end{align*}
$$

As before,

$$
\binom{k-m-m i-n j}{i+j+1}=0 \quad \text { if } i \geq 0, j \geq q_{(k-m)}
$$

and

$$
\binom{k-n-m i-n j}{i+j+1}=0 \quad \text { if } i \geq p_{(k-n)}, j \geq 0
$$

So, we can substitute $q_{(k)}$ for $q_{(k-m)}$ in (28) and $p_{(k)}$ for $p_{(k-n)}$ in (29) and

$$
\begin{align*}
& \mathrm{e}_{m n}^{B C(k-m)}=I+(B+C) \sum_{i=0, j=0}^{p(k)-1, q(k)-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1},  \tag{30}\\
& \mathrm{e}_{m n}^{B C(k-n)}=I+(B+C) \sum_{i=0, j=0}^{p(k)-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1} . \tag{31}
\end{align*}
$$

It is easy to see that, due to (3), formula (30) can also be used instead of (28) if $q_{(k-m)}<1$ and formula (31) can also be used instead of (29) if $p_{(k-n)}<1$.

Now it is possible to prove the formula

$$
\begin{align*}
\Delta \mathbf{e}_{m n}^{B C k}= & B \mathbf{e}_{m n}^{B C(k-m)}+C \mathrm{e}_{m n}^{B C(k-n)} \\
= & B\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right] \\
& +C\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}\right] . \tag{32}
\end{align*}
$$

By definition, we get

$$
\begin{aligned}
& \Delta \mathrm{e}_{m n}^{B C k}=\mathrm{e}_{m n}^{B C(k+1)}-\mathrm{e}_{m n}^{B C k} \\
& =I+(B+C) \sum_{i=0, j=0}^{p(k), q_{(k)}} B^{i} C^{j}\binom{i+j}{i}\binom{k+1-m i-n j}{i+j+1} \\
& -I-(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j+1} \\
& =(B+C)\left[\sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\left[\binom{k+1-m i-n j}{i+j+1}-\binom{k-m i-n j}{i+j+1}\right]\right. \\
& +\sum_{j=0}^{q_{(k)}} B^{p_{(k)}} C^{j}\binom{p_{(k)}+j}{p_{(k)}}\binom{k+1-m p_{(k)}-n j}{p_{(k)}+j+1} \\
& \left.+\sum_{i=0}^{p_{(k)}} B^{i} C^{q_{(k)}}\binom{i+q_{(k)}}{i}\binom{k+1-m i-n q_{(k)}}{i+q_{(k)}+1}\right] .
\end{aligned}
$$

With the aid of equations $k=p_{(k)}(m+1), k=q_{(k)}(n+1)$, we get

$$
\begin{aligned}
& \binom{k+1-m p_{(k)}-n j}{p_{(k)}+j+1}=\binom{p_{(k)}+1-n j}{p_{(k)}+1+j}=0 \quad \text { if } j>0, \\
& \binom{k+1-m i-n q_{(k)}}{i+q_{(k)}+1}=\binom{q_{(k)}+1-m i}{q_{(k)}+1+i}=0
\end{aligned}
$$

and

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & (B+C)\left[\sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j}+B^{p_{(k)}}+C^{q_{(k)}}\right] \\
= & (B+C)\left[I+\sum_{i=1}^{p_{(k)}-1} B^{i} C^{0}\binom{i}{i}\binom{k-m i}{i}+\sum_{j=1}^{q_{(k)}-1} B^{0} C^{j}\binom{j}{0}\binom{k-n j}{j}\right. \\
& +\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n j}{i+j}+B^{\left.p_{(k)}+C^{q_{(k)}}\right] .}
\end{aligned}
$$

By (10), we have

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & (B+C)\left[I+\sum_{i=1}^{p_{(k)}-1} B^{i} C^{0}\binom{i-1}{i-1}\binom{k-m i}{i}+\sum_{j=1}^{q_{(k)}-1} B^{0} C^{j}\binom{j-1}{0}\binom{k-n j}{j}\right. \\
& +\sum_{i=1, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i-1}\binom{k-m i-n j}{i+j} \\
& +\sum_{i=1, j=1}^{p_{(k)}^{-1, q_{(k)}-1}} B^{i} C^{j}\binom{i+j-1}{i}\binom{k-m i-n j}{i+j}+B^{\left.p_{(k)}+C^{q_{(k)}}\right]} \\
= & (B+C)\left[I+\sum_{i=1, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i-1}\binom{k-m i-n j}{i+j}\right. \\
& \left.+\sum_{i=0, j=1}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j-1}{i}\binom{k-m i-n j}{i+j}+B^{p_{(k)}}+C^{q_{(k)}}\right] .
\end{aligned}
$$

We replace in the first sum the summation index $i$ by $i+1$ and in the second sum we substitute the summation index $j$ by $j+1$. Then

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & (B+C)\left[I+\sum_{i=0, j=0}^{p_{(k)}^{-2, q_{(k)}-1}} B^{i+1} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right. \\
& \left.+\sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j+1}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}+B^{p_{(k)}}+C^{q_{(k)}}\right] \\
= & B+B(B+C) \sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}+B^{p_{(k)}(B+C)} \\
& +C+C(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}+C^{q_{(k)}(B+C)} \\
= & B\left[I+(B+C)\left(\sum_{i=0, j=0}^{p_{(k)}-2, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}+B^{p_{(k)}-1}\right)\right] \\
& +C\left[I+(B+C)\left(\sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-2} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}+C^{q_{(k)}-1}\right)\right] .
\end{aligned}
$$

Because $k=p_{(k)}(m+1)=q_{(k)}(n+1)$, we can express $B^{p_{(k)}-1}$ and $C^{q_{(k)}-1}$ in the form

$$
\begin{aligned}
& B^{p_{(k)}-1}=\sum_{j=0}^{q_{(k)}-1} B^{p_{(k)}-1} C^{j}\binom{p_{(k)}-1+j}{p_{(k)}-1}\binom{k-m\left(p_{(k)}-1+1\right)-n j}{p_{(k)}-1+j+1}, \\
& C^{q_{(k)}-1}=\sum_{i=0}^{p_{(k)}-1} B^{i} C^{q_{(k)}-1}\binom{i+q_{(k)}-1}{i}\binom{k-m i-n\left(q_{(k)}-1+1\right)}{i+q_{(k)}-1+1},
\end{aligned}
$$

where

$$
\begin{aligned}
& \binom{k-m\left(p_{(k)}-1+1\right)-n j}{p_{(k)}-1+j+1}=\binom{k-m p_{(k)}-n j}{p_{(k)}+j}=0 \quad \text { if } j>0, \\
& \binom{k-m i-n\left(q_{(k)}-1+1\right)}{i+q_{(k)}-1+1}=\binom{k-m i-n q_{(k)}}{i+q_{(k)}}=0 \quad \text { if } i>0 .
\end{aligned}
$$

Thus,

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C k}= & B\left[I+(B+C) \sum_{i=0, j=0}^{p(k)-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m(i+1)-n j}{i+j+1}\right] \\
& +C\left[I+(B+C) \sum_{i=0, j=0}^{p_{(k)}-1, q_{(k)}-1} B^{i} C^{j}\binom{i+j}{i}\binom{k-m i-n(j+1)}{i+j+1}\right] \\
= & B \mathrm{e}_{m n}^{B C(k-m)}+C \mathrm{e}_{m n}^{B C(k-n)} .
\end{aligned}
$$

Therefore, formula (32) is valid.
We proved that formula (6) holds in each of the considered cases I, II, III and IV for $k \geq 1$. If $k=0$, the proof can be done directly because $p_{(0)}=q_{(0)}=0, p_{(1)}=q_{(1)}=1$,

$$
\begin{aligned}
\Delta \mathrm{e}_{m n}^{B C 0}= & \mathrm{e}_{m n}^{B C 1}-\mathrm{e}_{m n}^{B C 0} \\
= & I+(B+C) \sum_{i=0, j=0}^{0,0} B^{i} C^{j}\binom{i+j}{i}\binom{1-m i-n j}{i+j+1} \\
& -I-(B+C) \sum_{i=0, j=0}^{-1,-1} B^{i} C^{j}\binom{i+j}{i}\binom{-m i-n j}{i+j+1}=I+B+C-I=B+C
\end{aligned}
$$

and

$$
B \mathrm{e}_{m n}^{B C(-m)}+C \mathrm{e}_{m n}^{B C(-n)}=B I+C I=B+C .
$$

Formula (6) holds again. Theorem 2 is proved.

## Open problems and concluding remarks

Formula (4) is valid for $k \in \mathbb{Z}_{-m}^{\infty}$. However, formula (6) holds for $k \in \mathbb{Z}_{0}^{\infty}$ only. Therefore, there is a difference between the definition domains of the formulas, and it is a challenge how to modify Definition 2 of discrete matrix delayed exponential for two delays in such
a way that formula (6) will hold for $k \in \mathbb{Z}_{-\max \{m, n\} \text {. In [1] formula (4) is used to get a repre- }}^{\infty}$ sentation of the solution of the problems (both homogeneous and nonhomogeneous)

$$
\begin{aligned}
& \Delta y(k)=B y(k-m)+f(k), \quad k \in \mathbb{Z}_{0}^{\infty}, \\
& y(k)=\varphi(k), \quad k \in \mathbb{Z}_{-m}^{0},
\end{aligned}
$$

where $f: \mathbb{Z}_{0}^{\infty} \rightarrow \mathbb{R}^{r}, y: \mathbb{Z}_{-m}^{\infty} \rightarrow \mathbb{R}^{r}$ and $\varphi: \mathbb{Z}_{-m}^{0} \rightarrow \mathbb{R}^{r}$.
It is an open problem how to use formula (6) to get a representation of the solution of the homogeneous and nonhomogeneous problems

$$
\begin{aligned}
& \Delta y(k)=B y(k-m)+C y(k-n)+f(k), \quad k \in \mathbb{Z}_{0}^{\infty}, \\
& y(k)=\varphi(k), \quad k \in \mathbb{Z}_{-s}^{0}, s=\max \{m, n\}
\end{aligned}
$$

if $B C=C B$.
Let us note that the first concept of matrix delayed exponential was given in [3] and the first concept of discrete matrix delayed exponential was given in [1]. Further development of the delayed matrix exponentials method and its utilization to various problems can be found, e.g., in [4-16].

## Competing interests

The authors declare that they have no competing interests.

## Authors' contributions

The authors have made the same contribution. Both authors have read and approved the final manuscript.

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## Acknowledgements

The first author was supported by Operational Programme Research and Development for Innovations, No CZ.1.05/2.1.00/03.0097, as an activity of the regional Centre AdMaS.

Received: 23 January 2013 Accepted: 29 April 2013 Published: 14 May 2013

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[^1]:    doi:10.1186/1687-1847-2013-139
    Cite this article as: Diblík and Morávková: Discrete matrix delayed exponential for two delays and its property. Advances in Difference Equations 2013 2013.139

