

A “three-sentence proof” of Hansson’s theorem

Henrik Petri¹

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Abstract We provide a new proof of Hansson’s theorem: every preorder has a complete preorder extending it. The proof boils down to showing that the lexicographic order extends the Pareto order.

Keywords Ordering extension theorem · Lexicographic order · Pareto order · Preferences

JEL Classification C65 · D01

1 Introduction

Two extensively studied binary relations in economics are the Pareto order and the lexicographic order. It is a well-known fact that the latter relation is an ordering extension of the former. For instance, in [Petri and Voorneveld \(2016\)](#), an essential ingredient is Lemma 3.1, which roughly speaking requires the order under consideration to be an extension of the Pareto order. The main message of this short note is that some fundamental order extension theorems can be reduced to this basic fact. An advantage of the approach is that it seems less abstract than conventional proofs and hence may offer a pedagogical advantage in terms of exposition. [Mandler \(2015\)](#) gives an elegant proof of Spzilaraj’s theorem ([1930](#)) that stresses the importance of the lexicographic

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✉ Henrik Petri
henrik.petri@hhs.se

¹ Department of Finance, Stockholm School of Economics, Box 6501, 113 83 Stockholm, Sweden

approach in proving ordering extension theorems. He shows that criteria can be built out of the relation \succ on the domain X in such a way that if the criteria are ordered lexicographically the corresponding relation extends \succ . At a technical level the proof presented here is quite similar. We use a result from [Evren and Ok \(2011\)](#), that uses criteria defined in a similar way as in [Mandler \(2015\)](#). However, in terms of exposition we argue that our proof is easier. Another difference compared to [Mandler \(2015\)](#) is that he gives a proof of Szpilrajn's theorem (1930), whereas this note presents a proof of Hansson's theorem (1968). Of course given a proof of one of them the other follows quite easily, but we believe that our proof of Hansson's theorem is easier to digest. Colloquially the proof in this note reads:

“Every preorder is essentially a Pareto order and the lexicographic order extends the Pareto order. Since the lexicographic order is complete every preorder has an ordering extension”.

We thus argue that abstract extension theorems of orders can be embodied in this simple principle.

The note is organized as follows. Section 2 contains notation and preliminaries. In Sect. 3, we present our proof of Hansson's theorem. Section 4 concludes. We gather some standard results in set theory in Appendix A.

2 Notation and preliminaries

As usual \mathbb{R} denotes the set of real numbers. A binary relation/order \succsim on a set X , is a subset $\succsim \subset X \times X$. If $(x, y) \in \succsim$ we write this as $x \succsim y$. An order \succsim on X is:

Reflexive If $x \succsim x$ for all $x \in X$.

Transitive If $x \succsim y, y \succsim z$ implies $x \succsim z$ for all $x, y, z \in X$.

Complete If $x \succsim y$ or $y \succsim x$ for all $x, y \in X$.

Antisymmetric If $x \succsim y$ and $y \succsim x$ implies $x = y$ for all $x, y \in X$.

A preorder \succsim is a reflexive, transitive order. If a preorder \succsim is complete, then we call \succsim a complete preorder. An order \succsim is a *linear order* if \succsim is antisymmetric, transitive and complete. Given a linear order \succsim on X and a subset A of X , $a \in A$ is a smallest element if $x \succsim a$ for all $x \in A$. A linear order on a set X is a well order if every nonempty subset A of X contains a smallest element. The *well ordering theorem* states that every set X admits a binary relation \succsim such that X with \succsim is well ordered.

If \succsim is an order we define a relation \succ on X by $x \succ y$ if and only if $x \succsim y$ and not $y \succsim x$. We also define an order \sim on X by $x \sim y$ if and only if $x \succsim y$ and $y \succsim x$. A complete preorder \succsim' on X *extends* (or is an ordering extension of) a preorder \succsim on X if for all $x, y \in X$: $x \succsim y$ implies that $x \succsim' y$ and $x \succ y$ implies that $x \succ' y$.

Given a set I and for each $i \in I$ a complete preorder \succsim_i on X the *Pareto order* \geq on X is defined by for all $x, y \in X$: $x \geq y$ if and only if $x \succsim_i y$ for all $i \in I$. We will sometimes refer to \succsim_i as a *coordinate relation*.

Let I be a set with linear order \leq on I and a collection of complete preorders \succsim_i on X for all $i \in I$. Define a *lexicographic relation* \geq_L on X by $x \geq_L y$ if and only if $x \succsim_i y$ for all $i \in I$, or there is a $j \in I$ such that $x \succ_j y$ and $x \sim_i y$ for all $i \in I$

with $i < j$. It is a standard fact in basic set theory that if \leq is a well order then \geq_L is a complete preorder on X . For example, [Mandler \(2015\)](#) alludes to this result, see also [Ciesielski \(1997\)](#). For completeness a proof is presented in Lemma [A.1](#) in the Appendix A.

3 The proof

We now give our proof of Hansson’s theorem. A first crucial ingredient in our proof is a corollary to a result in [Evren and Ok \(2011\)](#). It shows that every preorder essentially is a Pareto order. The result has the same universal character as a representation result by [Chipman \(1960\)](#), which shows that every complete preorder essentially is a lexicographic order:

Lemma 3.1 *For every preorder \succsim there is a Pareto order \geq with coordinate relations \succsim_i for all $i \in I$ such that $x \succsim y$ if and only if $x \geq y$ for all $x, y \in X$.*

Proof Let $I = X$ and for each $x \in X$ let $u_x(y) := \mathbf{1}_{\{z \in X \mid z \succsim_x y\}}(y)$ for all $y \in X$ (where $\mathbf{1}_A$ denotes the indicator function of a set A). Define a relation \succsim_x by $z \succsim_x y$ if and only if $u_x(z) \geq u_x(y)$. The result follows by [Evren and Ok \(2011, Proposition 1\)](#). \square

Another important observation in our proof is that the lexicographic order extends the Pareto order. This observation is recorded as Lemma [3.2](#).

Lemma 3.2 *Let \geq be a Pareto order on X with coordinate relations \succsim_i . Then the lexicographic order \geq_L on X with coordinate relations \succsim_i for all $i \in I$ extends \geq .*

Proof Let $x, y \in X$ with $x \geq y$. If $y >_L x$ then there is an $i \in I$ such that $y \succ_i x$, contradicting that $x \succsim_i y$ for all $i \in I$. Let $x, y \in X$ with $x > y$. Then $x \succ_j y$ for some $j \in I$ and $x \succsim_i y$ for all $i \in I$. If $y \geq_L x$ then either $x \sim_i y$ for all $i \in I$ or $y \succ_i x$ for some $i \in I$, a contradiction. \square

We are now ready for our proof of Hansson’s theorem:

Theorem 3.3 *Let \succsim be a preorder order on X . Then there exists a complete preorder \succsim' on X extending \succsim .*

Proof By Lemma [3.1](#) there is a Pareto relation \geq with coordinate relations \succsim_i such that $x \geq y$ if and only if $x \succsim y$ for all $x, y \in X$. Well order I by \leq and let \geq_L be the lexicographic relation with coordinate relations \succsim_i for all $i \in I$. Then \geq_L is a complete preorder by Lemma [A.1](#) and Lemma [3.2](#) implies that \geq_L extends \succsim . \square

4 Concluding remarks

A short “three sentences” proof of Hansson’s theorem is given. Some of the advantages of the approach are:

- (a) The theorem follows from intuitively plausible and simple principles. Once it is understood that the lexicographic order extends the Pareto order, the rest of the proof follows smoothly.

- (b) The use of the well ordering theorem is transparent. We only had to use the well ordering theorem once in the second sentence of Theorem 3.3. Hence if I is some set that is known to be well ordered (like the set of natural numbers or a finite set), we see that the proof of Hansson's theorem follows without invoking the well ordering Theorem 3.3.

Finally, some readers may object to the claim made in the title of this note, namely that the proof of Theorem 3.3 is only three sentences long. To prove the main result we did indeed have to invoke as many as three lemmas. However, it is our belief that the prerequisites are either well known or otherwise useful to know, and given that, the proof of Hansson's theorem is short.

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A Appendix: Lemma A.1

Lemma A.1 *Let I be an arbitrary set and \leq a well order on I . For each $i \in I$ let \succsim_i be a complete preorder on X . Then the lexicographic order \geq_L on X , with order \leq on I and relations \succsim_i on X for all $i \in I$, is a complete preorder on X .*

Proof Assume I is well ordered by \leq . For all $x, y \in X$, let $N(x, y) := \{i \in I : x \succ_i y\} \cup \{i \in I : y \succ_i x\}$. We show that \geq_L is complete and transitive.

Totality: Let $x, y \in X$. If $x \sim_i y$ for all $i \in I$ then $x \geq_L y$. Otherwise $x \succ_i y$ or $y \succ_i x$ for some $i \in I$ and hence $N(x, y) \neq \emptyset$. Since I is well ordered there is a smallest $j \in I$ such that $x \succ_j y$ or $y \succ_j x$ and hence $x \sim_i y$ for all $i < j$. If $x \succ_j y$ ($y \succ_j x$) it follows that $x \geq_L y$ ($y \geq_L x$).

Transitivity Let $x, y, z \in X$, $x \geq_L y$ and $y \geq_L z$. W.l.o.g. assume that $x >_L y$ and $y >_L z$ (the other cases are either trivial or follows by using similar arguments as in the present case). Then $N(x, y) \neq \emptyset$ and $N(y, z) \neq \emptyset$. Hence there is a smallest $j \in N(x, y)$ and a smallest $j' \in N(y, z)$. If $j = j'$, then we are done. Assume $j < j'$ then $x \succ_j y \sim_j z$ and $x \sim_i y \sim_i z$ for all $i < j$. Thus $x \geq_L z$. If $j' < j$ then $x \sim_j y \succ_{j'} z$ and $x \sim_i y \sim_i z$ for all $i < j'$ and hence $x \geq_L z$. \square

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