



Exponential tractability of L_2 -approximation with function values

David Krieg^{1,2} · Paweł Siedlecki³ · Mario Ullrich¹ · Henryk Woźniakowski^{3,4}

Received: 31 May 2022 / Accepted: 7 February 2023 / Published online: 7 March 2023
© The Author(s) 2023

Abstract

We study the complexity of high-dimensional approximation in the L_2 -norm when different classes of information are available; we compare the power of function evaluations with the power of arbitrary continuous linear measurements. Here, we discuss the situation when the number of linear measurements required to achieve an error $\varepsilon \in (0, 1)$ in dimension $d \in \mathbb{N}$ depends only poly-logarithmically on ε^{-1} . This corresponds to an exponential order of convergence of the approximation error, which often happens in applications. However, it does not mean that the high-dimensional approximation problem is easy, the main difficulty usually lies within the dependence on the dimension d . We determine to which extent the required amount of information changes if we allow only function evaluation instead of arbitrary linear information. It turns out that in this case we only lose very little, and we can even restrict to linear algorithms. In particular, several notions of tractability hold simultaneously for both types of available information.

Communicated by: Holger Rauhut

✉ Paweł Siedlecki
psiedlecki@mimuw.edu.pl

David Krieg
david.krieg@jku.at

Mario Ullrich
mario.ullrich@jku.at

Henryk Woźniakowski
hwozniak@mimuw.edu.pl

- ¹ Institut für Analysis, Johannes Kepler Universität, Linz, Austria
- ² Johann Radon Institute for Computational and Applied Mathematics (RICAM), Austrian Academy of Sciences, Linz, Austria
- ³ Faculty of Mathematics, Informatics and Mechanics, University of Warsaw, Warsaw, Poland
- ⁴ Department of Computer Science, Columbia University, New York, USA

Keywords Approximation · Multivariate problems · Tractability · Complexity

Mathematics Subject Classification (2010) 65Y20 · 41A25 · 41A65 · 41A63

1 Exposition

We want to approximate real- or complex-valued functions defined on some (nonempty) set \mathcal{D} , and belonging to a space F . We assume that F is a separable Banach space of functions defined on \mathcal{D} , such that function evaluation $f \mapsto f(x)$ is continuous on F for each $x \in \mathcal{D}$ and F is continuously embedded in $L_2 = L_2(\mathcal{D}, \mu)$ for some measure μ . Formally, the approximation problem is given as

$$\text{APP} : F \rightarrow L_2, \quad \text{APP}(f) := f,$$

which might be understood as a continuous embedding into L_2 . The class of all spaces F satisfying the assumptions above will be denoted by \mathcal{A} . In particular, for each $F \in \mathcal{A}$ we have some associated nonempty set \mathcal{D} , measure μ on \mathcal{D} and continuous embedding APP.

We approximate APP by using functionals from the class Λ^{std} consisting of all function evaluations, or from the class $\Lambda^{\text{all}} = F^*$ of all continuous linear functionals.

Below B_F denotes the closed unit ball in F . Let us define, for $n \in \mathbb{N}$, the

- *n*-th linear sampling width as

$$e_n(F, L_2) := \inf_{\substack{x_1, \dots, x_n \in \mathcal{D} \\ \varphi_1, \dots, \varphi_n \in L_2}} \sup_{f \in B_F} \left\| f - \sum_{i=1}^n f(x_i) \varphi_i \right\|_{L_2},$$

- *n*-th sampling width as

$$g_n(F, L_2) := \inf_{\substack{x_1, \dots, x_n \in \mathcal{D} \\ \phi : \mathbb{R}^n \rightarrow L_2}} \sup_{f \in B_F} \left\| f - \phi(f(x_1), \dots, f(x_n)) \right\|_{L_2},$$

- *n*-th linear width as

$$a_n(F, L_2) := \inf_{\substack{T : L_2 \rightarrow L_2 \\ \text{rank}(T) \leq n}} \sup_{f \in B_F} \|f - Tf\|_{L_2},$$

- *n*-th Gelfand width as

$$c_n(F, L_2) := \inf_{\substack{\phi : \mathbb{R}^n \rightarrow L_2 \\ N \in (F^*)^n}} \sup_{f \in B_F} \|f - \phi \circ N(f)\|_{L_2}.$$

These quantities represent the *minimal worst case errors* that can be achieved with linear or nonlinear algorithms using at most n function values or linear measurements, respectively.

We also define the information-based complexity of the problem APP for the classes Λ^{std} and Λ^{all} , respectively, as the minimal number of evaluations from Λ^{std} or Λ^{all} necessary to obtain the absolute precision of approximation at most ε , i.e., as

$$n^{\text{std}}(\varepsilon, F) := \min \{n : g_n(F, L_2) \leq \varepsilon\}$$

and

$$n^{\text{all}}(\varepsilon, F) := \min \{n : c_n(F, L_2) \leq \varepsilon\}.$$

Note that, since $g_n(F, L_2) \leq e_n(F, L_2)$, we have

$$n^{\text{std}}(\varepsilon, F) \leq \min \{n : e_n(F, L_2) \leq \varepsilon\} =: n^{\text{std-lin}}(\varepsilon, F),$$

and all our upper bounds are proven for $n^{\text{std-lin}}(\varepsilon, F)$. There is a lot of literature on the size of these quantities for specific classes F . We refer to the monographs [10, 33, 35, 36, 42, 43, 45] for more details and literature on the subject.

Here, we are specifically interested in the comparison of these quantities for general classes F . That is, since $n^{\text{all}}(\varepsilon, F) \leq n^{\text{std}}(\varepsilon, F)$ is obvious for all $F \in \mathcal{A}$, we ask for an upper bound on $n^{\text{std}}(\varepsilon, F)$ based on knowledge of the function $n^{\text{all}}(\varepsilon, F)$. However, it is known that such a bound cannot hold without certain assumptions on F , see [36, Chapter 26] and references therein, and even then, the involved ‘‘constants’’ depend in a non-trivial way on F . One approach to obtain qualitative statements on the relation of the complexities is to consider a whole sequence of spaces $(F_d)_{d \in \mathbb{N}}$, where d can be interpreted as the dimension of the underlying domain. We then assume a certain bound on $n^{\text{all}}(\varepsilon, F_d)$, depending only on $\varepsilon \in (0, 1)$ and $d \in \mathbb{N}$, and ask for an upper bound on $n^{\text{std}}(\varepsilon, F_d)$, hopefully not much worse than the bound on $n^{\text{all}}(\varepsilon, F_d)$.

In the present paper, we allow arbitrary Banach spaces of functions F_d , but we assume that $n^{\text{all}}(\varepsilon, F_d)$ depends only poly-logarithmically on ε^{-1} . That is, we assume that there exist $A_d, B_d > 0$ such that

$$n^{\text{all}}(\varepsilon, F_d) \leq A_d \left(1 + \ln \varepsilon^{-1}\right)^{B_d} \quad \text{for all } 0 < \varepsilon \leq 1, \tag{1}$$

and study how this translates into bounds on $n^{\text{std}}(\varepsilon, F_d)$. Note that the above bound (1) on the complexity implies that

$$c_n(F_d, L_2) \leq e \exp(-n/A_d)^{1/B_d} \quad \text{for all } n \geq A_d, \tag{2}$$

whereas (2) implies that (1) holds with $+1$ added on the right hand side. The assumption (1) is therefore equivalent to the existence of a (possibly nonlinear) algorithm based on arbitrary linear information that converges exponentially fast. We will show that in this case, we do not lose much when we only allow linear algorithms and function evaluations as information. One of our main results may be stated as follows.

Theorem (see Corollary 5) *Assume that $F_d \in \mathcal{A}$ for every $d \in \mathbb{N}$ and*

$$n^{\text{all}}(\varepsilon, F_d) \leq c d^q (1 + \ln \varepsilon^{-1})^p$$

for some $p, c > 0, q \geq 0$, and all $\varepsilon \in (0, 1)$. Then

$$n^{\text{std-lin}}(\varepsilon, F_d) \leq C d^q (1 + \ln d)^p (1 + \ln \varepsilon^{-1})^p$$

for all $\varepsilon \in (0, 1)$ and $d \in \mathbb{N}$, and some $C > 0$ that depends only on c, p and q .

This shows that every Banach space that is assumed to be *approximable in high dimensions* (in the above sense) with an exponential order by some algorithm and information can practically already be treated with linear algorithms based on function values. In particular, this improves upon Theorem 26.21 from [36] and solves Open Problem 128 therein. Let us add that we do not know if the additional $(1 + \ln d)^p$ is necessary.

There are many appearances of the assumption (1) in the literature. Besides the detailed study of certain weighted Hilbert spaces of analytic functions [8, 24, 25, 30, 47], it appears naturally in the context of approximation with (increasingly flat) Gaussian kernels [12, 17, 26, 41], or in tensor product approximations [14, 16], or for certain smoothness spaces on complex spheres [7]. Moreover, it is a typical assumption for the construction of greedy bases [4, 5, 15]. Let us also add that there is quite some study on the *stability* of algorithms that can achieve an exponential convergence, see [1–3, 39] for details.

When it comes to the study of the *tractability* of the problem, i.e., the precise dependence of the error on the dimension, especially when we only allow function evaluations, there is much less to cite and we are only aware of the Hilbert space references from above. As an explicit example, let us mention the Gaussian space on \mathbb{R}^d with reproducing kernel $K(x, y) = \exp(-\|x - y\|_2^d)$, which satisfies a relation of the form (1) for L_2 -approximation with respect to the Gaussian probability measure μ , see [41]. In the Hilbert case, there are some general results which make the situation somewhat simpler. For example, it is known that linear algorithms are always optimal and one may work with the singular value decomposition of the embedding APP. We refer to [33, Chapter 4] and [36, Chapter 26].

A bit more is known in the case of *algebraic tractability*, i.e., when the complexity depends polynomially on ε^{-1} instead of $\ln \varepsilon^{-1}$. In addition to general Hilbert space results, see [36, Chapter 26], and characterizations for weighted Korobov spaces, see [11], there are also quite sharp results for the classical smoothness spaces $C^k(\Omega_d)$ of k -times differentiable functions on certain d -dimensional domains, possibly for $k = \infty$. See [27, 34, 46] for details on approximation, and [20–23] for numerical integration in the same classes. However, a comparison as proven here in the case of exponential convergence is not possible in this case, see the end of Section 2. In any case, it is open to determine the precise behavior of $n^{\text{std}}(\varepsilon, F_d)$ for most classical spaces, while $n^{\text{all}}(\varepsilon, F_d)$ is more often known.

Our results are based on the following (special case of) Theorem 3 from [9], see also [18, 28, 29, 32, 44], which allows us to treat more general classes of functions.

Theorem 1 *For each $0 < r < 2$, there is a universal constant $b \in \mathbb{N}$, depending only on r , such that the following holds. For all $F \in \mathcal{A}$ and $n \geq 2$, we have*

$$e_{bn}(F, L_2) \leq \left(\frac{1}{n} \sum_{k \geq n} a_k(F, L_2)^r \right)^{1/r}.$$

Additionally, we use the following fundamental result from [38], see also [6, 31].

Theorem 2 For all $F \in \mathcal{A}$ and $n \geq 1$, we have

$$a_n(F, L_2) \leq (1 + \sqrt{n}) c_n(F, L_2).$$

Remark 3 We would like to stress that the proof of Theorem 1 in [9] is non-constructive, and we do not know how to explicitly construct evaluation points $x_1, \dots, x_{bn} \in \mathcal{D}$ together with some $\varphi_1, \dots, \varphi_{bn} \in L_2$ satisfying

$$\sup_{f \in B_F} \left\| f - \sum_{i=1}^{bn} f(x_i) \varphi_i \right\|_{L_2} \leq \left(\frac{1}{n} \sum_{k \geq n} a_k(F, L_2)^r \right)^{1/r}.$$

However, for problems with known operators T achieving the infimum as in the definition of linear widths $a_n(F, L_2)$, we are able to specify algorithms utilizing i.i.d. sampling of evaluation points from some known distribution, and satisfying inequalities similar to the one above with high probability, see Theorem 8 of [29].

2 Exponential tractability of approximation

The notions of tractability are defined as follows. Let us fix, for every $d \in \mathbb{N}$, some space $F_d \in \mathcal{A}$. For each $F_d \in \mathcal{A}$ we have some associated set \mathcal{D}_d equipped with a measure μ_d , and a continuous embedding $\text{APP}_d : F_d \rightarrow L_2(\mathcal{D}_d, \mu_d)$. The index $d \in \mathbb{N}$ is an arbitrary parameter, but it usually stands for the dimension of the domain \mathcal{D}_d . A *multivariate approximation problem* is simply a sequence of embeddings

$$\widetilde{\text{APP}} = (\text{APP}_d : F_d \rightarrow L_2(\mathcal{D}_d, \mu_d))_{d \in \mathbb{N}}.$$

Moreover, tractability notions are defined relative to the considered class of information operations, i.e., we can consider tractability for Λ^{std} or Λ^{all} . Therefore, for $x \in \{\text{std}, \text{all}\}$, we say that $\widetilde{\text{APP}}$ is

- *exponentially strongly polynomially tractable* (EXP-SPT) for the class Λ^x if and only if

$$n^x(\varepsilon, F_d) \leq C (1 + \ln \varepsilon^{-1})^p$$

for some $C, p > 0$ and for all $d \in \mathbb{N}$ and $\varepsilon \in (0, 1)$,

- *exponentially polynomially tractable* (EXP-PT) for the class Λ^x if and only if

$$n^x(\varepsilon, F_d) \leq C d^q (1 + \ln \varepsilon^{-1})^p$$

for some $C, p, q > 0$ and for all $d \in \mathbb{N}$ and $\varepsilon \in (0, 1)$,

- *exponentially quasi-polynomially tractable* (EXP-QPT) for the class Λ^x if and only if

$$n^x(\varepsilon, F_d) \leq C \exp(t(1 + \ln d)(1 + \ln(1 + \ln \varepsilon^{-1})))$$

for some $C, t > 0$ and for all $d \in \mathbb{N}$ and $\varepsilon \in (0, 1)$,

- *exponentially uniformly weakly tractable* (EXP-UWT) for the class Λ^x if and only if for all $\alpha, \beta > 0$ we have

$$\lim_{d + \varepsilon^{-1} \rightarrow \infty} \frac{\ln n^x(\varepsilon, F_d)}{d^\alpha + (1 + \ln \varepsilon^{-1})^\beta} = 0,$$

- exponentially weakly tractable (EXP-WT) for the class Λ^x if and only if

$$\lim_{d+\varepsilon^{-1} \rightarrow \infty} \frac{\ln n^x(\varepsilon, F_d)}{d + (1 + \ln \varepsilon^{-1})} = 0.$$

It is easy to see that we have the following logical relation between the tractability notions defined above

$$\text{EXP-SPT} \implies \text{EXP-PT} \implies \text{EXP-QPT} \implies \text{EXP-UWT} \implies \text{EXP-WT}.$$

For a multivariate approximation problem we prove that exponential strong polynomial tractability (EXP-SPT), exponential polynomial tractability (EXP-PT), exponential uniform weak tractability (EXP-UWT) and exponential weak tractability (EXP-WT) for the class Λ^{all} are each equivalent to the corresponding tractability property for the class Λ^{std} . Moreover, exponential quasi-polynomial tractability (EXP-QPT) for Λ^{all} implies exponential uniform weak tractability (EXP-UWT) for Λ^{std} , i.e, the next tractability notion in the tractability hierarchy considered here. Whether the equivalence of exponential quasi-polynomial tractability (EXP-QPT) for the classes Λ^{all} and Λ^{std} holds remains an open problem.

These equivalences are in sharp contrast to the results for algebraic tractability. See, e.g., [19, 37, 40] for examples where the problem is algebraically tractable for Λ^{all} but the curse of dimensionality holds for Λ^{std} . In particular, [37, Example 5] shows that for the tensor product $W_{2,d}^s$ of certain univariate periodic Sobolev spaces, $s > 1/2$, we have QPT for Λ^{all} , but the curse of dimensionality for Λ^{std} .

3 Results

We now present our results. The first results are concerned with EXP-(S)PT and EXP-QPT. Both are direct corollaries of the following theorem.

Theorem 4 Assume that $F \in \mathcal{A}$ satisfies

$$n^{\text{all}}(\varepsilon, F) \leq A (1 + \ln \varepsilon^{-1})^B$$

for some $B > 0$ and $A \geq 1$ and all $\varepsilon \in (0, 1)$. Then

$$n^{\text{std}}(\varepsilon, F) \leq n^{\text{std-lin}}(\varepsilon, F_d) \leq C (1 + \ln \varepsilon^{-1})^B$$

for all $\varepsilon \in (0, 1)$, where

$$C = 3b A \left(\ln(36A) (1 + B^3) \right)^B$$

and b is the absolute constant from Theorem 1 in the case $r = 1$.

Proof Observe that the inequality

$$n^{\text{all}}(\varepsilon, F) \leq A(1 + \ln \varepsilon^{-1})^B$$

implies that

$$c_n(F, L_2) \leq e \exp(-(n/A)^{1/B}).$$

We obtain from Theorem 2, and $1 + n^{1/2} \leq 2n^{1/2}$, that

$$a_n(F, L_2) \leq 2 e n^{1/2} \exp(-(n/A)^{1/B}).$$

Applying first Lemma 8 and then Lemma 9 from the [Appendix](#), we deduce that

$$\begin{aligned} \sum_{k \geq n} a_k(F, L_2) &\leq 2 e \sum_{k \geq n} k^{1/2} \exp(-(k/A)^{1/B}) \\ &\leq 6 A^{1/B} B \max(3B/2, 1) (n - 1)^{3/2-1/B} \exp(-((n - 1)/A)^{1/B}) \end{aligned}$$

for all $n \geq n_0(A, B) := A \max(3B/2, 1)^B + 1$. In particular, $(a_n(F, L_2)) \in \ell_1$. It follows from Theorem 1 that there exists an absolute constant $b \in \mathbb{N}$ such that

$$\begin{aligned} e_{bn}(F, L_2) &\leq n^{-1} \sum_{k \geq n} a_k(F, L_2) \\ &\leq 6 A^{1/B} B \max(3B/2, 1) (n - 1)^{1/2-1/B} \exp(-((n - 1)/A)^{1/B}) \end{aligned}$$

for all $n \geq n_0(A, B)$.

In the case $B \leq 2$ we have $(n - 1)^{1/2-1/B} \leq A^{1/2-1/B}$, and thus

$$e_{bn}(F, L_2) \leq 36 A^{1/2} \exp(-((n - 1)/A)^{1/B}).$$

If $B > 2$, then Lemma 10 with $u = 1/2 - 1/B$ yields for any $\delta > 0$ that

$$e_{bn}(F, L_2) \leq 9 A^{1/2} B^2 \delta^{1-B/2} \exp(((B/2 - 1)\delta - 1)((n - 1)/A)^{1/B})$$

and taking $\delta = 2/B$ yields

$$e_{bn}(F_d, L_2) \leq 36 A^{1/2} (B/2)^{B/2+1} \exp\left(-\frac{2}{B} \left(\frac{n - 1}{A}\right)^{1/B}\right).$$

If we put $B_0 := \max\{B/2, 1\}$, we have for all $B > 0$ and $n \geq n_0(A, B)$ the bound

$$e_{bn}(F_d, L_2) \leq 36 A^{1/2} B_0^{B_0+1} \exp\left(-\frac{1}{B_0} \left(\frac{n - 1}{A}\right)^{1/B}\right)$$

which is smaller than ε if

$$n \geq A B_0^B \left(\ln\left(36 A^{1/2} B_0^{B_0+1} \varepsilon^{-1}\right)\right)^B + 1.$$

Thus

$$\begin{aligned} n^{\text{std}}(\varepsilon, F_d) &\leq b \max \left\{ A B_0^B \left(\ln\left(36 A^{1/2} B_0^{B_0+1} \varepsilon^{-1}\right)\right)^B + 2, n_0(A, B) \right\} \\ &\leq 3b \max \left\{ A B_0^B R^B \left(1 + \ln(\varepsilon^{-1})\right)^B, A (3B_0)^B \right\} \\ &\leq 3b A B_0^B R^B \left(1 + \ln(\varepsilon^{-1})\right)^B \end{aligned}$$

with

$$R := \ln(36) + \frac{\ln A}{2} + (B_0 + 1) \ln B_0 \leq \ln(36A) B_0^2$$

which gives the desired estimate. □

Corollary 5 Assume that $F_d \in \mathcal{A}$ for every $d \in \mathbb{N}$ and

$$n^{all}(\varepsilon, F_d) \leq c d^q (1 + \ln \varepsilon^{-1})^p$$

for some $p, c > 0, q \geq 0$, and all $\varepsilon \in (0, 1)$. Then

$$n^{std}(\varepsilon, F_d) \leq n^{std\text{-}lin}(\varepsilon, F_d) \leq C d^q (1 + \ln d)^p (1 + \ln \varepsilon^{-1})^p$$

for all $\varepsilon \in (0, 1)$ and $d \in \mathbb{N}$, and some $C > 0$ that depends only on c, p and q .

In particular, if $\widetilde{\text{APP}}$ is exponentially (strongly) polynomially tractable for the class Λ^{all} then it is exponentially (strongly) polynomially tractable for Λ^{std} .

Proof We use Theorem 4 with $A = c d^q + 1$ and $B = p$. □

We now turn to the assumption that $\widetilde{\text{APP}}$ is exponentially quasi-polynomially tractable for the class Λ^{all} . This is the only case where we do not know if it implies the same property for Λ^{std} .

For convenience, let us write $\ln_+(x) := 1 + \ln(x)$.

Corollary 6 Assume that $F_d \in \mathcal{A}$ for every $d \in \mathbb{N}$ and

$$n^{all}(\varepsilon, F_d) \leq c \exp\left(t \cdot \ln_+ d \cdot \ln_+ \ln_+ \varepsilon^{-1}\right)$$

for some $c, t > 0$ and all $\varepsilon \in (0, 1)$. Then

$$n^{std}(\varepsilon, F_d) \leq c \widetilde{\exp}\left(t \cdot \ln_+ d \cdot \left(\ln_+ \ln_+ \varepsilon^{-1} + 4 \ln(t \ln_+ d) + C\right)\right)$$

for all $\varepsilon \in (0, 1)$ and $d > (e + \frac{1}{c})^{1/t} e^{-1}$, and some $C > 0$ that depends only on c .

In particular, if $\widetilde{\text{APP}}$ is exponentially quasi-polynomially tractable for the class Λ^{all} , then it is exponentially uniformly weakly tractable for the class Λ^{std} .

Proof Note that

$$c \exp\left(t(1 + \ln d)(1 + \ln(1 + \ln \varepsilon^{-1}))\right) = c e^t d^t \left(1 + \ln \varepsilon^{-1}\right)^{t(1 + \ln d)}.$$

Hence, we can apply Theorem 4 with $A = c e^t d^t$ and $B = t \ln_+ d$, i.e., $A = c e^B$. Note that $A, B \geq 1$ for $d > (e + \frac{1}{c})^{1/t} e^{-1}$. We obtain that there exists an absolute constant $b > 0$ such that

$$n^{std}(\varepsilon, F) \leq C (1 + \ln \varepsilon^{-1})^B$$

for all $\varepsilon \in (0, 1)$, with

$$\begin{aligned} C &= 3b A \left(\ln(36A)(1 + B^3)\right)^B \leq 3b A (2B)^{3B} (\ln(36c) + B)^B \\ &\leq c \exp\left(B\left(c' + 4 \ln(B)\right)\right) \end{aligned}$$

where $c' > 0$ only depends on c . This proves the bound.

Now, since $\ln n^{\text{std}}(\varepsilon, F_d)$ depends only logarithmically on d and double-logarithmically on ε^{-1} , we obtain

$$\lim_{d+\varepsilon^{-1} \rightarrow \infty} \frac{\ln n^{\text{std}}(\varepsilon, F_d)}{d^\alpha + (1 + \ln \varepsilon^{-1})^\beta} = 0$$

for all $\alpha, \beta > 0$, i.e., $\widetilde{\text{APP}}$ is exponentially uniformly weakly tractable for the class Λ^{std} . \square

We finally discuss EXP-UWT and EXP-WT.

Theorem 7 *Assume that $F_d \in \mathcal{A}$ for every $d \in \mathbb{N}$. If the problem $\widetilde{\text{APP}}$ is exponentially (uniformly) weakly tractable for the class Λ^{all} , then it is exponentially (uniformly) weakly tractable for the class Λ^{std} .*

Proof Assume that there are $0 < \alpha, \beta \leq 1$ such that

$$\lim_{d+\varepsilon^{-1} \rightarrow \infty} \frac{\ln n^{\text{all}}(\varepsilon, F_d)}{d^\alpha + (1 + \ln \varepsilon^{-1})^\beta} = 0.$$

It is enough to show that

$$\lim_{d+\varepsilon^{-1} \rightarrow \infty} \frac{\ln n^{\text{std}}(\varepsilon, F_d)}{d^\alpha + (1 + \ln \varepsilon^{-1})^\beta} = 0. \tag{3}$$

By assumption, for every $0 < h \leq 1/16$, there is some $v_0 \in \mathbb{N}$ such that

$$0 \leq \frac{\ln n^{\text{all}}(\varepsilon, F_d)}{d^\alpha + (1 + \ln \varepsilon^{-1})^\beta} \leq h$$

for all $\varepsilon \in (0, 1)$ and $d \in \mathbb{N}$ with $d^\alpha + (1 + \ln \varepsilon^{-1})^\beta \geq v_0$. It follows that

$$c_n(F_d, L_2) \leq e \exp \left(- \left(\frac{\ln n}{h} - d^\alpha \right)^{1/\beta} \right)$$

for all $n \geq \exp(hv_0)$. From Theorem 2, we get

$$a_n(F_d, L_2) \leq 2e \exp \left(- \left(\left(\frac{\ln n}{h} - d^\alpha \right)^{1/\beta} - \frac{1}{2} \ln n \right) \right).$$

For all $n \geq \exp(2hd^\alpha)$ and $h \leq 1/16$, we have

$$\left(\frac{\ln n}{h} - d^\alpha \right)^{1/\beta} - \frac{1}{2} \ln n \geq \frac{1}{2} \left(\frac{\ln n}{h} - d^\alpha \right)^{1/\beta} + \frac{1}{8} \frac{\ln n}{h}$$

and hence we have for all $n \geq \max\{\exp(hv_0), \exp(2hd^\alpha)\}$ that

$$a_n(F_d, L_2) \leq 2e \exp \left(- \frac{1}{2} \left(\frac{\ln n}{h} - d^\alpha \right)^{1/\beta} \right) \cdot n^{-1/(8h)}.$$

It follows from Theorem 1 that for some absolute constant $b \in \mathbb{N}$ and all $n \geq \max\{\exp(h\nu_0), \exp(2hd^\alpha)\}$, we have

$$\begin{aligned} e_{bn}(F_d, L_2) &\leq \frac{1}{n} \sum_{k \geq n} a_k(F_d, L_2) \leq \sum_{k \geq n} a_k(F_d, L_2) \\ &\leq 2e \exp\left(-\frac{1}{2} \left(\frac{\ln n}{h} - d^\alpha\right)^{1/\beta}\right) \sum_{k \geq n} k^{-1/(8h)} \\ &\leq 2e \exp\left(-\frac{1}{2} \left(\frac{\ln n}{h} - d^\alpha\right)^{1/\beta}\right), \end{aligned}$$

where we again used that $h \leq 1/16$. It follows that

$$n^{\text{std}}(\varepsilon, F_d) \leq D \exp\left(4h \left((1 + \ln \varepsilon^{-1})^\beta + d^\alpha\right)\right)$$

for some absolute constant $D > 0$ and all $\varepsilon \in (0, 1)$ and $d \in \mathbb{N}$ such that $d + \varepsilon^{-1}$ is sufficiently large. This implies

$$0 \leq \lim_{d+\varepsilon^{-1} \rightarrow \infty} \frac{\ln n^{\text{std}}(\varepsilon, F_d)}{d^\alpha + (1 + \ln \varepsilon^{-1})^\beta} \leq 4h.$$

Since $h \in (0, 1/16)$ can be chosen arbitrarily close to 0, we obtain (3).

This allows us to conclude our statement. Indeed, for uniform weak tractability we take arbitrary α and β from $(0, 1)$, and for weak tractability we take $\alpha = \beta = 1$. \square

Appendix Technical Lemmas

The following lemmas are used in the proofs of our results.

Lemma 8 *Let A and B be arbitrary positive real numbers. For $n \geq A(B/2)^B$ we have the following inequality*

$$\sum_{k \geq n+1} k^{1/2} \exp(-(k/A)^{1/B}) \leq \int_n^\infty t^{1/2} \exp(-(t/A)^{1/B}) dt.$$

Proof It is enough to show that the function $f : (0, \infty) \rightarrow \mathbb{R}$ given by $f(t) = t^{1/2} \exp(-(t/A)^{1/B})$ is decreasing on $(A(B/2)^B, \infty)$. Indeed, for $t > A(B/2)^B$ we have

$$\begin{aligned} f'(t) &= \left(\exp\left(\frac{1}{2} \ln(t) - (t/A)^{1/B}\right)\right)' = \\ &= \exp\left(\frac{1}{2} \ln(t) - (t/A)^{1/B}\right) \left(\frac{1}{2t} - (1/AB)(t/A)^{1/B-1}\right) < 0. \end{aligned}$$

\square

Lemma 9 *Let A and B be arbitrary positive real numbers. For every $n \geq A \max(3B/2, 1)^B$ we have the following inequality*

$$\int_n^\infty t^{1/2} \exp(-(t/A)^{1/B}) dt \leq A^{1/B} B \max(3B/2, 1) n^{3/2-1/B} \exp(-(n/A)^{1/B}).$$

Proof Using integration by substitution, with $u = (t/A)^{1/B}$, we obtain that

$$\int_n^\infty t^{1/2} \exp(-(t/A)^{1/B}) dt = A^{3/2} B \Gamma(3B/2, (n/A)^{1/B})$$

where, for $a \in \mathbb{R}$ and $x > 0$, $\Gamma(a, x) = \int_x^\infty v^{a-1} \exp(-v) dv$ is the incomplete gamma function.

It is known (see, e.g., Satz 4.4.3 in [13]) that for $a \geq 1$ and $x > a$ we have

$$\Gamma(a, x) \leq a x^{a-1} \exp(-x).$$

If, on the other hand, $0 < a < 1$ and $x > 1$ then since $v^{a-1} \leq x^{a-1}$ for $v \geq x$ we have

$$\Gamma(a, x) = \int_x^\infty v^{a-1} \exp(-v) dv \leq x^{a-1} \int_x^\infty \exp(-v) dv = x^{a-1} \exp(-x).$$

Therefore, for every $a > 0$ and $x > \max(a, 1)$, the following bound holds

$$\Gamma(a, x) \leq \max(a, 1) x^{a-1} \exp(-x).$$

Thus for $n > A \max(3B/2, 1)^B$, and taking $a = 3B/2$ and $x = (n/A)^{1/B}$, we have

$$\int_n^\infty t^{1/2} \exp(-(t/A)^{1/B}) dt \leq A^{1/B} B \max(3B/2, 1) n^{3/2-1/B} \exp(-(n/A)^{1/B}).$$

□

Lemma 10 *For every $A, B, n, \delta, u > 0$ we have the following inequality*

$$n^u \exp\left(- (n/A)^{1/B}\right) \leq A^u \delta^{-uB} \exp\left((uB\delta - 1)(n/A)^{1/B}\right).$$

Proof Let $x = \delta(n/A)^{1/B}$. Then $n^u = A^u \delta^{-uB} x^{uB}$. Using the fact that $\ln(x) \leq x$ for all $x > 0$ we obtain that

$$\begin{aligned} \ln(n^u) &= \ln(A^u \delta^{-uB}) + uB \ln(x) \leq \\ &\leq \ln(A^u \delta^{-uB}) + uBx = \ln(A^u \delta^{-uB}) + uB\delta(n/A)^{1/B} \end{aligned}$$

Hence, taking exponentials of both sides we derive that

$$n^u \leq A^u \delta^{-uB} \exp(uB\delta(n/A)^{1/B})$$

and thus

$$n^u \exp(- (n/A)^{1/B}) \leq A^u \delta^{-uB} \exp((uB\delta - 1)(n/A)^{1/B})$$

as claimed. □

Funding David Krieg is supported by the Austrian Science Fund (FWF) Project F5506, which is part of the Special Research Program ‘‘Quasi-Monte Carlo Methods: Theory and Applications’’.

Declarations

Conflict of interest The authors declare no competing interests.

Open Access This article is licensed under a Creative Commons Attribution 4.0 International License, which permits use, sharing, adaptation, distribution and reproduction in any medium or format, as long as you give appropriate credit to the original author(s) and the source, provide a link to the Creative Commons licence, and indicate if changes were made. The images or other third party material in this article are included in the article's Creative Commons licence, unless indicated otherwise in a credit line to the material. If material is not included in the article's Creative Commons licence and your intended use is not permitted by statutory regulation or exceeds the permitted use, you will need to obtain permission directly from the copyright holder. To view a copy of this licence, visit <http://creativecommons.org/licenses/by/4.0/>.

References

1. Adcock, B., Hansen, A.C., Shadrin, A.: A stability barrier for reconstructions from Fourier samples. *SIAM J. Numer. Anal.* **52**(1), 125–139 (2014)
2. Adcock, B., Huybrechs, D., Martin-Vaquero, J.: On the numerical stability of Fourier extensions. *Found. Comput. Math.* **14**, 635–687 (2014)
3. Adcock, B., Platte, R.B., Shadrin, A.: Optimal sampling rates for approximating analytic functions from pointwise samples. *IMA J. Numer. Anal.* **39**(3), 1360–1390 (2018)
4. Binev, P., Cohen, A., Dahmen, W., DeVore, R., Petrova, G., Wojtaszczyk, P.: Convergence rates for greedy algorithms in reduced basis methods. *SIAM J. Math. Anal.* **43**(3), 1457–1472 (2011)
5. Buffa, A., Maday, Y., Patera, A.T., Prud'homme C., Turinici, G.: A priori convergence of the greedy algorithm for the parametrized reduced basis method. *ESAIM: Mathematical modelling and numerical analysis – modélisation mathématique et analyse numérique* **46**(3), 595–603 (2012)
6. Creutzig, J., Wojtaszczyk, P.: Linear vs. nonlinear algorithms for linear problems. *J. Complex.* **20**, 807–820 (2004)
7. Deimer, J.J.A., Sergio, A.T.: Estimates for n-widths of sets of smooth functions on complex spheres. *J. Complex.* **64**, 101537 (2021)
8. Dick, J., Kritzer, P., Pillichshammer, F., Woźniakowski, H.: Approximation of analytic functions in Korobov spaces. *J. Complex.* **30**(2), 2–28 (2014)
9. Dolbeault, M., Krieg, D., Ullrich, M.: A sharp upper bound for sampling numbers in l_2 . *Appl. Comput. Harmon. Anal.* **63**, 113–134 (2023)
10. Dũng, D., Temlyakov, V.N., Ullrich, T.: *Hyperbolic Cross Approximation. Advanced Courses in Mathematics - CRM Barcelona.* Springer, Heidelberg (2018)
11. Ebert, A., Pillichshammer, F.: Tractability of approximation in the weighted Korobov space in the worst-case setting – a complete picture. *J. Complex.* **67**, 101571 (2021)
12. Fasshauer, G.E., Hickernell, F.J., Woźniakowski, H.: On dimension-independent rates of convergence for function approximation with Gaussian kernels. *SIAM J. Numer. Anal.* **50**(1), 247–271 (2012)
13. Gabcke, W.: *Neue Herleitung Und Explizite Restabschätzung Der Riemann-Siegel-Formel* PhD thesis. University of Göttingen (1979)
14. Griebel, M., Oettershagen, J.: On tensor product approximation of analytic functions. *J. Approx. Theory* **207**, 348–379 (2016)
15. Haasdonk, B.: Convergence rates of the pod–greedy method. *ESAIM: Mathematical Modelling and Numerical Analysis - Modélisation Mathématique et Analyse Numérique* **47**(3), 859–873 (2013)
16. Hackbusch, W., Khoromskij, B.N.: Tensor-product approximation to operators and functions in high dimensions. *J. Complex.* **23**(4), 697–714 (2007). *Festschrift for the 60th Birthday of Henryk Woźniakowski*
17. Hangelbroek, T., Ron, A.: Nonlinear approximation using Gaussian kernels. *J. Funct. Anal.* **259**(1), 203–219 (2010)
18. Hinrichs, A., Krieg, D., Novak, E., Prochno, J., Ullrich, M.: Random sections of ellipsoids and the power of random information. *Trans. Am. Math. Soc.* **374**(12), 8691–8713 (2021)
19. Hinrichs, A., Krieg, D., Novak, E., Vyřbřal, J.: Lower bounds for the error of quadrature formulas for Hilbert spaces. *J. Complex.* **65**, 101544 (2021)

20. Hinrichs, A., Novak, E., Ullrich, M.: On weak tractability of the clenshaw–Curtis Smolyak algorithm. *J. Approx. Theory* **183**, 31–44 (2014)
21. Hinrichs, A., Novak, E., Ullrich, M., Woźniakowski, H.: The curse of dimensionality for numerical integration of smooth functions. *Math. Comput.* **83**(290), 2853–2863 (2014)
22. Hinrichs, A., Novak, E., Ullrich, M., Woźniakowski, H.: Product rules are optimal for numerical integration in classical smoothness spaces. *J. Complex.* **38**, 39–49 (2017)
23. Hinrichs, A., Prochno, J., Ullrich, M.: The curse of dimensionality for numerical integration on general domains. *J. Complex.* **50**, 25–42 (2019)
24. Irrgeher, C., Kritzer, P., Pillichshammer, F., Woźniakowski, H.: Approximation in Hermite spaces of smooth functions. *J. Approx. Theory* **207**, 98–126 (2016)
25. Irrgeher, C., Kritzer, P., Pillichshammer, F., Woźniakowski, H.: Tractability of multivariate approximation defined over Hilbert spaces with exponential weights. *J. Approx. Theory* **207**, 301–338 (2016)
26. Karvonen, T., Särkkä, S.: Worst-case optimal approximation with increasingly flat Gaussian kernels. *Adv. Comput. Math.* **46**, 1–17 (2020)
27. Krieg, D.: Uniform recovery of high-dimensional c^r -functions. *J. Complex.* **50**, 116–126 (2019)
28. Krieg, D., Ullrich, M.: Function values are enough for l_2 -approximation. *Found. Comput. Math.* **21**(4), 1141–1151 (2021)
29. Krieg, D., Ullrich, M.: Function values are enough for l_2 -approximation: Part II. *J. Complex.* **66**, 101569 (2021)
30. Kritzer, P., Woźniakowski, H.: Simple characterizations of exponential tractability for linear multivariate problems. *J. Complex.* **51**, 110–128 (2019)
31. Mathé, P.: S-numbers in information-based complexity. *J. Complex.* **6**(1), 41–66 (1990)
32. Nagel, N., Schäfer, M., Ullrich, T.: A new upper bound for sampling numbers. [arXiv:2010.00327](https://arxiv.org/abs/2010.00327) (2020)
33. Novak, E., Woźniakowski, H.: Tractability of Multivariate Problems. Vol 1: Linear information, volume 6 of EMS tracts in mathematics. European Mathematical Society (EMS), Zürich (2008)
34. Novak, E., Woźniakowski, H.: Approximation of infinitely differentiable multivariate functions is intractable. *J. Complex.* **25**, 398–404 (2009)
35. Novak, E., Woźniakowski, H.: Tractability of multivariate problems volume II : Standard information for functionals, volume 12 of EMS tracts in mathematics. European Mathematical Society, (EMS), Zürich (2010)
36. Novak, E., Woźniakowski, H.: Tractability of multivariate problems volume III : Standard information for operators, volume 18 of EMS tracts in mathematics. European Mathematical Society, (EMS), Zürich (2012)
37. Novak, E., Woźniakowski, H.: Tractability of multivariate problems for standard and linear information in the worst case setting: part i. *J. Approx. Theory* **207**, 177–192 (2016)
38. Pietsch, A.: Operator ideals. Deutscher Verlag, Wiss. Berlin (1978)
39. Platte, R.B., Trefethen, L.N., Kuijlaars, A.B.J.: Impossibility of fast stable approximation of analytic functions from equispaced samples. *SIAM Rev.* **53**(2), 308–318 (2011)
40. Vybíral, J.: A variant of Schur’s product theorem and its applications. *Adv. Math.* **368**, 107140 (2020)
41. Sloan, I., Woźniakowski, H.: Multivariate approximation for analytic functions with Gaussian kernels. *J. Complex.* **45**, 1–21 (2018)
42. Temlyakov, V.N.: Approximation of periodic functions. Computational Mathematics and Analysis Series, Nova Science Publishers, Inc, Commack, NY (1993)
43. Temlyakov, V.N.: Multivariate Approximation, volume 32 of Cambridge Monographs on Applied and Computational Mathematics Cambridge University Press (2018)
44. Ullrich, M.: On the worst-case error of least squares algorithms for l_2 -approximation with high probability. *Journal of Complexity*, p. 60 (2020)
45. Wendland, H.: Scattered data approximation, vol. 17 of Cambridge monographs on applied and computational mathematics. Cambridge University Press, Cambridge (2005)
46. Xu, G.: On weak tractability of the Smolyak algorithm for approximation problems. *J. Approx. Theory* **192**, 347–361 (2015)
47. Zhang, J.: A note on EC-tractability of multivariate approximation in weighted Korobov spaces for the standard information class. *Journal of Complexity*, p. 67 (2021)