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A Proof of a Frankl–Kupavskii Conjecture on Intersecting Families

Agnijo Banerjee¹

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Abstract

A family $\mathcal{F} \subset \mathcal{P}(n)$ is *r*-wise *k*-intersecting if $|A_1 \cap \cdots \cap A_r| \ge k$ for any $A_1, \ldots, A_r \in \mathcal{F}$. It is easily seen that if \mathcal{F} is *r*-wise *k*-intersecting for $r \ge 2, k \ge 1$ then $|\mathcal{F}| \le 2^{n-1}$. The problem of determining the maximum size of a family \mathcal{F} that is both r_1 -wise k_1 -intersecting and r_2 -wise k_2 -intersecting was raised in 2019 by Frankl and Kupavskii (Combinatorica 39:1255–1266, 2019). They proved the surprising result that, for $(r_1, k_1) = (3, 1)$ and $(r_2, k_2) = (2, 32)$ then this maximum is at most 2^{n-2} , and conjecture d the same holds if k_2 is replaced by 3. In this paper we shall not only prove this conjecture but we shall also determine the exact maximum for $(r_1, k_1) = (3, 1)$ and $(r_2, k_2) = (2, 3)$ for all n.

Keywords Intersecting families \cdot Intersection properties \cdot Frankl-Kupavskii conjecture

1 Introduction

We say that a family $\mathcal{F} \subset \mathcal{P}(n)$ is *r*-wise *k*-intersecting if any *r* sets in \mathcal{F} have common intersection of size at least *k*. (If *r* is omitted it is assumed to be 2, and if *k* is omitted it is assumed to be 1.) We define the collections of families $\mathcal{A}_k(n)$ and $\mathcal{B}_k(n)$ to consist of all 3-wise *k*-intersecting families in $\mathcal{P}(n)$ and all *k*-intersecting families in $\mathcal{P}(n)$ respectively. We also define, for a family $\mathcal{F} \subset \mathcal{P}(n)$, the function $w(\mathcal{F}) = \frac{|\mathcal{F}|}{2n}$.

Frankl and Kupavskii [1] started the investigation of the function $W_{k_1,k_2}(n) = \max\{w(\mathcal{F}) : \mathcal{F} \in \mathcal{A}_{k_1}(n) \cap \mathcal{B}_{k_2}(n)\}$. This is the maximum proportion of $\mathcal{P}(n)$ that can be occupied by a family that is both 3-wise k_1 -intersecting and (2-wise) k_2 -intersecting. Equivalently, $W_{k_1,k_2}(n) = \max\{w(\mathcal{F} \cap \mathcal{G}) : \mathcal{F} \in \mathcal{A}_{k_1}(n), \mathcal{G} \in \mathcal{B}_{k_2}(n)\}$. We shall concentrate on the most important case, $k_1 = 1, k_2 = 3$. Frankl and Kupavskii conjectured the following:

Agnijo Banerjee ab2558@cam.ac.uk

¹ Department of Pure Mathematics and Mathematical Statistics (DPMMS), University of Cambridge, Wilberforce Road, Cambridge CB3 0WA, UK

Conjecture 1 If $\mathcal{F} \subset \mathcal{P}(n)$ is both 3-wise intersecting and 3-intersecting, then $|\mathcal{F}| \leq 2^{n-2}$. Equivalently, $W_{1,3}(n) \leq \frac{1}{4}$ for all n.

If \mathcal{F} is *r*-wise *k*-intersecting then so is the up-set generated by \mathcal{F} . Thus, we may assume throughout this paper that \mathcal{F} is an up-set. Also, applying left-compressions to \mathcal{F} preserves the property of being *r*-wise *k*-intersecting so we may also assume that, except where otherwise stated, \mathcal{F} is left-compressed.

Also, for a family $\mathcal{F} \subset \mathcal{P}(n)$, and for $k \geq 1$, we denote by $\mathcal{F} \times \{0, 1\}^k$ the family $\{A \subset \mathcal{P}(n+k) : A \cap [n] \in \mathcal{F}\}$. (This arises from considering subsets of [n] as binary sequences of length n). We observe that $w(\mathcal{F} \times \{0, 1\}^k) = w(\mathcal{F})$ and thus that $W_{k_1,k_2}(n)$ is non-decreasing for all k_1 and k_2 .

We note that $\max\{w(\mathcal{F}) : \mathcal{F} \in \mathcal{A}_1(n)\} = \frac{1}{2}$ (achieved by, for instance, letting \mathcal{F} consist of all subsets of [n] containing 1), and for every fixed $k_2 \geq 3$ we have $\max\{w(\mathcal{F}) : \mathcal{F} \in \mathcal{B}_{k_2}(n)\}$ tends to $\frac{1}{2}$ as $n \to \infty$. The Harris–Kleitman inequality thus tells us that

$$\liminf_{n\to\infty} W_{1,k_2}(n) \ge \frac{1}{4}.$$

In view of this trivial inequality, the result of Frankl and Kupavskii that $W_{1,32}(n) \le \frac{1}{4}$ for all *n* is surprising. In this paper we shall prove considerably more, namely that $W_{1,3}(n) < \frac{1}{4}$ for all *n*. Moreover, we shall determine the exact value of $W_{1,3}(n)$ for all *n*, and find the unique left-compressed families \mathcal{F} with $w(\mathcal{F}) = W_{1,3}(n)$. We write W(n) for $W_{1,3}(n)$.

For small values of *n*, the families of maximum size can be found by simply taking the maximum 3-intersecting families (as in [3]), since these will also be 3-wise intersecting. For n = 3, this is {123} and $W(3) = \frac{1}{8}$. Likewise, for n = 4, we get {123, 1234} and $W(4) = \frac{1}{8}$. For n = 5, we can take all sets of size at least 4 and $W(5) = \frac{3}{16}$. Likewise, for n = 6, we can take all sets whose intersection with [5] has size at least 4 and $W(6) = \frac{3}{16}$. Our main result will thus consider $n \ge 7$.

For odd $n \ge 7$, we define the following family:

$$\mathcal{F}_n = \{A \subset [n] : 1 \in A, |A| \ge \frac{n+3}{2}\} \cup \{A \subset [n] : 1 \notin A, |A| \ge n-2\}.$$

We have that \mathcal{F}_n is 3-intersecting and 3-wise intersecting. We also have that

$$w(\mathcal{F}_n) = \frac{1}{4} + 2^{-n} \left(-\frac{1}{2} \binom{n-1}{\frac{n-1}{2}} + n \right).$$

We note that $\lim_{n\to\infty} w(\mathcal{F}_n) = \frac{1}{4}$, and $w(\mathcal{F}_n) \le \frac{1}{4}$ for all $n \ge 7$. Also, for $n \ge 11$, we have $w(\mathcal{F}_{n+2}) > w(\mathcal{F}_n)$. However, we have $w(\mathcal{F}_7) > w(\mathcal{F}_9) > w(\mathcal{F}_{11})$, and indeed $w(\mathcal{F}_7) > w(\mathcal{F}_n)$ for odd $n, 9 \le n \le 71$. We shall prove the following theorem.

Theorem 1 For $n \ge 7$, the following hold.

If $7 \le n \le 72$, $W(n) = w(\mathcal{F}_7)$, and the unique left-compressed family \mathcal{F} attaining $w(\mathcal{F}) = W(n)$ is $\mathcal{F} = \mathcal{F}_7 \times \{0, 1\}^{n-7}$.

If $n \ge 73$ is odd, $W(n) = w(\mathcal{F}_n)$, and the unique left-compressed family \mathcal{F} attaining $w(\mathcal{F}) = W(n)$ is $\mathcal{F} = \mathcal{F}_n$

If $n \ge 74$ is even, $W(n) = w(\mathcal{F}_{n-1})$, and the unique left-compressed family \mathcal{F} attaining $w(\mathcal{F}) = W(n)$ is $\mathcal{F} = \mathcal{F}_{n-1} \times \{0, 1\}$.

Since $w(\mathcal{F}_n) \leq \frac{1}{4}$ for all $n \geq 7$, Theorem 1 completes the proof of Conjecture 1.

2 Determining the Values of W(n)

The proof of Theorem 1 is loosely inspired by the proof of the Ahlswede-Khachatrian Theorem [2]. We assume that \mathcal{F} is a left-compressed up-set. We also say that a family of subsets of [n] is trivial if every element contains 1 and almost-trivial if every element of size $\leq n - 3$ contains 1. (We use triviality in the sense of [1].) The main idea is to consider the generating set of minimal elements of the up-set \mathcal{F} , and attempt to transform this into a generating set for a family of subsets of [n - 1], without reducing the value of w. Lemmas 1 and 2 will show that this is always possible unless \mathcal{F} is almost-trivial, and Lemma 3 will establish the upper bound for $w(\mathcal{F})$ if it is almost-trivial.

For $A \neq B \subset [n]$ and |A| = |B| = r, we write $A \prec B$ if, for all $1 \le i \le |A|$, the *i*th element of *A* is at most the *i*th element of *B*. This is equivalent to saying that *A* can be obtained from *B* by left-compressions so any left-compressed family containing *B* must contain *A*.

For any up-set \mathcal{F} , we let its generating set $\mathcal{G} = \mathcal{G}(\mathcal{F})$ be the family of all minimal elements of \mathcal{F} , and we say that \mathcal{G} generates \mathcal{F} . We have that \mathcal{G} is an antichain and $\mathcal{F} = \{A \subset [n] : \exists B \in \mathcal{G}, B \subset A\}$. Also, for \mathcal{F} left-compressed, \mathcal{G} has the property that if $A \prec B$ and $B \in \mathcal{G}$ then $\exists C \in \mathcal{G}$ with $C \subset A$ (as $A \prec B$ implies $A \in \mathcal{F}$). If a generating set \mathcal{G} has this property, we say it is a left-compressed generating set (and in fact it then generates a left-compressed family).

For \mathcal{G} a left-compressed generating set for \mathcal{F} and $E \in \mathcal{G}$, we define

 $\mathcal{D}(E) = \{A \subset [n] : A \cap [\max(E)] = E\}.$

For distinct *E*, the families $\mathcal{D}(E)$ must be disjoint, since if $A \in \mathcal{D}(E) \cap \mathcal{D}(E')$ for $E \neq E' \in \mathcal{G}$, we may assume w.l.o.g. that $\max(E) \leq \max(E')$ in which case we get $E \subset E'$, contradicting the fact that \mathcal{G} is an antichain. Also, since every element of $\mathcal{D}(E)$ is a superset of *E*, we have that $\mathcal{D}(E) \subset \mathcal{F}$ for all $E \in \mathcal{G}$.

For $A \in \mathcal{F}$, let $E \in \mathcal{G}$ such that $E \subset A$, with |E| minimal. Let E' consist of the first |E| elements of A. Then $E' \prec E$ and $E' \subset A$. If $E' \notin \mathcal{G}$ then there a set $C \in \mathcal{G}$ with $C \subset E'$ and |C| < |E'| = |E|. Since $C \subset E'$, we have that $C \subset A$, contradicting minimality of |E|. Thus $E' \in \mathcal{G}$ and in fact $A \in \mathcal{D}(E')$.

We have that the families $\mathcal{D}(E)$ are all subfamilies of \mathcal{F} and are disjoint for distinct E, and that every set in \mathcal{F} is in $\mathcal{D}(E)$ for some $E \in \mathcal{G}$. Thus \mathcal{F} is the disjoint union of the $\mathcal{D}(E)$ for $E \in \mathcal{G}$.

We thus have

$$|\mathcal{F}| = \sum_{E \in \mathcal{G}} |\mathcal{D}(E)|.$$

For a given $E \in \mathcal{G}$, we have $|\mathcal{D}(E)| = 2^{n - \max(E)}$ so we obtain

$$w(\mathcal{F}) = \sum_{E \in \mathcal{G}} 2^{-\max(E)}.$$

Now, we split \mathcal{G} into two parts: $\mathcal{G}(n) = \{A \in \mathcal{G} : n \in A\}$ and $\mathcal{G}(\overline{n}) = \mathcal{G} \setminus \mathcal{G}(n)$. If $\mathcal{G}(n)$ is empty, we may consider $\mathcal{G}(\overline{n})$ as a generating set for a family $\mathcal{F}' \subset \mathcal{P}(n-1)$, for which $\mathcal{F} = \mathcal{F}' \times \{0, 1\}$. In this case, $w(\mathcal{F}) = w(\mathcal{F}')$.

We thus aim to transform \mathcal{F} into a family in which $\mathcal{G}(n)$ is empty, and we may do so by two means. The first is to simply remove some element A of $\mathcal{G}(n)$, which decreases $w(\mathcal{F})$ by 2^{-n} . The second is to replace an element $A \in \mathcal{G}(n)$ with $A' = A \setminus \{n\}$, and possibly remove some other elements of \mathcal{G} that are supersets of A'. This does not remove any elements of \mathcal{F} but it does add $A \setminus \{n\}$ (which was not previously in \mathcal{F} as otherwise A would not be minimal) so it increases $w(\mathcal{F})$ by at least 2^{-n} . (In fact, this has the effect of adding $A \setminus \{n\}$ to \mathcal{F} but leaving it otherwise unchanged, so it increases $w(\mathcal{F})$ by precisely 2^{-n}). We refer to this operation as shortening A.

If every element of $\mathcal{G}(n)$ is either removed or shortened, the resulting family will be of the form $\mathcal{F}' \times \{0, 1\}$ for some family $\mathcal{F}' \subset \mathcal{P}(n-1)$, so we seek to remove or shorten every element of $\mathcal{G}(n)$. However, shortening may cause a violation of the intersection properties (as if *A* is in $\mathcal{G}(n)$, *B* and *C* are in \mathcal{G} , and $A' = A \setminus \{n\}$, it does not follow that $|A' \cap B| \ge 3$ or that $A' \cap B \cap C \ne \phi$) so we must avoid this situation when shortening.

Suppose that some $A \in \mathcal{G}(n)$ cannot be shortened. Then there are two cases:

Case 1: $\exists B \in \mathcal{G}$ such that $|A' \cap B| < 3$. Case 2: $\exists B, C \in \mathcal{G}$ such that $A' \cap B \cap C = \phi$.

We consider Case 1 first. We have $A' \cap B = (A \cap B) \setminus n$ so the only way this can occur is if $|A \cap B| = 3$ and $n \in A \cap B$. Suppose there was some $i \notin A \cup B$. Then we consider $A^* = A \setminus \{n\} \cup \{i\}$. Since $A^* \prec A$, we have $A^* \in \mathcal{F}$ but $|A^* \cap B| = 2$, contradicting that \mathcal{F} is 3-intersecting. Thus $A \cup B = [n]$. We will refer to a pair (A, B)of sets in $\mathcal{G}(n)$ with $A \cup B = [n]$ and $|A \cap B| = 3$ as a sharp pair. If i < j < n and $A \cap B = \{i, j, n\}$, we say (A, B) is an (i, j)-sharp pair.

We now consider Case 2. Again, we have $A' \cap B \cap C = (A \cap B \cap C) \setminus \{n\}$ so we require $A \cap B \cap C = \{n\}$. Suppose there was some *i* in at most one of *A*, *B*, and *C*. Then either $i \notin B$ or $i \notin C$, and w.l.o.g. we can assume the former. We consider $B^* = B \setminus \{n\} \cup \{i\}$, and again $B^* \prec B$ so $B^* \in \mathcal{F}$, and $A \cap B^* \cap C = \phi$, contradicting that \mathcal{F} is 3-wise intersecting. Thus every *i* appears in exactly two of *A*, *B*, and *C*, except for i = n which appears in all three. We will refer to a triple (A, B, C) of sets in $\mathcal{G}(n)$ with every element other than *n* appearing in exactly two of *A*, *B*, and *C* as a sharp triple.

We have thus shown that A can be shortened unless it is part of a sharp pair or a sharp triple, so for $w(\mathcal{F})$ maximum we may assume that every element of $\mathcal{G}(n)$ is in at least one sharp pair or triple. In fact, shortening and removing some elements will preserve the intersection properties provided that in every sharp pair or triple in which at least one element is shortened, we also have at least one element removed.

We have the following two lemmas, which together imply that either $w(\mathcal{F}) \leq W(n-1)$ or \mathcal{F} is almost-trivial. This is helpful as it is much easier to analyse the maximum size of an almost-trivial family.

Lemma 1 If there is no (i, n-1)-sharp pair in $\mathcal{G}(n)$ for some i < n-1 then $w(\mathcal{F}) \le W(n-1)$.

Proof Assume there is no (i, n-1)-sharp pair. We consider the sharp pairs and triples in $\mathcal{G}(n)$. Every sharp triple contains exactly two sets containing n-1 and one not containing n-1. In a sharp pair (A, B), we know that (A, B) is not (i, n-1)-sharp so n-1 is not in $A \cap B$. Thus exactly one of A and B contains n-1 and one does not.

We can now partition $\mathcal{G}(n)$ into $\mathcal{G}(n-1,n)$ and $\mathcal{G}(\overline{n-1},n)$, with $\mathcal{G}(n-1,n)$ consisting of those sets in $\mathcal{G}(n)$ containing n-1 and $\mathcal{G}(\overline{n-1},n)$ those not containing n-1. Then we may shorten all elements of whichever of $\mathcal{G}(n-1,n)$ and $\mathcal{G}(\overline{n-1},n)$ is larger (or either if they are of the same size) and remove all elements of the other, producing a generating set \mathcal{G}' for a family \mathcal{F}' . Since every sharp pair and triple contains both an element of $\mathcal{G}(n-1,n)$ and an element of $\mathcal{G}(\overline{n-1},n)$, this will preserve the intersection properties regardless of which is shortened and which is removed. Because shortening a set increases $w(\mathcal{F})$ by 2^{-n} and removing a set decreases $w(\mathcal{F})$ by 2^{-n} , we have

$$w(\mathcal{F}') = w(\mathcal{F}) + 2^{-n} \left| |\mathcal{G}(n-1,n)| - |\mathcal{G}(\overline{n-1},n)| \right|.$$

In particular, $w(\mathcal{F}') \ge w(\mathcal{F})$. Also, since every element of $\mathcal{G}(n)$ was either removed or shortened, no element of \mathcal{G}' contains n, so \mathcal{G}' also generates a family $\mathcal{F}'' \subset \mathcal{P}(n-1)$ with $\mathcal{F}' = \mathcal{F}'' \times \{0, 1\}$. Thus $w(\mathcal{F}') \le W(n-1)$ so $w(\mathcal{F}) \le W(n-1)$. \Box

Lemma 2 If there is an (i, n - 1)-sharp pair in $\mathcal{G}(n)$ for some i < n - 1 then \mathcal{F} is almost-trivial.

Proof We have that there is an (i, n - 1)-sharp pair, for some i < n - 1. Let this sharp pair be (A, B). Suppose there was $C \in \mathcal{F}$ with $|C| \le n - 3$ and $1 \notin C$. Then there is also some other $j \notin C$ with $1 < j \le n - 1$, so if $n \in C$ we can take a compression to obtain $C' = C_{jn}(C)$ which does not contain n (and is in \mathcal{F} since it is left-compressed), so we may assume w.l.o.g. $n \notin C$. There is still a (possibly different) $j \notin C$ with $1 < j \le n - 3$. We now take $A' = C_{j(n-1)}(C_{1i}(A))$ and $B' = C_{j(n-1)}(C_{1i}(B))$. Since \mathcal{F} is left-compressed, both A' and B' are in \mathcal{F} . However, $A' \cap B' = \{1, j, n\}$ so $A' \cap B' \cap C' = \phi$, contradicting that \mathcal{F} is 3-wise intersecting. Thus there is no such C, so \mathcal{F} is almost-trivial.

The next lemma concerns the maximum size of an almost-trivial family.

Lemma 3 Let \mathcal{F} be almost-trivial.

If n is odd, then $w(\mathcal{F}) \leq w(\mathcal{F}_n)$ with equality iff $\mathcal{F} = \mathcal{F}_n$. If n is even, then $w(\mathcal{F}) < w(\mathcal{F}_{n-1})$.

Proof First, the only elements of \mathcal{F} that could possibly fail to contain 1 are those of size at least n - 2, and there are at most n of them. Thus we may remove them to form a trivial family \mathcal{F}' with $w(\mathcal{F}') \ge w(\mathcal{F}) - n2^{-n}$. We may remove 1 from each set in \mathcal{F}' to form a family \mathcal{F}'' of subsets of [2, n], with $|\mathcal{F}'| = |\mathcal{F}''|$. Since \mathcal{F}' is 3-intersecting, \mathcal{F}'' is 2-intersecting.

As proven in [3], if *n* is odd, say n = 2l + 1, then $|\mathcal{F}''| \leq 2^{n-2} - \frac{1}{2} {n-1 \choose l}$ with equality iff $\mathcal{F}'' = [2, n]^{(\geq l+1)}$. In this case we get precisely that $w(\mathcal{F}) \leq w(\mathcal{F}_n)$, and the equality case we get is precisely \mathcal{F}_n .

Also, from [3], if *n* is even, say n = 2l + 2, then $|\mathcal{F}''| \leq 2^{n-2} - \binom{n-2}{l}$. In this case we get that $w(\mathcal{F}) \leq \frac{1}{4} + 2^{-n}(-\binom{n-2}{l} + n) = w(\mathcal{F}_{n-1}) - (n-2)2^{-n}$. The equality case we obtain is $\mathcal{F} = (\mathcal{F}_{n-1} \times \{0, 1\}) \setminus \{A \subset [n] : 1, n \notin A, |A| = n - 3\}$. Thus, as \mathcal{F} is simply $\mathcal{F}_{n-1} \times \{0, 1\}$ with some elements (namely, all sets of size n - 3 not containing 1 or *n*) excluded, we have $w(\mathcal{F}) < w(\mathcal{F}_{n-1})$.

We can now prove Theorem 1 by induction. For n = 7, we know that \mathcal{F} is 3-intersecting so, again from [3], the 3-intersecting family of subsets of [7] of maximum size is in fact \mathcal{F}_7 , and this is unique. Since \mathcal{F}_7 is also 3-wise intersecting, it is the unique maximum family so $W(7) = w(\mathcal{F}_7)$.

Assume Theorem 1 holds for n - 1. By Lemmas 1 and 2, either $w(\mathcal{F}) \leq W(n - 1)$ or \mathcal{F} is almost-trivial. If n is even, then by Lemma 3, if \mathcal{F} is almost-trivial then $w(\mathcal{F}) < w(\mathcal{F}_{n-1})$, so \mathcal{F} is not optimal. Thus the former case must hold so W(n) = W(n - 1).

If *n* is odd, then we must have either $w(\mathcal{F}) \leq W(n-1)$ or $w(\mathcal{F}) \leq w(\mathcal{F}_n)$. Thus we have $W(n) = \max(W(n-1), w(\mathcal{F}_n))$. Since we assume Theorem 1 holds for n-1, we have that W(n) = W(n-1) for $n \leq 71$, since $W(n-1) = w(\mathcal{F}_7) \geq w(\mathcal{F}_n)$. However, for $n \geq 73$, we instead have $W(n) = w(\mathcal{F}_n)$, and in this case the unique maximum left-compressed up-set is \mathcal{F}_n . Thus, by induction, we have that the value of W(n) is precisely as stated in Theorem 1 for all $n \geq 7$. This suffices to prove Conjecture 1.

3 Uniqueness of Left-Compressed Families in Theorem 1

It remains to demonstrate uniqueness of the left-compressed families \mathcal{F} satisfying $w(\mathcal{F}) = W(n)$ in the cases $n \leq 72$ or n even. In this case, we know that \mathcal{F} cannot be almost-trivial so must not have an (i, n - 1)-sharp pair. We can then partition $\mathcal{G}(n)$ into $\mathcal{G}(n - 1, n)$ and $\mathcal{G}(\overline{n-1}, n)$ as in the proof of Lemma 1, and may shorten one of these and remove the other. If $|\mathcal{G}(n - 1, n)| \neq |\mathcal{G}(\overline{n-1}, n)|$ then doing so will strictly increase w, which is impossible since \mathcal{F} is maximum. Thus $|\mathcal{G}(n - 1, n)| = |\mathcal{G}(\overline{n-1}, n)|$ and so we may shorten either one of $\mathcal{G}(n - 1, n)$ and $\mathcal{G}(\overline{n-1}, n)$ and remove the other without changing the value of w.

If we shorten $|\mathcal{G}(n-1, n)|$, we remove from \mathcal{F} those elements of \mathcal{G} containing both n and n-1, and add elements $A \setminus \{n\}$ for all $A \in \mathcal{G}$ containing n but not n-1. Suppose

that shortening produces the family \mathcal{F}' and we have $A \in \mathcal{F}'$, $B \notin \mathcal{F}'$ with $B \prec A$. Since \mathcal{F} is left-compressed, either *B* has been removed from \mathcal{F} or *A* has been added. If *B* was removed, since $B \prec A$, we must have that *A* contains *n* and *n* – 1 and thus was not added. If $A \in \mathcal{G}$ it would also have been removed, so $A \notin \mathcal{G}$. Then there is an $A' \subset A \in \mathcal{F}$ with |A'| < |A|, and a corresponding subset $B' \subset B$ consisting of the initial |A'| elements of *B*. Since $B \prec A$, we have $B' \prec A'$ so $B' \in \mathcal{F}$, contradicting $B \in \mathcal{G}$. Otherwise, if *A* was added, we must have $A' = A \cup \{n\} \in \mathcal{G}$ and $n - 1 \notin A$. Then we have $B' = B \cup \{n\} \prec A'$ so $B' \in \mathcal{F}$. If $B' \notin \mathcal{G}$ then there is a proper subset $B'' \subset B'$ in \mathcal{F} with |B''| = |A|, and then $B \prec B''$ contradicting $B \notin \mathcal{F}'$. Thus $B' \in \mathcal{G}$, in which case we must have $n - 1 \notin B'$ so B' is also shortened, once again contradicting $B \notin \mathcal{F}'$. Thus, there can be no such *A* and *B*, so \mathcal{F}' is left-compressed.

In this case, by the inductive hypothesis, we conclude that \mathcal{F}' is an extension of the unique maximum left-compressed family for n - 1.

In the case $n \leq 72$, the family \mathcal{F}' must be an extension of \mathcal{F}_7 , i.e. $\mathcal{F}' = \mathcal{F}_7 \times \{0, 1\}^{n-7}$. Thus for every $A \subset [7]$ with |A| = 5, exactly one of A and $A \cup \{n\}$ is in \mathcal{G} , and these are the only elements of $\mathcal{G}(n-1,n)$. If $n \geq 9$, for every $B \subset [7]$ with |B| = 3 we can choose $A_1, A_2 \subset [7]$ with $|A_1| = |A_2| = 5$ and $A_1 \cap A_2 = B$. If both correspond to sets in $\mathcal{G}(n-1,n)$ containing n, we can perform a left-compression to one of them to replace the n by n-1, to get two sets in \mathcal{F} whose intersection is B. Then, for any $C \in \mathcal{F}$, if $|C \cap [7]| \leq 4$, then we can choose $B \subset [7]$ with |B| = 3 and $B \cap C = \phi$, contradicting that \mathcal{F} is 3-wise intersecting. Thus, for all $C \in \mathcal{F}$, we have $|C \cap [7]| \geq 5$ so we have $\mathcal{F} \subset \mathcal{F}_7 \times \{0, 1\}^{n-7}$. Thus this is the unique left-compressed family \mathcal{F} with $w(\mathcal{F}) = W(n)$. For n = 8, again from [3] we have that this is the unique left-compressed and thus is again the unique left-compressed family \mathcal{F} with $w(\mathcal{F}) = W(n)$.

The final remaining case is $n \ge 74$ even. Again, as before, \mathcal{F}' must be an extension of the unique maximum left-compressed family for n-1. By the inductive hypothesis, \mathcal{F}' is an extension of \mathcal{F}_{n-1} . Thus, for $A \subset [n-1]$ with $|A| = \frac{n+2}{2}$ and $1 \in A$, either A or $A \cup \{n\}$ is in \mathcal{G} , and if $A \cup \{n\}$ is in \mathcal{G} it is in $\mathcal{G}(\overline{n-1}, n)$. In particular, if $n-1 \in A$ then $A \cup \{n\}$ cannot be in $\mathcal{G}(\overline{n-1}, n)$ so A must be in \mathcal{G} . Now suppose there was some $B \in \mathcal{G}(n-1,n)$ with $|B| \le \frac{n+2}{2}$. Then we can choose $A \subset [n-1]$ with $|A| = \frac{n+2}{2}$ and $1, n-1 \in A$ such that $|A \cap B| = 2$ (if $1 \in B$ we must have $A \cap B = \{1, n-1\}$, otherwise it contains n-1 and some other element). Thus every element B of $\mathcal{G}(n-1,n)$ has $|B| \ge \frac{n+4}{2}$. If $1 \in B$ then $B \setminus \{n\}$ contains 1 and has size at least $\frac{n+2}{2}$ so it is a superset of a set in \mathcal{G} , contradicting $B \in \mathcal{G}$. Thus, for all $B \in \mathcal{G}(n-1,n)$, we must also have that $1 \notin B$.

However, we can take two different subsets of [n-1] of size $\frac{n+2}{2}$ containing 1 and n-1, which are thus both in \mathcal{G} , and whose intersection is precisely $\{1, n-2, n-1\}$, so every set in \mathcal{F} must intersect this set. For $B \in \mathcal{G}(n-1,n)$, if $|B| \leq n-3$, we can perform left-compressions to obtain some $B' \in \mathcal{F}$ with $1, n-2, n-1 \notin B'$, thus contradicting that \mathcal{F} is 3-wise intersecting. Hence, for all $B \in \mathcal{G}(n-1, n)$, we have $|B| \geq n-2$. But we also have in \mathcal{F}' all sets that do not contain 1 and have size n-3, and all such B are supersets of some such set. However, the only set that could have arisen in this manner from shortening a set in $\mathcal{G}(n-1, n)$ is [2, n-2] (since sets in $\mathcal{G}(n-1, n)$ must exclude n-1) and so any B may only be a superset

of this set and no other set of size n - 3 not containing 1. This is impossible as all such *B* contain n - 1. Thus there are no such *B* so $\mathcal{G}(n - 1, n)$ is empty, and since $|\mathcal{G}(n - 1, n)| = |\mathcal{G}(\overline{n - 1}, n)|$ we have $\mathcal{G}(\overline{n - 1}, n)$ is also empty and so $\mathcal{F} = \mathcal{F}'$, completing the proof of Theorem 1.

4 Concluding Remarks and Open Problems

In [1], Frankl and Kupavskii define monotone properties Π_1 and Π_2 to be incompatible if

$$\liminf_{n \to \infty} \left(\max_{\mathcal{F}_1 \subset \mathcal{P}(n) \text{ has } \Pi_1} w(\mathcal{F}_1) \right) \left(\max_{\mathcal{F}_2 \subset \mathcal{P}(n) \text{ has } \Pi_2} w(\mathcal{F}_2) \right)$$
$$= \liminf_{n \to \infty} \left(\max_{\mathcal{F} \subset \mathcal{P}(n) \text{ has } \Pi_1, \Pi_2} w(\mathcal{F}) \right).$$

By the Harris–Kleitman inequality, the latter is always at least the former, so Π_1 and Π_2 are incompatible if equality holds in the limit as $n \to \infty$ in the Harris–Kleitman inequality.

Conjecture 1 is equivalent to the properties of being 3-wise intersecting and 3intersecting being incompatible. In [1], Frankl and Kupavskii further conjecture the following.

Conjecture 2 For every integer $s \ge 1$ there exists an integer $t_0(s)$ such that for all $t \ge t_0(s)$, the properties of being 3-wise s-intersecting and t-intersecting are incompatible.

The s = 1 case of this conjecture was proved by Frankl and Kupavskii [1], and Theorem 1 implies that $t_0(1) = 3$. However, this conjecture remains open for all other values of *s*.

We may attempt to analyse the behaviour of $W_{k_1,k_2}(n)$ for $k_2 > k_1 > 1$ to attempt to prove this conjecture for $k_1 = s > 1$. However, the methods in this paper are not sufficient to do so, as the proof of the analogue of Lemma 1 fails for $k_1 > 1$. In this case, the analogue of a sharp triple is a triple A, B, C of sets in $\mathcal{G}(n)$ in which $|A \cap B \cap C| = k_1$ and all elements of [n] are contained in at least two of A, B, and C. Unlike in the case $k_1 = 1$, if $k_1 > 1$ it is possible for all three of these to contain n - 1, in which case it is not possible to shorten $\mathcal{G}(n - 1, n)$ and remove $\mathcal{G}(\overline{n - 1}, n)$. Likewise, if $k_2 > k_1 + 2$, the proof of the analogue of Lemma 2 also fails.

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Data Availability This paper doesn't use any raw data.

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