Discrete Comput Geom 34:129–141 (2005) DOI: 10.1007/s00454-004-1153-z



## **Covering with Fat Convex Discs\***

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**Abstract.** According to a theorem of L. Fejes Tóth [4], if non-crossing congruent copies of a convex disc *K* cover a convex hexagon *H*, then the density of the discs relative to *H* is at least area  $K/f_K(6)$  where  $f_K(6)$  denotes the maximum area of a hexagon contained in *K*. We say that a convex disc is *r*-fat if it is contained in a unit circle *C* and contains a concentric circle *c* of radius *r*. Recently, Heppes [7] showed that the above inequality holds without the non-crossing assumption if *K* is a 0.8561-fat ellipse. We show that the non-crossing assumption can be omitted if *K* is an  $r_0$ -fat convex disc with  $r_0 = 0.933$  or an  $r_1$ -fat ellipse with  $r_1 = 0.741$ .

#### 1. Introduction

By a *convex disc* we mean a compact convex set with non-empty interior and by an *n*-gon we mean a polygon with at most *n* sides in the Euclidean plane. We denote the interior, boundary, convex hull, and area of a disc *K* by int *K*, bd *K*, conv *K*, and area *K*, respectively. Further, let  $f_K(n)$  denote the maximum area of an *n*-gon contained in *K*. To simplify our notation, we omit the subscript if *K* is a unit circle. Thus

$$f(n) = \frac{n}{2}\sin\frac{2\pi}{n}$$

stands for the area of a regular *n*-gon inscribed in a unit circle.

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We consider coverings of a convex hexagon H by congruent copies  $K_1, \ldots, K_N$  of a convex disc K. The density of these discs relative to H is defined as  $\sum_{i=1}^{N} \operatorname{area} K_i/\operatorname{area} H$ . We say that two convex discs *cross* if removing their intersection from them each disc becomes non-connected.

<sup>\*</sup> This work was done at the Mathematical Sciences Research Institute at Berkeley, CA, where the author was participating in a semester-long program on Discrete and Computational Geometry. The research was also supported by OTKA Grants T 030012, T 038397, and T 043520.

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According to a theorem of L. Fejes Tóth [4], (see also p. 167 in [5]) if N non-crossing congruent copies of a convex disc K cover a convex hexagon H, then the density of the discs relative to H is at least

area  $K/f_K(6)$ .

The bound is exact for N = 1. If K is centrally symmetric, then it is also asymptotically sharp, or in other words, it is sharp for the whole plane.

The assumption that the discs do not cross seems to be superfluous. However, as examples by Heppes and Wegner show, we cannot simply exclude the possibility that crossing pairs occur. Wegner [11] constructed for all N > 1 a convex hexagon H and a convex disc K such that N congruent copies of K cover H but they cannot be rearranged to obtain a crossing-free covering of H.

The use of crossing pairs in a covering is particularly wasteful if *K* is close to a circle. L. Fejes Tóth suggested that it might be possible to eliminate the crossing-free condition at least for "round" convex discs. A first step in this direction was recently achieved by Heppes [7]. He called ellipses with half-axis a = 1 and  $1 \ge b \ge b_0 = 0.8561$  fat, and showed that if a convex hexagon is covered by congruent fat ellipses, then the density of the ellipses relative to the hexagon is at least  $2\pi/\sqrt{27}$ . Quite against the convention, here and throughout the paper we use the term ellipse to denote a convex disc, rather than a curve. Extending the notion of fatness, we say that a convex disc *K* is *r*-fat if it is contained in a unit circle *C* and contains a concentric circle *c* of radius *r*. We call the set  $C \setminus c$  containing bd *K* an *annulus associated to K*. The common center of the circles *C* and *c* is referred to as the *center of K*. Our main result is the following

**Theorem 1.** Let  $r_0 = 0.933$ . If congruent copies of an  $r_0$ -fat convex disc K cover a convex hexagon H, then the density of the discs relative to H is at least

 $\frac{\operatorname{area} K}{f_K(6)}$ 

We also weaken a little the fatness condition in Heppes's result for ellipses:

**Theorem 2.** Let  $r_1 = 0.741$ . If congruent  $r_1$ -fat ellipses cover a convex hexagon H, then the density of the ellipses relative to H is at least

$$\frac{\pi}{f(6)} = \frac{2\pi}{\sqrt{27}}.$$

If a countable system of convex discs covers the whole plane, their density can be defined by a suitable limit (see, e.g., [6]). Recall that the *covering density*  $\vartheta(K)$  of a convex disc *K* is defined as the infimum of the densities of all coverings of the whole plane by congruent copies of *K*. Before Heppes's result, besides the trivial examples of space-fillers, the circle was the only convex body for which  $\vartheta(K)$  was known. Theorem 1 implies the following:

Corollary. We have

$$\vartheta(K) = \frac{\operatorname{area} K}{f_K(6)}$$

for all centrally symmetric  $r_0$ -fat convex discs.

Principally, this enables us to determine the covering density of all  $r_0$ -fat convex discs. Of course, we still face the problem of calculating  $f_K(6)$ , which is a difficult task in general, but it can be done in special cases. We give some examples in Section 5.

We prove Theorems 1 and 2 in Section 4. Their proof is prepared in the next section by stating some known results which we shall need in the proof and in Section 3 by proving a lemma claiming that crossing fat discs are close.

#### 2. Some Auxiliary Results

In this section we gather some known results needed in the proof of Theorems 1 and 2.

**Proposition 1** [7]. If  $E_1$  and  $E_2$  are ellipses with half-axes  $a_1 \ge b_1$  and  $a_2 \ge b_2$  such that

$$\min_{i=1,2} \frac{a_i^2}{b_i} \le \max_{i=1,2} \frac{b_i^2}{a_i},$$

then  $E_1$  and  $E_2$  do not cross.

The special case when one of the ellipses is a circle is stated on p. 480 in [7]. The proof of the general case is similar.

**Proposition 2** [7]. Let q be a point on the boundary of an ellipse of half-axes a and  $b, a \ge b$ , let  $l_1$  be a line touching the ellipse at q and let  $l_2$  be the line orthogonal to  $l_1$  and containing the center of the ellipse. Then the maximum of the distance from q to  $l_2$  is a - b.

This is Lemma 1 in [7] (with the notation slightly changed in order to fit to ours).

**Proposition 3** [7], [1]. Let  $K_1, \ldots, K_N$  be convex discs covering a convex hexagon H. Suppose that no pair of the discs  $K_1, \ldots, K_N$  cross and no proper subset of them covers H. Then it is possible to construct convex polygons  $D_1, \ldots, D_N$  with the number of sides  $n_1, \ldots, n_N$  such that

$$D_i \subset K_i \cap H$$
 for  $i = 1, \ldots, N$ ,

$$\bigcup_{i=1}^N D_i = H,$$

$$(\operatorname{int} D_i) \cap (\operatorname{int} D_j) = \emptyset$$
 for  $i, j = 1, \dots, N, i \neq j$ ,

and

$$\sum_{i=1}^{N} n_i \le 6N.$$

The construction of polygons with the above properties was described first in [4], however, the presentation there is very vague. A more thorough treatment can be found on p. 170 in [5]. A very detailed description of the construction is given in [1]. A more general statement is proved in [2].

**Proposition 4** [10]. We have, for all convex discs K and for all integers  $n \ge 3$ ,

$$\frac{\operatorname{area} K}{\pi} f(n) \le f_K(n)$$

For a proof see [10], [6, pp. 36–37], or [9, pp. 14–15].

**Proposition 5** [3]. The sequence  $f_K(n)$  is concave for all convex discs K:

$$f_K(n+1) - f_K(n) \le f_K(n) - f_K(n-1)$$
 for  $n \ge 4$ .

See [3], [5, p. 169], [6, pp. 34–35], or [9, pp. 11–13] for a proof.

#### 3. Crossing Fat Discs are Close

Intuitively it is clear that if two fat discs cross, then they are close in some sense. The following lemma describes this precisely for different types of crossing fat discs.

#### Lemma.

(i) If two r-fat discs cross, then the distance between their centers is at most

$$2\sqrt{1-r^2}$$
.

(ii) If two ellipses with half-axes  $a_1$ ,  $b_1$  and  $a_2$ ,  $b_2$  ( $a_1 \ge b_1$ ,  $a_2 \ge b_2$ ) cross, then the distance between their centers is at most

$$\sqrt{(a_1+a_2-b_1-b_2)^2+\left(\max_{i=1,2}a_i-\min_{i=1,2}b_i\right)^2}.$$

(iii) Let K be an r-fat disc and let E be an ellipse with half-axes a and b such that  $b \le a \le b^2$ . If K and E cross, then the distance between their centers is at most

$$\sqrt{\left(a-b+\sqrt{1-r^2}\right)^2+(a-r)^2}.$$

Part (ii) of the lemma is due to Heppes. The proposition in [7] refers to the special case when  $a_1 = a_2 = 1$  and  $b_i \ge b_0$ , however he actually proves claim (ii) of our lemma.

For the proof of (i) we consider two *r*-fat discs  $K_1$  and  $K_2$  with associated annuli  $C_i \setminus c_i$  centered at  $p_i$ , i = 1, 2 (see Fig. 1). We introduce Cartesian coordinates so that the coordinates of  $p_1$  and  $p_2$  are (0, -a) and (0, a), respectively. Suppose that

$$a > \sqrt{1 - r^2}.\tag{1}$$





We shall show that then  $K_1$  and  $K_2$  do not cross.

Let

$$S = \{(x, y) | -r < x < r, -a < y < a\}$$

and observe that (1) implies, on one hand, that

$$(\operatorname{bd} K_1) \cap (\operatorname{bd} K_2) \subset (C_1 \setminus c_1) \cap (C_2 \setminus c_2) \subset S,$$

and, on the other hand, that  $\operatorname{bd} K_1 \cap S$  is the graph of a concave function f(x) and  $\operatorname{bd} K_2 \cap S$  is the graph of a convex function g(x). It is then clear that  $(\operatorname{bd} K_1) \cap (\operatorname{bd} K_2)$ , if not empty, consists of a single point, of two points, or a line segment. In neither case can  $K_1$  and  $K_2$  cross. This settles part (i) of the lemma.

Consider now an *r*-fat disc *K* with associated annulus  $C \setminus c$  centered at *p*, and an ellipse *E* with half-axes *a* and *b*,  $b \le a \le b^2$  (see Fig. 2). Since  $b^2/a \ge 1$ , Proposition 1 implies that *E* and *C* do not cross. Suppose that *K* and *E* cross. Then the set  $E \setminus K$  consists of at least two connected components. As *E* and *C* do not cross, one of these components is contained in *C* (actually in  $C \setminus c$ ).

Let *M* be a connected component of  $E \setminus K$  such that  $M \subset C$ . Let  $q_1$  and  $q_2$  be two points on bd *M* dividing bd *M* into two arcs, one of which is contained in bd *K* and the other one in bd *E*. Let  $l_1$  be the line parallel to  $q_1q_2$  tangent to bd  $M \cap$  bd *E* at a point, say *q*. Let  $l_2$  be the line orthogonal to  $l_1$  through the center of *E* and let  $l_3$  be the line parallel to  $l_2$  through *p*. According to Proposition 2 the distance from *q* to  $l_2$  is at most a - b. The distance from *q* to  $l_3$  is at most  $\sqrt{1 - r^2}$ , hence the distance between the lines  $l_2$  and  $l_3$  is at most  $a - b + \sqrt{1 - r^2}$ . Now (iii) follows by noting that the difference of the distances of the centers of *E* and *K* from  $l_1$  is at most a - r.





# 4. Proof of the Theorems

We recall that the proof of Theorem 1 in the case when there are no crossing pairs is based on Proposition 3. The construction mentioned there cannot be carried out if crossings occur. We shall construct a new, crossing-free covering by successively replacing crossing pairs by other discs.

Let  $K_1, \ldots, K_N$  be congruent copies of an  $r_0$ -fat convex disc K covering the hexagon H. It will help to understand some steps in the proof if we recall the proof for the case when crossings do not occur. Then, considering the polygons  $D_i$  described in Proposition 3 and using Proposition 5, we have

area 
$$H = \sum_{i=1}^{N} \operatorname{area} D_i \le \sum_{i=1}^{N} f_K(n_i) \le \sum_{i=1}^{N} (f_K(6) + (n_i - 6)(f_K(6) - f_K(5))) \le N f_K(6),$$

which is exactly, what we have to prove.

Suppose now that two of the discs, say  $K_1$  and  $K_2$  cross. Let  $C_i \setminus c_i$  with center  $p_i$  be an annulus associated to  $K_i$  (i = 1, 2). According to the lemma, the distance between  $p_1$  and  $p_2$  is at most  $2\sqrt{1-r_0^2}$ .

Let  $E_1$  be an ellipse with half-axes

$$a_1 = 1 + \sqrt{1 - r_0^2} = 1.3598\dots$$

and

$$b_1 = \frac{r_0^2}{a_1} \left( 1 + \frac{f(6)}{\pi} \right) = 1.1695 \dots$$
 (2)

Then a congruent copy of  $E_1$ , placed so that its center is the midpoint of the segment  $p_1p_2$  and its longer axis lies on the line  $p_1p_2$ , covers  $conv(C_1 \cap C_2)$ , hence it covers  $K_1 \cap K_2$ . To see this it is enough to observe that the minimum radius of curvature of  $E_1$ ,

$$\varrho_1 = b_1^2/a_1 = 1.0057\dots$$

is greater than 1. For later reference we note that the maximum radius of curvature of  $E_1$  is

$$\overline{\varrho}_1 = a_1^2/b_1 = 1.5812\dots$$

The choice of  $a_1$  is natural. We also need the restriction that the minimum radius of curvature of  $E_1$  is at least 1. Another condition specifying the particular choice of  $b_1$  is the following. Our goal is to construct a crossing-free covering of H in which a copy of  $E_1$  substitutes two copies of K. We apply Proposition 3 to this new covering. If to a copy of  $E_1$  a polygon D with n sides is associated, we would like to have that its area is at most  $f_K(n) + f_K(6)$ . This is guaranteed by (2). Indeed, (2), Proposition 4, and the obvious inequalities  $f(n) < \pi$  and area  $K \ge \pi r_0^2$  imply

area 
$$D \le f_{E_1}(n) = a_1 b_1 f(n) = r_0^2 f(n) \left(1 + \frac{f(6)}{\pi}\right) < f_K(n) + f_K(6).$$

Replacing the discs  $K_1$  and  $K_2$  by a congruent copy of  $E_1$  covering their union we obtain again a covering of H. We repeat this process until no two discs from the original covering cross. Of course, it can occur that now two ellipses or an ellipse and one of the original discs cross. We consider first the latter case.

If a copy of  $E_1$  and a copy of K cross, then according to part (iii) of the lemma the distance between their centers is at most

$$\sqrt{\left(a_1-b_1+\sqrt{1-r_0^2}\right)^2+(a_1-r_0)^2}=0.6964\dots$$

Let  $E_2$  be an ellipse with half-axes

$$a_{2} = \frac{1}{2} \left( a_{1} + 1 + \sqrt{\left(a_{1} - b_{1} + \sqrt{1 - r_{0}^{2}}\right)^{2} + (a_{1} - r_{0})^{2}} \right) = 1.5281 \dots$$

and

$$b_2 = \frac{r_0^2}{a_2} \left( 1 + \frac{2f(6)}{\pi} \right) = 1.5118\dots$$
 (3)

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We note that the extreme values of the radius of curvature of  $E_2$  are

$$\varrho_2 = b_2^2/a_2 = 1.4956...$$
 and  $\overline{\varrho}_2 = a_2^2/b_2 = 1.5446...$ 

As  $\underline{\rho}_2 > a_1$ , the union of a unit circle and a circle of radius  $a_1$  whose centers are at a distance of  $\sqrt{(a_1 - b_1 + \sqrt{1 - r_0^2})^2 + (a_1 - r_0)^2}$  apart can be covered by a congruent copy of  $E_2$ . Therefore, the union of a crossing pair of copies of  $E_1$  and K can be covered by a copy of  $E_2$ . Using this, we successively eliminate all such crossing pairs by replacing them by an appropriate copy of  $E_2$ .

In the covering obtained thus, four different types of crossings can occur: Two copies of  $E_1$  or two copies of  $E_2$ , as well as a copy of  $E_1$  and a copy of  $E_2$  can cross. Finally, a copy of  $E_2$  and a copy of K can cross. Using the lemma and considerations similar to the above, one can see that the union of any of these types of crossing pairs of discs can be covered by a circle of radius

$$R = r_0 \sqrt{1 + \frac{3f(6)}{\pi}} = 1.7407\dots$$
 (4)

We continue to replace all crossing pairs of the system by circles of radius *R*. We end up with a covering of *H* consisting of four types of discs, in which, as can be easily checked, crossings can occur only among an  $r_0$ -fat disc from the original covering and a circle of radius *R*. If such a crossing occurs, say a copy of *K* with associated annulus  $C \setminus c$  centered at *p* and a circle  $\overline{C}$  of radius *R* centered at *q* cross, then the distance between *p* and *q* is at most

$$d = \sqrt{1 - r_0^2 + (R - r_0)^2} = 0.8842...$$

An easy computation shows that then the length of the common chord of  $\overline{C}$  and C is at most

$$\frac{1}{d}\sqrt{4R^2d^2 - (R^2 + d^2 - 1)^2} = 1.4168\ldots < 2r_0$$

and, moreover,  $\overline{C}$  contains the greater of the two segments of *C* determined by this common chord. Thus  $C \setminus \overline{C}$  can be covered by a circle of radius  $r_0$ . Now, in our covering whenever a circle of radius *R* and a copy of *K* cross, we discard the latter and introduce a circle of radius  $r_0$  to cover the part of the discarded copy of *K* that was left uncovered by the circle of radius *R*.

A circle of radius  $r_0$  introduced thus cannot cross an ellipse or a circle. Still, it can occur that such a circle and one of the original discs cross. We observe that a circle of radius  $r_0$  is itself an  $r_0$ -fat disc, therefore the union of such a crossing pair can be covered by a congruent copy of  $E_1$ , and we can start doing the whole procedure over again. Since the number of the original discs decreases in each cycle, the procedure ends in finitely many steps.

The final result of the procedure is a new covering of H consisting of, besides some of the original discs, ellipses congruent to  $E_1$  or  $E_2$ , as well as of circles of radius  $r_0$  and R. It might occur that one of these discs is contained in the union of the others. If this happens, we discard it.

As the new covering is crossing-free, we can now apply Proposition 3 to them. Let  $D_{11}, \ldots, D_{1N_1}, D_{21}, \ldots, D_{2N_2}, D_{31}, \ldots, D_{3N_3}, D_{41}, \ldots, D_{4N_4}$ , and  $D_{51}, \ldots, D_{5N_5}$  be the polygons associated to the original discs, the circles of radius  $r_0$ , the ellipses congruent to  $E_1$ , the ellipses congruent to  $E_2$ , and the circles of radius R, respectively. Denoting by  $n_{ij}, 1 \le i \le N_j, 1 \le j \le 5$ , the number of sides of  $D_{ij}$  we have

$$\sum_{i=1}^{5} \sum_{j=1}^{N_i} n_{ij} \le 6 \sum_{i=1}^{5} N_i.$$
(5)

We note that ultimately in the new covering a copy of  $E_1$  replaces two original discs, a copy of  $E_2$  replaces three copies of K, and, finally a circle of radius R replaces at least four discs from the original covering. Hence,

$$N_1 + N_2 + 2N_3 + 3N_4 + 4N_5 \le N.$$
(6)

Our next goal is to give upper bounds for the areas of polygons contained in different types of discs. We have, by definition,

area 
$$D_{1j} \leq f_K(n_{1j})$$
 for  $1 \leq j \leq N_1$ . (7)

Obviously, area  $K \ge r_0^2 \pi$ , hence in view of Proposition 4

$$r_0^2 f(n) \le f_K(n).$$
 (8)

Therefore

area 
$$D_{2j} \le r_0^2 f(n_{2j}) \le f_K(n_{2j})$$
 for  $1 \le j \le N_2$ . (9)

Further, it follows from (2), (3), (4), (8), and the obvious inequality  $f(n) < \pi$ , that

area 
$$D_{3j} \le f_{E_1}(n_{3j}) = a_1 b_1 f(n_{3j}) \le f_K(n_{3j}) + f_K(6)$$
 for  $1 \le j \le N_3$ , (10)  
area  $D_{4j} \le f_{E_2}(n_{4j}) = a_2 b_2 f(n_{4j}) \le f_K(n_{4j}) + 2f_K(6)$  for  $1 \le j \le N_4$ , (11)

and

area 
$$D_{5j} \le R^2 f(n_{5j}) \le f_K(n_{5j}) + f_K(6)$$
 for  $1 \le j \le N_5$ . (12)

Inequalities (7), (9), (10), (11), and (12) imply that

area 
$$H \le \sum_{i=1}^{5} \sum_{j=1}^{N_i} \operatorname{area} D_{ij} \le \sum_{i=1}^{5} \sum_{j=1}^{N_i} f_K(n_{ij}) + (N_3 + 2N_4 + 3N_5) f_K(6).$$
 (13)

By Proposition 5 and inequality (5) it follows that

$$\sum_{i=1}^{5} \sum_{j=1}^{N_i} f_K(n_{ij}) \le \sum_{i=1}^{5} N_i f_K(6).$$
(14)

Combining (13) and (14) and taking into account (6) we get

area 
$$H \leq N f_K(6)$$
.

Multiplying both sides by area K and rearranging we get the claim of Theorem 1.  $\Box$ 

The proof of Theorem 2 is similar. We introduce two ellipses and a circle. In their definition the function

$$g(k,n) = \frac{kf(6) + (n-6)(f(6) - f(5))}{f(n)},$$
(15)

defined for integers  $k \ge 1$  and  $n \ge 3$ , plays an important role. It is easy to check that

$$\min_{n>3} g(2,n) = g(2,7) = 1.9745\dots,$$
(16)

$$\min_{n \ge 3} g(3, n) = g(3, 8) = 2.9115...,$$
(17)

and

$$\min_{n \ge 3} g(4, n) = g(4, 9) = 3.8214\dots$$
(18)

Let  $E_1$  be an ellipse with half-axes

$$a_1 = 1 + \frac{\sqrt{5}}{2}(1 - r_1) = 1.2895..$$

and

$$b_1 = \frac{r_1}{a_1}g(2,7) = 1.1374\dots$$
 (19)

Let  $E_2$  be an ellipse with half-axes

$$a_2 = \frac{1}{2} \left( 1 + a_1 + \sqrt{(1 + a_1 - r_1 - b_1)^2 + (a_1 - r_1)^2} \right) = 1.4875 \dots$$

and

$$b_2 = \frac{r_1}{a_2}g(3,8) = 1.4503\dots$$
 (20)

Finally, let  $E_3$  be a circle of radius

$$R = \sqrt{r_1 g(4, 9)} = 1.6827\dots$$
 (21)

We note that the extreme values of the radius of curvature of  $E_1$  are

$$\underline{\varrho}_1 = b_1^2/a_1 = 1.0032...$$
 and  $\overline{\varrho}_1 = a_1^2/b_1 = 1.4620..$ 

and those for  $E_2$  are

$$\underline{\varrho}_2 = b_2^2/a_2 = 1.4140\ldots$$
 and  $\overline{\varrho}_2 = a_2^2/b_2 = 1.5257\ldots$ 

Let *E* be an  $r_1$ -fat ellipse with half-axes 1 and *b*,  $r_1 \le b \le 1$ . Using the lemma it can be checked that if two congruent copies of *E* cross, then their union can be covered by a copy of  $E_1$ , and if a copy of *E* and a copy of  $E_1$  cross, then their union can be covered by a copy of  $E_2$ . Further, if any two copies of *E*,  $E_1$ , or  $E_2$  cross, then their union can be covered by a copy of  $E_3$ .

Consider now *N* congruent copies of *E* covering a convex hexagon *H*. We construct a new, crossing-free covering of *H* consisting of congruent copies of *E*,  $E_1$ ,  $E_2$ , and  $E_3$  as follows. We start by replacing step-by-step all crossing pairs of copies of *E* by appropriate copies of  $E_1$  covering their union. We continue by successively replacing all crossing pairs of a copy of *E* and a copy of  $E_1$  by appropriate copies of  $E_2$ . If there are still crossing pairs of ellipses in the covering, we replace them, again step-by-step, by copies of  $E_3$ . Since *R* is greater than the maximum radius of curvature of any of the ellipses *E*,  $E_1$ , or  $E_2$ , no copy of  $E_3$  can cross a copy of these ellipses. Thus, the resulting system is crossing-free. Finally, we reduce the covering by discarding any disc which is contained in the union of some others.

To the resulting covering of H, consider the convex polygons described in Proposition 3. Let  $D_{11}, \ldots, D_{1N_1}, D_{21}, \ldots, D_{2N_2}, D_{31}, \ldots, D_{3N_3}$ , and  $D_{41}, \ldots, D_{4N_4}$  be the polygons associated to the copies of E,  $E_1$ ,  $E_2$ , and  $E_3$ , respectively. Denoting by  $n_{ij}$ ,  $1 \le i \le N_j$ ,  $1 \le j \le 4$ , the number of sides of  $D_{ij}$  we have

$$\sum_{j=1}^{4} \sum_{i=1}^{N_j} n_{ij} \le 6 \sum_{j=1}^{4} N_j.$$
(22)

The construction of the new covering readily implies that

$$N_1 + 2N_2 + 3N_3 + 4N_4 \le N. \tag{23}$$

We continue to give upper bounds for the areas of the polygons  $D_{ij}$ . The definition of  $f_E(n)$ , together with Proposition 4, yields that

area  $D_{1j} \le f_E(n_{1j}) \le f_E(6) + (n_{1j} - 6)(f_E(6) - f_E(5))$  for  $1 \le j \le N_1$ . (24)

Using relations (15)–(21) we get

area 
$$D_{2j} \leq f_{E_1}(n_{2j}) = a_1 b_1 f(n_{2j}) = r_1 f(n_{2j}) g(2,7) \leq b f(n_{2j}) g(2,n_{2j})$$
  
=  $2 f_E(6) + (n_{2j} - 6) (f_E(6) - f_E(5))$  for  $1 \leq j \leq N_2$ , (25)

area 
$$D_{3j} \leq f_{E_2}(n_{3j}) = a_2 b_2 f(n_{3j}) = r_1 f(n_{3j}) g(3, 8) \leq b f(n_{3j}) g(3, n_{3j})$$
  
=  $3 f_E(6) + (n_{3j} - 6) (f_E(6) - f_E(5))$  for  $1 \leq j \leq N_3$ , (26)

and

area 
$$D_{4j} \leq f_{E_3}(n_{4j}) = R^2 f(n_{4j}) = r_1 f(n_{4j}) g(4,9) \leq b f(n_{4j}) g(4,n_{4j})$$
  
=  $4 f_E(6) + (n_{4j} - 6) (f_E(6) - f_E(5))$  for  $1 \leq j \leq N_4$ . (27)

The combination of inequalities (22)–(27) readily yields

area 
$$H = \sum_{i=1}^{4} \sum_{j=1}^{N_i} \operatorname{area} D_{ij} \le f_E(6) \sum_{j=1}^{4} \sum_{i=1}^{N_j} j + (f_E(6) - f_E(5)) \sum_{j=i}^{4} \sum_{j=1}^{N_i} (n_{ij} - 6) \le N f_E(6).$$

This completes the proof of Theorem 2.

## 5. Remarks

In our theorems we gave  $r_0$  and  $r_1$  to three decimals. In fact, these are rounded up values of the solutions of the equations

$$\left(1 + \frac{f(6)}{\pi}\right)^2 r_0^4 = \left(1 + \sqrt{1 - r_0^2}\right)^3$$

and

$$(g(2,7))^2 r_1^2 = \left(1 + \frac{\sqrt{5}}{2}(1-r_1)\right)^3,$$

respectively, with which the theorems still hold. More accurate values are

$$r_0 = 0.93242333...$$
 and  $r_1 = 0.74039619...$ 

The fatness condition in Theorem 1 can be weakened a little further. It is easy to see that the bound for the distance of the centers of crossing *r*-fat discs in part (i) of the lemma is best possible. However, it can also be seen that if two congruent copies of an *r*-fat disc *K* cross and the distance between their centers is  $d \le 2\sqrt{1-r^2}$ , then area  $K \ge rd/2 + r^2(\pi - \arctan(d/2r))$ . In the proof of Theorem 1 we used the lower bound  $\pi r_0^2$  for area *K*. Using instead the information above, we can show that the density bound of Theorem 1 holds for 0.93125458-fat convex discs. We omit the proof, since it is more involved and the improvement it gives is very slight.

Let  $P_n$  denote a regular *n*-gon inscribed in a unit circle.  $P_n$  is  $r_0$ -fat for  $n \ge 9$ . The regular octagon  $P_8$  is not  $r_0$ -fat, and it misses the fatness-bound 0.93125458, as well. However, repeating the argument of the proof of Theorem 1 and taking into account that the area of  $P_8$  is considerably greater than the area of the circle inscribed into it, we can see that the density-bound of the theorem holds for  $P_8$  without the assumption that the octagons do not cross. It is easy to find the hexagon of maximum area contained in  $P_n$ . Its vertices are vertices of  $P_n$  and the difference between the lengths of its sides is as small as possible under this condition. Using this we get that

$$\vartheta(P_{6k}) = \frac{k\sin(\pi/3k)}{\sin(\pi/3)}$$

and

$$\vartheta(P_{6k\pm 2}) = \frac{(3k\pm 1)\sin(\pi/(3k\pm 1))}{2\sin(k\pi/(3k\pm 1)) + \sin((k\pm 1)\pi/(3k\pm 1))}$$

for all  $k \ge 1$ .

Mount and Silverman [8] gave an algorithm which determines the value of  $f_K(6)$  in O(n) time if K is a convex *n*-gon. Their algorithm can now be applied to determine the covering density of a centrally symmetric  $r_0$ -fat *n*-gon in O(n) time.

Finally we note, that using the method developed by Böröczky Jr. in [2], a slight modification of our proof yields that the bounds for the density in Theorems 1 and 2 hold if H is an arbitrary convex disc, provided that the number of discs is sufficiently large.

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Received December 17, 2003, and in revised form July 22, 2004, Online publication January 21, 2005.