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The Euclidean Distortion of Complete Binary Trees*

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Abstract. Bourgain [1] showed that every embedding of the complete binary tree of depth n into l_2 has metric distortion $\geq \Omega(\sqrt{\log n})$. An alternative proof was later given by Matousek [3]. This note contains a short proof for this fact.

A mapping $\varphi\colon X \to Y$ between metric spaces (X, d) and (Y, ρ) has *distortion* $\leq \gamma$ if there is a real $a > 0$, such that

$$\forall x_1, x_2 \in X, \qquad \gamma a \cdot \rho(\varphi(x_1), \varphi(x_2)) \geq d(x_1, x_2) \geq a \cdot \rho(\varphi(x_1), \varphi(x_2)).$$

Every graph G induces a metric d_G on its vertex set, where $d_G(u, v)$ is the length of the shortest path in G joining u and v. In this note, $\|\cdot\|$ denotes the l_2 norm.

Here we give a short proof of:

Theorem 1 [1]. *Every mapping of T_n, the complete binary tree of depth n, into l_2 has distortion $\geq \Omega(\sqrt{\log n})$.*

This was previously proved (in a more general form) in [1] and [3]. The bound is tight; an l_2-embedding of T_n with distortion $O(\sqrt{\log n})$ appears in [1]. For a broader discussion of graph embeddings and distortion see, e.g., [2].

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Our key tool is a simple geometric inequality. For a positive integer n, define $\Gamma = \Gamma_n$ to be the set $\{(p, q)\colon\ 0 \le p < q \le n, q - p = 2^i \text{ for some } i \ge 1\}$.

Lemma 1. *Let $x_0, \ldots, x_n$ be real vectors. Then*

$$\sum_{(p,q)\in\Gamma} \frac{\|x_p - 2x_{(p+q)/2} + x_q\|^2}{(p-q)^2} \le \sum_{p=0}^{n-1} \|x_{p+1} - x_p\|^2.$$

Proof. For real vectors, a, b, c the parallelogram identity $P(a, b, c)$ says $\|a - 2b + c\|^2 + \|a - c\|^2 = 2\|a - b\|^2 + 2\|b - c\|^2$. Summing $(1/(p-q)^2)P(x_p, x_{(p+q)/2}, x_q)$ over $(p, q) \in \Gamma$ yields

$$\sum_{(p,q)\in\Gamma} \left(\frac{\|x_p - 2x_{(p+q)/2} + x_q\|^2}{(p-q)^2} + \frac{\|x_p - x_q\|^2}{(p-q)^2} \right)$$
$$= \sum_{(p,q)\in\Gamma} \left(\frac{2\|x_p - x_{(p+q)/2}\|^2}{(p-q)^2} + \frac{2\|x_{(p+q)/2} - x_q\|^2}{(p-q)^2} \right).$$

For a, b with $1 \le a < b \le n$, the first summand on the right is $\|x_a - x_b\|^2/2(b-a)^2$ if and only if $(p, q) = (a, 2b - a) \in \Gamma$ and the second summand is $\|x_a - x_b\|^2/2(b-a)^2$ if and only if $(p, q) = (2a - b, b) \in \Gamma$. In each case, $b - a$ must be a power of 2. Therefore, the right-hand side does not exceed $\sum(\|x_a - x_b\|^2/(a-b)^2)$ where the sum is over pairs (a, b) such that $b - a$ is a power of 2. Separating the terms where $b = a + 1$, we bound the right-hand side from above by

$$\sum_{a=0}^{n-1} \|x_{a+1} - x_a\| + \sum_{(a,b)\in\Gamma} \frac{\|x_a - x_b\|^2}{(a-b)^2}.$$

Comparing this with the summation on the left-hand side yields the lemma. □

Proof of Theorem 1. Let f map $V(T_n)$ into l_2. We may assume f is nonexpansive, i.e., for every two vertices $\|f(x) - f(y)\| \le d_T(x, y)$. We seek a pair of vertices w, w' for which $\|f(w) - f(w')\|/d_T(w, w')$ is small. A *fork* in T is a quadruple of vertices $\Phi = (u, v, w, w')$, where v is a descendant of u, the least common ancestor of w, w' is v and $d_T(u, v) = d_T(v, w) = d_T(v, w')$. We let

$$\delta(\Phi) = \frac{\|f(u) - 2f(v) + f(w)\|}{d_T(u, w)} \quad \text{and} \quad \delta'(\Phi) = \frac{\|f(u) - 2f(v) + f(w')\|}{d_T(u, w')}.$$

By the triangle inequality:

$$\frac{\|f(w) - f(w')\|}{d_T(w, w')} \le \delta(\Phi) + \delta'(\Phi).$$

As in [1] and [3], the theorem follows by exhibiting a fork Φ for which $\delta(\Phi) + \delta'(\Phi) \le O(1/\sqrt{\log n})$. We do this by a simple averaging argument. We define a probability distribution over forks Φ and show that the expectation $\mathbf{E}[(\delta(\Phi))^2 + (\delta'(\Phi))^2] \le O(1/\log n)$. Hence, $\min(\|f(w) - f(w')\|/d_T(w, w')) \le O(1/\sqrt{\log n})$, as claimed.

As usual, we identify the vertices of T_n with binary strings of length $\le n$. (The root is the empty string and the two children of vertex α are $\alpha 0$ and $\alpha 1$.) Let $\beta(j)$ denote the jth prefix of $\beta \in \{0,1\}^n$, the jth node on the path from the root to the leaf β.

To select a fork randomly, independently choose β uniformly from $\{0,1\}^n$ and (p, q) uniformly from Γ_n. Define the fork $\Phi = (\beta(p), \beta((p+q)/2), \beta(q), \beta'(q))$, where $\beta'(q)$ is obtained from $\beta(q)$ by complementing the bit indexed by $1 + (p+q)/2$. By symmetry, $\delta(\Phi)$ and $\delta'(\Phi)$ are identically distributed. For any $\alpha \in \{0,1\}^n$, Lemma 1 with $x_i = f(\alpha(i))$ implies

$$\mathbf{E}[(\delta(\Phi))^2 \mid \beta = \alpha] \le \frac{1}{|\Gamma_n|} \sum_{i=0}^{n-1} \| f(\alpha(i+1)) - f(\alpha(i)) \|^2 \le O\left(\frac{1}{\log n}\right).$$

The last inequality follows since f is nonexpansive and $|\Gamma_n| = \Omega(n \log n)$. Averaging over α gives $\mathbf{E}[(\delta(T))^2 + (\delta(T'))^2] = O(1/\log n)$, as required. □

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