

Parallel LL parsing

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The original version of this article unfortunately contained two mistakes. The presentation of Definition 6 and Algorithm 8, step 3a were incorrect. The corrected parts of Definition 6 and Algorithm 8 are given below.

In Definition 6 it should read:

Definition 6 Let $G = (N, T, P, S)$ be a context-free grammar. The function $PSLS(x, y)$ for a pair of strings $x, y \in T^*$ is defined as follows:

$$\text{PSLS}(x, y) = \left\{ \alpha : \exists S \Rightarrow_{lm}^* wuA\beta \Rightarrow wxB\gamma \Rightarrow^* wxy\delta, \right. \\ \left. \begin{array}{l} w, u \in T^*, A, B \in N, \alpha, \beta, \gamma, \delta \in (N \cup T)^*, u \neq x, \\ \alpha \text{ is the shortest prefix of } B\gamma \text{ such that } \text{FIRST}_1(y) \subseteq \text{FIRST}_1(\alpha) \end{array} \right\} \\ [\dots]$$

Algorithm 8, step 3a) should read:

- (a) $D_j := \{ [Y \rightarrow \alpha.X\beta, u_j, v_j, \gamma], \text{ where } [Y \rightarrow \alpha.X.\beta, u_i, v_i, \delta] \in D_i, u_j \in \text{LAST}_q(\text{BEFORE}_q(Y)\alpha), v_j \in \text{FIRST}_k(Xv_i), \text{ and } \gamma \text{ is the shortest prefix of } X\delta \text{ such that } \gamma \Rightarrow^* a\omega, a \text{ is the first symbol of } v_j \}$.

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