# A G-Equivariant String-Net Construction 

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We dedicate this article to the memory of Krzysztof Gawedzki, in admiration for the mathematical and physical depth and breadth of his work. He pioneered higher structures in quantum field theories and used, in particular, equivariant structures and orbifold constructions to get beautiful insights [10, 18, 19].


#### Abstract

We develop a string-net construction for the (2,1)-dimensional part of a $G$-equivariant three-dimensional topological field theory based on a $G$-graded spherical fusion category. In this construction, a $G$-equivariant generalization of the Ptolemy groupoid enters. We compute the associated cylinder categories and show that, as expected, the model is closely related to the $G$-equivariant Turaev-Viro theory.


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## 1. Introduction

String-nets originated in physics from a description of topological phases of matter [23]. A mathematical formulation for string-nets was later given in [21]. The idea is to consider a vector space generated by graphs embedded into a surface and labeled with data from a spherical fusion category. Relations are given by the graphical calculus in the category, which is considered locally on the surface. This can be understood as an example of a topological field theory in terms of generators and relations in the sense of [33].

Using the description of the bicategory of $(3,2,1)$-cobordisms in terms of generators and relations [4], in [3,17] it was shown that the string-net construction of [21] can be extended to a once extended three-dimensional TQFT and that the string-net TQFT is equivalent to the once extended Turaev-Viro TQFT of $[1,2,5]$.

A natural question is whether this equivalence can be extended to a $G$ equivariant setting, for $G$ any finite group. There are several different points that need to be addressed when trying to answer this question. In [20], a version of the Levin-Wen model on surfaces with $G$-bundles was defined. However, a rigorous mathematical formulation for $G$-equivariant string-nets, in the spirit of [21], was not given. If there was a suitable $G$-equivariant string-net construction on surfaces, possibly with boundary, one should then compare it to a $G$-equivariant version of the Turaev-Viro construction. TQFTs on manifolds with $G$-bundles were defined in [29], where they are called homotopy quantum field theories (HTQFTs). A construction of a (once extended) Turaev-Viro HTQFT with input a $G$-graded spherical fusion category is given in [30,32]. Thus, there is a natural candidate to compare an equivariant string-net construction to.

In this paper, we give a mathematical definition for $G$-equivariant stringnets on compact surfaces, which are allowed to have non-empty boundary. As an algebraic input, we need a $G$-graded spherical fusion category $\mathcal{C}$. Furthermore, we show that our construction indeed reproduces the (2,1)-part of a once extended HTQFT as defined in [26]. We show that its value on objects of a suitable $G$-bordism bicategory is equivalent to the $G$-center of $\mathcal{C}$. In addition, we are able to compute string-net spaces on surfaces in terms of purely algebraic data, i.e., we will show in Sect. 6.2.3 that the string-net space on a surface is isomorphic to a certain hom-space in the $G$-center. Comparing our string-net construction to the Turaev-Viro theory of [32] is subtle, as we formulate, for reasons explained in [26, Remark 2.4], our construction in the language of bicategories and use related, but slightly different, geometric inputs. Our assumptions allow us to compute cylinder categories rather than to postulate them, cf. Remark 5.17.

An extension to three-dimensional bordisms, though, is beyond the scope of this paper. Since the results of [4] are not available for $G$-equivariant HTQFTs, such an extension would require different methods from the ones on [17].

From the point of view of applications, our construction might be interesting for the construction of correlators in orbifold rational conformal theory (RCFT). Constructions of correlators in (orbifold) RCFTs in terms of threedimensional TQFT were given in a series of papers [9,11-14]. Due to the close connection of string-nets with three-dimensional TQFTs, in [27] a string-net constructions for closed RCFT correlators based on Cardy-bulk algebra field content was given. The construction was extended to arbitrary open-closed RCFTs with fixed open-closed field content in [28] and to arbitrary openclosed RCFTs with defects in [15]. Using the string-net construction given in this paper, it seems reasonable to obtain a construction of orbifold correlators very similar to the previous ones.

This paper is organized as follows. In Sect. 2, we recall some facts about spherical fusion categories as well as about $G$-graded categories. In particular, we give in Proposition 2.3 an explicit expression for the $G$-crossing of the $G$-center of a $G$-graded spherical fusion category. In Sect. 3, we recall once extended $G$-equivariant HTQFTs and explain our definition for an equivariant bordism bicategory. Sections 4,5 and 6 are the main parts of the paper. In Sect. 4, a $G$-labeled version of the Ptolemy groupoid is introduced. This is the central technical tool for our string-net construction. The main result in this section is Theorem 4.9, where we show that the $G$-enhanced Ptolemy complex inherits from the ordinary Ptolemy complex the property of being connected and simply connected. Using the $G$-Ptolemy groupoid, in Sect. 5 we finally lay out our construction of the $G$-equivariant string-net space and show that it indeed is the $(2,1)$-part of a once extended HTQFT. In Sect. 6, we compute string-net spaces for a cylinder, pair of pants and a genus two surface with three boundary components. By doing so, we will see how the HTQFT structure of our string-net construction induces the $G$-crossing as well as the monoidal product in the $G$-center. A higher genus computation will then connect equivariant string-nets to the equivariant Turaev-Viro HTQFT.

### 1.1. Miscellaneous Notation

We fix some notation, which will be used throughout the whole paper. First, $\mathbb{K}$ will be an algebraically closed field of characteristic zero, $G$ a finite group and $B G$ its classifying space, which is an Eilenberg-MacLane space of type $K(G, 1)$.

For a small category $\mathcal{C}$, the set of objects is denoted $\mathcal{C}_{0}$ and the set of morphisms $\mathcal{C}_{1}$.

Given a graph $\Gamma$, its sets of vertices and edges will be $V(\Gamma)$ and $E(\Gamma)$. The set of half-edges incident to a vertex $v$ will be denoted $H(v)$. Given a surface $\Sigma$ and an embedded graph $\Gamma \hookrightarrow \Sigma$, the connected components of $\Gamma^{[2]}:=\Sigma \backslash \Gamma$ will be called 2-faces of $\Gamma$. If a graph $\Gamma$ is oriented, its set of oriented edges will be $E^{o r}(\Gamma)$. In addition, an edge $e$ with an orientation will be written in bold
symbols, i.e $\mathbf{e}:=(e, o r)$. The same edge with the opposite orientation will get an additional overline $\overline{\mathbf{e}}:=(e,-o r)$. A graph is called finite, if it has finitely many vertices and edges.

## 2. Categorical Preliminaries

### 2.1. Spherical Fusion Categories

We will work exclusively with spherical fusion categories. For the reader's convenience, we recall some definitions and facts. Proofs for the statements can be found in many sources, and an exhaustive textbook treatment is given in [8].

Categories $\mathcal{C}$ in this paper are always enriched in the symmetric monoidal category of finite-dimensional $\mathbb{K}$-vector spaces and abelian. In this case, we speak of $\mathbb{K}$-linear categories. Since we fix the ground field from the start, we will just speak of linear categories. A linear category is monoidal, if there is a bilinear functor $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ and an unit object $\mathbb{1}$ in $\mathcal{C}$ together with the usual associativity and unitality constraints. Without loss of generality we can assume that monoidal categories are strict, meaning that the following objects are identical $\mathbb{1} \otimes c=c=c \otimes \mathbb{1}$ and $(a \otimes b) \otimes c=a \otimes(b \otimes c)$. Assuming these strictness conditions, associativity and unitality constraints become trivial.

In addition, we take categories to be finitely semi-simple. That is, there is a finite set of isomorphism classes of simple objects, for which we choose a set $\mathrm{I}(\mathrm{C})$ of representatives, which includes the monoidal unit. Every object decomposes as a finite direct sum of simple objects and $\operatorname{End}_{\mathcal{C}}(\mathbb{1}) \simeq \mathbb{K}$. A category satisfying these properties is called a fusion category.

Furthermore, a monoidal category has right (resp. left) duals, if for any object $c \in \mathcal{C}$ there exists an object $c^{*}$ (resp. $c^{*}$ ) and morphisms

$$
\begin{equation*}
\operatorname{ev}_{c}: c^{*} \otimes c \rightarrow \mathbb{1}, \quad \operatorname{coev}_{c}: \mathbb{1} \rightarrow c \otimes c^{*} \tag{1}
\end{equation*}
$$

for right duals and

$$
\begin{equation*}
\widetilde{\mathrm{ev}}_{c}: c \otimes c^{*} \rightarrow \mathbb{1}, \quad \widetilde{\operatorname{coev}}_{c}: \mathbb{1} \rightarrow c^{*} \otimes c \tag{2}
\end{equation*}
$$

for left duals. The evaluation and coevaluation morphisms have to satisfy the snake identities

$$
\begin{equation*}
\left(\mathrm{Id}_{c} \otimes \mathrm{ev}_{c}\right) \circ\left(\operatorname{coev}_{c} \otimes \operatorname{Id}_{c}\right)=\mathrm{Id}_{c}, \quad\left(\widetilde{\mathrm{ev}}_{c} \otimes \mathrm{Id}_{c}\right) \circ\left(\mathrm{Id}_{c} \otimes \widetilde{\operatorname{coev}_{c}}\right)=\mathrm{Id}_{c} \tag{3}
\end{equation*}
$$

The category is called rigid if every object has a left and right dual. A pivotal structure on a rigid category is a monoidal natural isomorphism $\pi: \operatorname{Id} \bullet \Rightarrow(\bullet)^{* *}$. Similar to the monoidal structure, we can assume $[24$, Theorem 2.2] pivotal structures to be strict if they exist, meaning $\pi_{c}=\mathrm{Id}_{c}$. In a category with a (strict) pivotal structure, i.e., a pivotal category, we can form left and right traces of morphisms $f \in \operatorname{End}_{\mathcal{C}}(c)$

$$
\begin{equation*}
\operatorname{tr}_{\ell}(f):=\operatorname{ev}_{c^{*}} \circ\left(f \otimes \operatorname{Id}_{c^{*}}\right) \circ \operatorname{coev}_{c}, \quad \operatorname{tr}_{r}(f):=\widetilde{\mathrm{ev}}_{c^{*}} \circ\left(\operatorname{Id}_{c^{*}} \otimes f\right) \circ \widetilde{\operatorname{coev}}_{c} \tag{4}
\end{equation*}
$$

A pivotal category is called spherical if left and right traces coincide. Note that in a spherical category one can identify left and right duals, which we will implicitly do. As left and right traces agree in spherical categories, we simply speak of the trace and drop the distinction between left and right traces from notation.

In a spherical category, we can associate with $c \in \mathcal{C}_{0}$ its dimension $d_{c} \in$ $\operatorname{End}_{\mathcal{C}}(\mathbb{1}) \simeq \mathbb{K}$ defined as

$$
\begin{equation*}
d_{c}:=\operatorname{tr}\left(\operatorname{Id}_{c}\right) \tag{5}
\end{equation*}
$$

The whole category has a global dimension

$$
\begin{equation*}
D:=\sum_{i \in I(\mathcal{C})} d_{i}^{2} \tag{6}
\end{equation*}
$$

which is nonvanishing [8, Theorem 7.21.12].

### 2.2. Graphical Calculus

The graphical calculus for spherical fusion categories plays a prominent role in the string-net construction. We discuss it quickly to fix some conventions. For $c \in \mathcal{C}_{0}$, we represent $\mathrm{Id}_{c}$ by a straight line in the plane, oriented from bottom to top and labeled with $c$. Similar $\mathrm{Id}_{c^{*}}$ is represented by a $c$-labeled straight line oriented from top to bottom

$$
\operatorname{Id}_{c}=\left.\right|_{c} \quad \operatorname{Id}_{c^{*}}=\left.\right|_{c}
$$

A morphism $c \xrightarrow{f} d$ is drawn as

and composition of morphism is simply concatenation of string-diagrams. The monoidal product is represented by drawing strands form left to right. Evaluation and coevaluation morphisms are given by oriented caps and cups In the string-net construction for oriented surfaces, it turns out to be convenient to label vertices not directly by morphisms in the category, but rather by


$\mathrm{ev}_{c}$


$\widetilde{\mathrm{ev}}_{c}$
elements in a vector space that treats inputs and outputs on the same footing and depends only on the cyclic order of inputs and outputs. This vector space is constructed as a direct limit. For this limit, we need to identify different hom-spaces; for $\mathcal{C}$ strictly pivotal, such identification maps are constructed using canonical isomorphisms $\tau_{a, b}: \operatorname{Hom}_{\mathcal{C}}(\mathbb{1}, a \otimes b) \xrightarrow{\simeq} \operatorname{Hom}_{\mathcal{C}}(\mathbb{1}, b \otimes a)$ for all $a, b \in \mathcal{C}_{0}$


Pivotality implies $\tau_{a, b} \circ \tau_{b, a}=\mathrm{Id}$. Therefore, instead of looking at each morphism set $\operatorname{Hom}_{\mathcal{C}}\left(\mathbb{1}, c_{1} \otimes \cdots \otimes c_{n}\right)$ on its own, we take the limit over the diagram $\rightarrow \operatorname{Hom}_{\mathcal{C}}\left(\mathbb{1}, c_{i+1} \otimes \cdots \otimes c_{i-1} \otimes c_{i}\right)^{\tau_{c_{i+1}}, c_{i+2} \otimes \cdots \otimes c_{i}} \operatorname{Hom}_{\mathcal{C}}\left(\mathbb{1}, c_{i+2} \otimes \cdots \otimes c_{i} \otimes c_{i+1}\right) \rightarrow$

For later use, we denote the limit by $\mathcal{C}\left(c_{1}, \ldots, c_{n}\right)$. An element $f \in$ $\mathcal{C}\left(c_{1}, \ldots, c_{n}\right)$ will be represented by a circular coupon rather than a box, since it only depends on the cyclic order of $c_{1}, \ldots, c_{n}$.


There is a partial composition map for elements in the limit, induced by the maps

$$
\begin{align*}
\operatorname{Hom}_{\mathcal{C}}(\mathbb{1}, a \otimes b) \otimes \operatorname{Hom}_{\mathcal{C}}\left(\mathbb{1}, b^{*} \otimes c\right) & \rightarrow \operatorname{Hom}_{\mathcal{C}}(\mathbb{1}, a \otimes c) \\
(f, g) & \mapsto\left(\operatorname{Id}_{a} \otimes \widetilde{\mathrm{ev}}_{b} \otimes \operatorname{Id}_{c}\right) \circ(f \otimes g) \tag{7}
\end{align*}
$$

The composition will still be represented by concatenating strands in stringdiagrams.


Figure 1. Completeness relation in the semi-simple category $C$, where $d i=\operatorname{tr}(I d i)$
Using semi-simplicity for any $c \in \mathcal{C}_{0}$, we can pick a basis $\left\{\alpha_{c, i}^{k}\right\}_{k}$ in $\mathcal{C}\left(c, i^{*}\right)$ and $\left\{\alpha_{m}^{c, i}\right\}_{m}$ in $\mathcal{C}\left(c^{*}, i\right)$ for any $i \in I(\mathcal{C})$ with

and thus we have the completeness relation shown in Fig. 1. Finally we discuss 6j-symbols in $\mathcal{C}$. In a pivotal finite semi-simple category, by choosing bases in the spaces of three point couplings, the vector space $\mathcal{C}(i, j, k, \ell)$ has two distinct bases. One stems from the decomposition $\mathcal{C}(i, j, k, \ell) \simeq \bigoplus_{r} \mathcal{C}(i, j, r) \otimes$ $\mathcal{C}\left(r^{*}, k, \ell\right)$, whereas the other corresponds to the splitting $\mathcal{C}(i, j, k, \ell) \simeq \bigoplus_{s}$ $\mathcal{C}(j, k, s) \otimes \mathcal{C}\left(s^{*}, \ell, i\right)$. The entries of the transformation matrix between the two bases are the 6 j -symbols


The 6 j -symbols satisfy the usual pentagon relation.

### 2.3. G-Categories

This section is a recollection of the relevant definitions associated with categories which are $G$-graded and possibly carry a $G$-action. The reader can find
detailed accounts of $G$-categories in the literature, e.g., [29, Appendix 5] [31] and references therein. Besides recalling definitions, we give in Proposition 2.3 an alternative definition for a $G$-crossing on the $G$-center of a $G$-graded category. Though our definition is equivalent to the one in [31], we still introduce it, since it makes the discussion of string-nets on cylinders more transparent later.

Throughout the section, $\mathcal{C}$ is a $\mathbb{K}$-linear, strict monoidal category. $\bar{G}$ denotes the category having as objects the elements of $G$ and only identity morphisms. Multiplication in $G$ endows $\bar{G}$ with a strict monoidal structure.

Definition 2.1. The monoidal, linear category $\mathcal{C}$ is $G$-graded, if it decomposes into a direct sum of pairwise disjoint, full $\mathbb{K}$-linear subcategories $\left\{\mathcal{C}_{g}\right\}_{g \in G}$, such that
(i) any object $c \in \mathcal{C}$ decomposes as a finite direct sum $c=c_{g_{1}} \oplus \cdots \oplus c_{g_{n}}$, with $c_{g_{i}} \in \mathcal{C}_{g_{i}}$.
(ii) for $c \in \mathcal{C}_{g}, d \in \mathcal{C}_{h}$, it holds $c \otimes d \in \mathcal{C}_{g h}$.
(iii) $\operatorname{Hom}(c, d)=0$ for $c \in \mathcal{C}_{g}$ and $d \in \mathcal{C}_{h}$ with $g \neq h$.
(vi) $\mathbb{1} \in \mathcal{C}_{e}$.

The subcategories $\mathcal{C}_{g}$ are called homogeneous components of $\mathcal{C}$. The component $\mathcal{C}_{e}$ for the neutral element $e \in G$ is called the neutral component. There is an obvious forgetful functor $U: \mathcal{C} \rightarrow \dot{\mathcal{C}}$, which simply forgets the $G$-grading. A $G$-graded category $\mathcal{C}$ is rigid/pivotal/spherical if its underlying monoidal linear category $\dot{\mathcal{C}}$ is rigid/pivotal/spherical. It is fusion, if its underlying category is a fusion category, such that any homogeneous component contains at least one simple object. In a $G$-graded fusion category, any set of representing objects for isomorphism classes of simple objects splits as a disjoint union $I=\bigsqcup_{g \in G} I_{g}$ with $I_{g}$ the set of isomorphism classes of simples in $\mathcal{C}_{g}$. Thus, simple objects are always homogeneous.

In a $G$-graded category, the group $G$ serves as an index set, but does not act on the category so far. A $G$-action will come in the form of a $G$-crossing. For a monoidal category $\mathcal{D}$, recall that $\operatorname{Aut}_{\otimes}(\mathcal{D})$ is the category having monoidal equivalences $F: \mathcal{D} \xrightarrow{\simeq} \mathcal{D}$ as objects and monoidal natural isomorphisms as morphisms. Composition of functors equips it with a strict monoidal structure whose monoidal unit is the identity functor.

Definition 2.2. A $G$-crossed category is a $G$-graded category $\mathcal{C}$ together with a strong monoidal functor $\rho: \bar{G} \rightarrow \operatorname{Aut}_{\otimes}(\mathcal{C})$ such that $\rho_{h}\left(\mathcal{C}_{g}\right) \subset \mathcal{C}_{h^{-1} g h}$. The functor $\rho$ is called a $G$-crossing.

For any $g \in G, \rho_{g}$ is a strong monoidal functor. In addition, as $\rho$ is strong monoidal, a $G$-crossed category comes with natural isomorphisms $\left\{\eta_{h, g}(\bullet): \rho_{g} \rho_{h}(\bullet) \stackrel{\simeq}{\leftrightarrows} \rho_{h g}(\bullet)\right\}$ and $\eta_{0}(\bullet): \mathrm{Id}_{\mathcal{C}} \stackrel{\simeq}{\Longrightarrow} \rho_{e}(\bullet)$. These maps satisfy the usual coherence diagrams, which can be found in [31, section 3].

Just as it is sometimes necessary to equip a given monoidal category with the additional structure of a braiding, there is a meaningful notion of a braiding for G-crossed categories. This cannot be an ordinary braiding for
the underlying monoidal category since for non-abelian $G$ it holds in general that $\operatorname{Hom}(c \otimes d, d \otimes c)=0$ for $c, d$ in different homogeneous components of $\mathcal{C}$. However, $\operatorname{Hom}\left(c \otimes d, d \otimes \rho_{g}(c)\right)$ for $d \in \mathcal{C}_{g}$ and $c \in \mathcal{C}_{h}$ has a chance to be nonzero, as both $c \otimes d$ and $d \otimes \rho_{g}(c) \in \mathcal{C}_{h g}$. Thus, a $G$-braiding is a natural isomorphism

$$
\begin{equation*}
\left\{\beta_{c, d}: c \otimes d \rightarrow d \otimes \rho_{g}(c)\right\} \tag{8}
\end{equation*}
$$

defined for homogeneous elements $d \in \mathcal{C}_{g}, c \in \mathcal{C}_{h}$ for all $g, h \in G$ and linearly extended to all objects of $\mathcal{C}$. The natural isomorphism has to satisfy three coherence diagrams for which we refer to [31, section 3].

To a $G$-graded category, we associate its $G$-center $\mathrm{Z}_{G}(\mathcal{C})$, which is defined to be the relative center with respect to $\mathcal{C}_{e}$. To be a bit more explicit, the $G$ center has objects pairs $\left(c, \gamma_{c, \bullet}\right)$, where $c \in \mathcal{C}$ and $\gamma_{c}$ is a relative half-braiding for objects in the neutral component $\mathcal{C}_{e}$, i.e.,

$$
\begin{equation*}
\gamma_{c, \bullet}=\left\{\gamma_{c, X}: c \otimes X \rightarrow X \otimes c\right\}_{X \in C_{e}} \tag{9}
\end{equation*}
$$

is a natural isomorphism satisfying the usual hexagon relation. A morphism $\left(c, \gamma_{c}\right) \rightarrow\left(d, \gamma_{d}\right)$ is a morphism $f: c \rightarrow d$ in $\mathcal{C}$ satisfying $\left(\operatorname{Id}_{X} \otimes f\right) \circ \gamma_{c, X}=$ $\gamma_{d, X} \circ\left(f \otimes \operatorname{Id}_{X}\right)$. The $G$-center obviously is a $G$-graded category. Similar to the non-graded case, for $\mathcal{C}$ a $G$-graded spherical fusion category, the center has the structure of a $G$-modular category. Of course, for a $G$-graded category there also exists its Drinfeld center $\mathbf{Z}(\mathcal{C})$. The $G$-center can be very different from the Drinfeld center. However, the two are related via orbifolding [16, Theorem 3.5]. Since we do not need the full $G$-modular structure on $\mathrm{Z}_{G}(\mathcal{C})$, we simply refer to [31] for the definition of a $G$-modular category. We only need that $\mathrm{Z}_{G}(\mathcal{C})$ is $G$-crossed and we explicitly give the $G$-crossing.

Similar to the non-equivariant case, there is an adjunction

$$
\begin{equation*}
I: \mathcal{C} \rightleftharpoons \mathrm{Z}_{G}(\mathcal{C}): F, \quad I \dashv F \tag{10}
\end{equation*}
$$

with $F$ the forgetful functor and its adjoint functor $I$, the induction functor, defined on objects

$$
\begin{gather*}
I(c):=\bigoplus_{i \in I_{e}} i^{*} \otimes c \otimes i,  \tag{11}\\
\gamma_{I(c), X}:=\sum_{i, j \in I_{e}} d_{j} \underbrace{\alpha}_{i} \text { \& }
\end{gather*}
$$

where $I_{e}$ is a set of simple objects in $\mathcal{C}_{e}$. Its action on morphisms is the obvious one. To have the structure of an element in the graded center, the object $I(c)$ is equipped with the standard non-crossing half-braiding where $X \in \mathcal{C}_{e}$. (Recall from Fig. 1 that the pairwise appearance of $\alpha$ implies a summation over dual bases.) We will continue with this notation, as figures tend to become overloaded with notation otherwise. Due to the compatibility condition with the half-braiding, $\operatorname{Hom}_{Z_{G}(\mathcal{C})}(a, b) \subset \operatorname{Hom}_{\mathcal{C}}(a, b)$ is a proper subspace. (Here, by abuse of notation, we identify $a \in \mathrm{Z}_{G}(\mathcal{C})$ with the underlying object in $\mathcal{C}$.) An idempotent $P: \operatorname{Hom}_{\mathcal{C}}(a, b) \rightarrow \operatorname{Hom}_{\mathcal{C}}(a, b)$ with $\operatorname{Im}(P)=\operatorname{Hom}_{Z_{G}(\mathcal{C})}(a, b)$ is given by

where we introduced the cloaking circle in $\mathcal{C}$

and the crossings of the cloaking circle with the $a$ - and $b$-labeled lines correspond to the half-braidings of $a$ and $b$. The proof that $P$ is an idempotent is exactly the same as in the non-graded case (c.f. [22]). The existence of a $G$-crossing on $\mathrm{Z}_{G}(\mathcal{C})$ for $\mathcal{C}$ a $G$-fusion category was proved in [16]. An explicit expression for a $G$-crossing on $\mathrm{Z}_{G}(\mathcal{C})$ is given in [31], where $\mathcal{C}$ only needs to be a non-singular ${ }^{1}$, pivotal $G$-graded category. We only work with spherical $G$ fusion categories, which are automatically non-singular. To define a $G$-crossing, we generalize (11) and consider a functor $I^{h}: \mathrm{Z}_{G}(\mathcal{C}) \rightarrow \mathcal{C}$, which acts on objects as follows $I^{h}(c)=\bigoplus_{i \in I_{h}} i^{*} \otimes c \otimes i$. The action on morphisms is the obvious one. We want to construct a $G$-crossing on $\mathrm{Z}_{G}(\mathcal{C})$ from $I^{h}$. In order to do so, we consider the idempotent

[^0]$$
\pi_{c}^{h}: I^{h}(c) \longrightarrow I^{h}(c)
$$
$$
\pi_{c}^{h}:=\sum_{i, j \in I_{h}} \frac{d_{i}}{D} \not \uparrow c
$$
where we again use the half-braiding $\gamma_{c, i^{*} \otimes j}$ to braid the strands. It is shown in [16] that $\sum_{i \in I_{e}} d_{i}^{2}=\sum_{U \in I_{h}} d_{U}^{2}$ for all $h \in G$. Using this, it is straightforward to show that $\pi_{c}$ is indeed an idempotent. Its image is denoted $P(c):=\operatorname{Im}\left(\pi_{c}\right)$, restriction and inclusion maps are depicted as follows

$r_{c}: I^{h}(c) \longmapsto P(c)$


The image $P(c)$ has half-braiding

where $X \in \mathcal{C}_{e}$ and we use again the half-braiding $\gamma_{c, i \otimes X \otimes j^{*}}$. Thus, $\left(P(c), \gamma_{P(c), \bullet}\right) \in \mathbf{Z}_{G}(\mathcal{C})_{h^{-1} g h}$. From that, we can give an explicit $G$-crossing for $\mathrm{Z}_{G}(\mathcal{C})$.

Proposition 2.3. The maps

$$
\begin{align*}
\phi_{h}: \mathrm{Z}_{G}(\mathcal{C})_{g} & \rightarrow \mathrm{Z}_{G}(\mathcal{C})_{h^{-1} g h} \\
c & \mapsto\left(P(c), \gamma_{P(c), \bullet}\right) \tag{12}
\end{align*}
$$

constitute a $G$-crossing on $\mathbf{Z}_{G}(\mathcal{C})$.
Up to a reordering of tensor factors, the proof is the same as the one given in [31, section 4]; hence, we skip it here.

Remark 2.4. The $G$-crossing we defined is not the same as the $G$-crossing given in [31]. However, in the semi-simple case, it is not hard to show that the center equipped with the two different $G$-crossings are equivalent as $G$-crossed categories. Our definition is motivated by string-net constructions on cylinders (cf. Sect. 6.2.1).

## 3. Once Extended $G$-Equivariant HTQFTs

The modern formulation of once extended HTQFTs uses the language of bicategories and 2 -functors. A suitable definition of a symmetric monoidal bicategory $G \mathcal{B}$ ord ( $n, n-1, n-2$ ) of once extended $G$-equivariant bordisms was given in [26]. For technical reasons, we need a slight modification of this bicategory and consider pointed maps to $B G$ for objects. For convenience, we will also choose base points for one-dimensional manifolds in the following definition. As we will explain in Remark 3.2, these choices are not really essential

In the following, a manifold $M$ is pointed, if for each of its connected components a distinguished basepoint has been chosen. A map between pointed manifolds $M \xrightarrow{f} N$ is a continuous, basepoint-preserving map. In addition, once and for all we choose a basepoint $\star \in B G$.

Definition 3.1. The symmetric monoidal bicategory $\operatorname{GBord}_{\star}(3,2,1)$ is given by:
(0) Objects of $G \mathcal{B}$ ord ${ }_{\star}(3,2,1)$ are pairs $(M, f)$ of a pointed, closed, oriented one-dimensional manifold $M$ and a pointed map $f: M \rightarrow(B G, \star)$.
(1) A 1-morphism is a pair $(\Sigma, \zeta):\left(M_{0}, f_{0}\right) \rightarrow\left(M_{1}, f_{1}\right)$ consisting of an oriented, compact, collared two-dimensional manifold with boundary $\Sigma$, orientation preserving diffeomorphisms $\iota_{0}: M_{0} \times(-1,0] \rightarrow \Sigma_{0}, \iota_{1}: M_{1} \times$ $[0,1) \rightarrow \Sigma_{1}$ and a map $\zeta: \Sigma \rightarrow B G$, such that the following diagram commutes. Here $\Sigma_{0} \cup \Sigma_{1}$ is a collar for $\Sigma$. Note that no basepoint is

chosen for $\Sigma$.
(2) A 2-morphism $(W, \phi):\left(\Sigma_{0}, \zeta_{0}\right) \Rightarrow\left(\Sigma_{1}, \zeta_{1}\right)$ in $G \mathcal{B}$ ord ${ }_{\star}(3,2,1)$ is a diffeomorphism class of a three-dimensional, collared, compact oriented bor$\operatorname{dism} W: \Sigma_{0} \rightarrow \Sigma_{1}$ together with a map $\phi: W \rightarrow B G$ satisfying a compatibility diagram for the maps into $B G$.
For the precise definition of 2-morphisms, composition of morphisms and the symmetric monoidal structure in $\operatorname{GBord}_{\star}(3,2,1)$, we refer to $[26$, Definition 2.3]. 1-morphisms in $G \mathcal{B}$ ord ${ }_{\star}(3,2,1)$ are also referred to as $G$-surfaces.

Remark 3.2. Since $B G$ is path connected, there is an equivalence of symmetric monoidal bicategories $G \mathcal{B}$ ord $d_{\star}(3,2,1)$ and $\operatorname{GB} \operatorname{Bord}(3,2,1)$. The equivalence is given by the forgetful functor $G \mathcal{B}$ ord ${ }_{\star}(3,2,1) \rightarrow G \mathcal{B} \operatorname{ord}(3,2,1)$, forgetting the base point on objects. Since any map $S^{1} \rightarrow B G$ is homotopy equivalent to a basepoint-preserving map, the forgetful functor is essentially surjective on objects. Note that 1- and 2 -morphisms are literally the same for $G \mathcal{B}$ ord ${ }_{\star}(3,2,1)$ and $\operatorname{G\mathcal {B}ord}(3,2,1)$, thus the two bicategories are equivalent as symmetric monoidal bicategories.

Definition 3.3. A three-dimensional, once extended $G$-equivariant HTQFT with values in a symmetric monoidal bicategory $\mathcal{S}$ is a symmetric monoidal 2 -functor

$$
\begin{equation*}
Z: G \mathcal{B} \operatorname{ord}_{\star}(3,2,1) \rightarrow \mathcal{S} \tag{13}
\end{equation*}
$$

with the additional requirement that $Z$ depends only on the homotopy class relative to the boundary of the map $\phi: W \rightarrow B G$ for a 2-morphism $(W, \phi)$.

In our definition of the $G$-equivariant bordism bicategory and once extended $G$-equivariant HTQFTs, one has pointed maps to the classifying space as data and all compatibility diagrams strictly commute. Homotopy invariance is built in as a property of the functor $Z$. In [32], a different formulation is given. The authors set up a category $G C o b$ with objects pointed surfaces equipped with homotopy classes of pointed maps to the classifying space. Morphisms in GCob are three-dimensional cobordisms, which again come with a homotopy class of pointed maps to $B G$. Diagrams of maps to $B G$ then only have to commute up to homotopy. A comparison between the two formulations is subtle. We chose the bicategorical framework, since we want to construct the (2,1)-part of a once extended HTQFT using string-nets. Bicategories are the natural algebraic framework for such a construction.

Though the definition of a once extended HTQFT is formulated for arbitrary symmetric monoidal bicategories as targets, we will only work with the category $\mathcal{B i M o d}_{\mathbb{K}}$ of bimodules as targets. This symmetric monoidal bicategory is standard [7, Proposition 7.8.2], see, e.g., [4, Definition 2.8] for the $\mathcal{V e c t}_{\mathbb{K}}$-enriched version. Its objects are linear categories. A 1-morphism $\mathcal{C} \xrightarrow{F} \mathcal{D}$ is a linear functor

$$
\begin{equation*}
\mathcal{D}^{o p} \boxtimes \mathcal{C} \rightarrow \mathcal{V}_{\text {ect }}^{\mathbb{K}}, \tag{14}
\end{equation*}
$$

where $\mathcal{D}^{o p} \boxtimes \mathcal{C}$ has objects pairs $(d, c) \in \mathcal{D}^{o p} \times \mathcal{C}$ and $\operatorname{Hom}_{\mathcal{D}^{o p} \boxtimes \mathcal{C}}\left((d, c),\left(d^{\prime}, c^{\prime}\right)\right):=$ $\operatorname{Hom}_{\mathcal{D}^{o p}}\left(d, d^{\prime}\right) \otimes_{\mathbb{K}} \operatorname{Hom}_{\mathcal{C}}\left(c, c^{\prime}\right)$. Another name for 1-morphisms are bimodules,


Figure 2. One configuration of possible arcs in an ideal triangulation. By cutting along the edges one obtains an honest triangle
or profunctors. A 2-morphism simply is a natural transformation between lin-
 structure. Given profunctors $F: \mathcal{E}^{o p} \boxtimes \mathcal{D} \rightarrow \mathcal{V e c t}_{\mathbb{K}}$ and $G: \mathcal{D}^{o p} \boxtimes \mathcal{C} \rightarrow \mathcal{V e c t}_{\mathbb{K}}$, the composition $F \circ G: \mathcal{E}^{o p} \boxtimes \mathcal{C} \rightarrow \mathcal{V e c t}_{\mathbb{K}}$ is given by the coend

$$
\begin{equation*}
(F \circ G)(e, c):=\int^{d \in \mathcal{D}} F(e, d) \otimes G(d, c) \tag{15}
\end{equation*}
$$

## 4. G-Equivariant Ptolemy Groupoid

In this section, we introduce $G$-triangulations on surfaces, which are an enhancement of ideal triangulations of a surface. The main result is Theorem 4.9. It allows us to treat any $G$-surface as combinatorial object.

We consider an oriented, compact smooth surface $\Sigma$ with $r>0$ boundary components. We choose a distinguished point on each connected component of the boundary and denote by $\delta=\left\{\delta_{1}, \ldots, \delta_{r}\right\}$ the chosen set of points. In addition, we fix an arbitrary finite set $M \subset \Sigma \backslash \partial \Sigma$ of marked points. In case $\Sigma$ is a disk, $M$ has to contain at least one element, and in all other cases, $M$ may be empty.

Definition 4.1. [25, Definition 1.19] An ideal triangulation T of $(\Sigma, \delta, M)$ consists of a collection $\left\{\alpha_{i}\right\}_{i \in I}$ of isotopy classes of embedded arcs, having endpoints in $\delta \cup M$, such that for every boundary component $b$, its isotopy class relative to $\delta \cup M$ is contained in $\left\{\alpha_{i}\right\}_{i \in I}$ and $\Sigma-\cup_{i \in I} \alpha_{i}$ is a disjoint union of triangles.

The dual graph to an ideal triangulation is a uni-trivalent fat graph, whose cyclic order at vertices is induced by the orientation of the boundaries of dual triangles. The orientation of the boundary of a triangle is, of course, induced by the orientation of $\Sigma$.

The 1-skeleton $\mathrm{T}^{[1]}$ of T is an isotopy class of an embedded graph, and when speaking of vertices, edges and faces of an ideal triangulation, we always mean vertices, edges and faces of $\mathrm{T}^{[1]}$. Ideal triangulations are not triangulations in general, e.g., we may encounter bubble graphs like in Fig. 2.


Figure 3. Flip move along the edge $f$. The red lines show the dual fat graph

It is well known that any two triangulations of a surface can be transformed into each other by a finite sequence of $2-2$ and $3-1$ Pachner moves. Though ideal triangulations are not triangulations, there is an analog of the 2 - 2-Pachner move for ideal triangulations.

Definition 4.2. Let $f \in E(\mathrm{~T})$ be such that $f$ is adjacent to two different 2-faces of T and $f$ is not a boundary edge. A flip along $f$ is the move shown in Fig. 3.

In contrast to triangulations, there is no $3-1$ move for ideal triangulation, since we work with a fixed set of vertices. To be more precise, there exists a twodimensional CW-complex $\mathcal{P}(\Sigma, \delta, M)$, the Ptolemy complex, describing ideal triangulations on $(\Sigma, \delta, M)$. Like the Lego-Teichmüller complex, the Ptolemy complex is additional structure om $\Sigma$, which enables us to treat the surface $\Sigma$ in a combinatorial way.

Definition 4.3. The Ptolemy complex $\mathcal{P}(\Sigma, \delta, M)$ is the two-dimensional CWcomplex with

0 -cells: Vertices of $\mathcal{P}(\Sigma, \delta, M)$ are ideal triangulations.
1-cells: There is an edge between two vertices for any flip $F_{e}: \mathrm{T} \rightarrow \mathrm{T}^{\prime}$ as in Fig. 3.
2-cells: There are three types of 2-cells:
P1: For $f \in E(\mathbf{T})$ such that the flip $F_{f}$ exists, there is the 2-cell


The edge $f^{\prime}$ is the new edge appearing in the flip (see figure 3).
P2: For any two edges $e, f \in E(\mathrm{~T})$ with disjoint endpoints such that $F_{e}$ and $F_{f}$ are defined, the flips commute, i.e. there is a quadrilateral 2-cell


The flips $F_{e}^{\prime}, F_{f}^{\prime}$ are the flips performed at the edges $e$, respectively, $f$, which are now part of the new ideal triangulations $T_{1}^{\prime}$ and $T_{2}^{\prime}$.
P3: Given two edges $e, f \in E(\mathbf{T})$ sharing exactly one endpoint and $F_{e}, F_{f}$ exist, there is a pentagonal 2 -cell


Theorem 4.4 [25, Corollary V.1.1.2]. The Ptolemy complex $\mathcal{P}(\Sigma, \delta, M)$ is connected and simply connected.

Connectedness is a classical result and can be proved by purely combinatorial methods. However, the proof that $\mathcal{P}(\Sigma, \delta, M)$ is simply connected uses a fair amount of Teichmüller theory. The crucial point is that ideal triangulations, or fat graphs, give a cell decomposition of the decorated Teichmüller space $\mathcal{T}(\Sigma, \delta, M)$, which is a contractible space. In [25, Chapter 5 , Definition 1.1], the Ptolemy groupoid is defined as the path groupoid of $\mathcal{T}(\Sigma, \delta, M)$; thus, it is the fundamental groupoid of the Ptolemy complex.


Figure 4. Due to the cyclicity condition (16), the $G$-labels $g$ and $g^{\prime}$ are uniquely fixed by the $G$-labels $g_{1}, g_{2}, g_{3}, g_{4}$

So far, we discussed the classical situation of just a surface with ideal triangulations. However, in this paper, we need a Ptolemy complex which also accounts for $G$-bundles over the surface. Thus, we introduce the notion of a $G$-Ptolemy complex and show that it is still connected and simply connected. In spirit, this is a discrete realization of the holonomy functor for the $G$-bundle determined by $\zeta: \Sigma \rightarrow B G$.

Definition 4.5. A $G$-labeled ideal triangulation on $(\Sigma, \delta, M)$ is an isotopy class of an embedded oriented graph $\mathrm{T} \hookrightarrow \Sigma$ together with a map $g: E^{o r}(\mathrm{~T}) \rightarrow G$, such that
(i) the underlying graph of T obtained by forgetting the orientation is an ideal triangulation of $(\Sigma, \delta, M)$.
(ii) the map $g$ satisfies $g(\mathbf{e})=g(\overline{\mathbf{e}})^{-1}$.
(iii) if oriented edges $\mathbf{e}_{1}, \mathbf{e}_{2}, \mathbf{e}_{3}$ form the counterclockwise oriented boundary of a triangle of T the following relation holds in $G$

$$
\begin{equation*}
g\left(\mathbf{e}_{1}\right) g\left(\mathbf{e}_{2}\right) g\left(\mathbf{e}_{3}\right)=e \tag{16}
\end{equation*}
$$

The map $g$ is defined using oriented edges. Due to the condition in ii), giving the map on one orientation of the edge uniquely fixes the map on the edge with opposite orientation.

Definition 4.6. Given an edge $\mathbf{f}$ of a $G$-triangulation T , which is adjacent to two different 2-cells the G-flip along $\mathbf{f}$ is given by the transformation shown in Fig. 4.

Remark 4.7. The edge $\mathbf{f}$ at which a $G$-flip is performed cannot be a boundary edge.
$G$-triangulations are a combinatorial tool to describe 1-morphisms in the bicategory $G \mathcal{B o r d}{ }_{\star}(3,2,1)$. In $G \mathcal{B o r d} d_{\star}(3,2,1)$, all three layers of data are equipped with a map to $B G$. Equivalently one can regard $G \mathcal{B} o r d_{\star}(3,2,1)$ as the bordism 2-category of bordisms together with the datum of a $G$-principal bundle. To make this precise, recall that the mapping groupoid $\Pi(M, N)$, for
topological spaces $M, N$, has as objects continuous maps $M \xrightarrow{f} N$ and morphisms $f \rightarrow g$ are homotopy classes relative to $M \times\{0,1\}$ of homotopies $f \sim_{h} g$. For $N=B G$ the classifying space of $G$, there is a canonical equivalence of groupoids

$$
\begin{equation*}
\Pi(M, B G) \simeq \operatorname{PB}^{\operatorname{B}} n_{G}(M), \tag{17}
\end{equation*}
$$

where $\operatorname{PBB}_{G}(M)$ is the groupoid of $G$-principal bundles on $M$. By choosing a basepoint $\bullet \in M$ and restricting objects to pointed maps from $M$ to $B G$ but keeping as morphisms equivalence classes of unpointed homotopies, one gets an equivalent groupoid $\Pi_{\star}(M, B G)$. So from the pointed map $f:(M, \bullet) \rightarrow(B G, \star)$, we get that objects of $G \mathcal{B} o r d_{\star}(3,2,1)$ can be described as unions of circles with fixed $G$-principal bundles. The underlying manifold of a connected object $((M, \bullet), f)$ is diffeomorphic to a circle $S^{1}$ and its classifying map determines a homomorphism $f_{*}: \pi_{1}\left(S^{1}, \bullet\right) \simeq \mathbb{Z} \rightarrow G \simeq \pi_{1}(B G, \star)$, i.e., an element $f_{*}(1)=: g$ in $G$. In this sense, we can identify objects of $G \mathcal{B}$ ord ${ }_{\star}(M, B G)$ with finitely many circles which are labeled by a group element.

Let $(\Sigma, \zeta):\left(\left(M_{0}, \bullet\right), f_{0}\right) \rightarrow\left(\left(M_{1}, \bullet\right), f_{1}\right)$. The surface $\Sigma$ itself is not pointed; however, if $* \in M_{0}$ is a basepoint of a connected component, it holds $\zeta\left(\iota_{0}(*)\right)=f_{0}(*)=\star$. Thus, the images of the basepoints in the boundary of $\Sigma$ all get mapped to the basepoint in $B G$ by $\zeta$.

To the data of a 1-morphism $(\Sigma, \zeta):\left(\left(M_{0}, m_{0}\right), f_{0}\right) \rightarrow\left(\left(M_{1}, m_{1}\right), f_{1}\right)$, we want to associate a $G$-triangulation of $\Sigma$. We begin by choosing a finite set of points $K \subset \Sigma \backslash \partial \Sigma$, which for $\Sigma$ homeomorphic to a disk needs to be non-empty. Let $\delta:=\iota_{0}\left(m_{0}\right) \cup \iota_{1}\left(m_{1}\right) \in \partial \Sigma$ be the image of the basepoints and T an ideal triangulation based at $\delta \cup K$. Let $b \subset \partial \Sigma$ be a connected component of the boundary. Assume that the circle which is mapped to $b$ is labeled by $g \in G$. By construction, $\zeta_{*}(b) \in \pi_{1}(B G, \star)$ is the loop corresponding to $g$. Thus, $G$-labels of boundary edges in T are uniquely fixed by source and target objects of the 1-morphism. Since all images of basepoints in $\Sigma$ get mapped to $\star$, the image of a non-boundary edge $\mathbf{e} \in E^{o r}(\mathbf{T})$ with both endpoints in $\delta$ is an oriented loop in $\pi_{1}(B G, \star)$ and hence determines a group element $g_{\mathrm{e}}$. The $G$-color of all other edges is determined up to certain gauge transformations. Let $p=\left(\mathbf{e}_{1}, \ldots, \mathbf{e}_{r}\right)$ be an oriented edge path in T with endpoints in $\delta$. The path $p$ gets mapped to a loop in $B G$, based at $\star$. Hence, there exists $g_{p} \in G$, such that $\operatorname{Im}\left(\left.\zeta\right|_{p}\right) \simeq g_{p}$ holds. We can assume that only the first and last edge of the path have endpoints in $\delta$. To each edge $\mathbf{e}_{i} \in p$, we assign a group element $g_{i}$, such that

$$
\begin{equation*}
g_{1} g_{2} \cdots g_{r}=g_{p} \tag{18}
\end{equation*}
$$

This coloring is not unique. We can change $g_{i} \mapsto g_{i} h$ and $g_{i+1} \mapsto h^{-1} g_{i+1}$ and the new color still satisfies (18). These are so-called gauge transformations. A specific labeling of the oriented edges of T with group elements is compatible with $\zeta$, if (18) holds for all possible edge paths with endpoints in $\delta$. The set of all $G$-colorings of T which are compatible with $\zeta$ is denoted by $\operatorname{Col}_{G}(\mathrm{~T})$. A
gauge transformation then is a map

$$
\begin{equation*}
\lambda: M \rightarrow \operatorname{Maps}\left(G, \operatorname{End}\left(\operatorname{Col}_{G}(\mathrm{~T})\right)\right), \tag{19}
\end{equation*}
$$

where for $p \in M$ and $g \in G, \lambda_{g}(p)$ acts as


Multiplication in $G$ endows the set of gauge transformations with a group structure such that the gauge group acts transitively on $\mathrm{Col}_{G}(\mathrm{~T})$. To summarize, for any 1 -morphism $(\Sigma, \zeta)$ in $G \mathcal{B} o r d_{\star}(3,2,1)$ and any ideal triangulation T of $\Sigma$, based at the images of the basepoints and a possibly empty set $M$, we get a $G$-triangulation, unique up to gauge transformation.

Definition 4.8. Let $(\Sigma, \zeta):\left(M_{0}, f_{0}\right) \rightarrow\left(M_{1}, f_{0}\right)$ be 1 -morphism in $G \mathcal{B o r d}_{\star}(3,2,1)$ and $M \subset \Sigma \backslash \partial \Sigma$ a finite set of points. The $G$-equivariant Ptolemy groupoid $\mathcal{P}^{G}(\Sigma, \zeta, M)$ is a two-dimensional CW-complex defined as follows.

0-cells: Vertices of $\mathcal{P}^{G}(\Sigma, \zeta, M)$ are (ideal) $G$-triangulations induced by $\zeta$.
1-cells: There is an edge $\mathrm{T} \rightarrow \mathrm{T}^{\prime}$ in $\mathcal{P}^{G}(\Sigma, \zeta, M)$ for any $G$-flip and any gauge transformation.
2-cells: Since the $G$-label for $G$-flips is uniquely determined, we lift the relations P1-P3 of the Ptolemy groupoid to $\mathcal{P}^{G}(\Sigma, \zeta, M)$. This yields three relations GP1, GP2 and GP3. In addition, there is a mixed relation as well as two pure gauge relations.


Figure 5. Flip commutes with gauge transformation

GP4: For $\mathbf{e}$ an edge of a G-triangulation and $v$ a vertex incident to $\mathbf{e}$, the flip along e and a gauge transformation at $v$ commute. That is, there is the quadrilateral 2-cell shown in Fig. 5
GP5: Gauge transformations at different vertices commute, i.e., for $v \neq u$ vertices of a $G$-triangulation, there are quadrilateral 2 -cells corresponding to $\lambda(v) \lambda(u)=\lambda(u) \lambda(v)$.
GP6: Finally we add triangular 2-cells $\lambda_{g}(v) \cdot \lambda_{h}(v)=\lambda_{g h}(v)$, implementing the gauge action.


Theorem 4.9. The $G$-equivariant Ptolemy groupoid $\mathcal{P}^{G}(\Sigma, \zeta, M)$ from Definition 4.8 is connected and simply connected.

We defer the proof of Theorem 4.9 to appendix, since it does not offer any deeper insight for the main results of this paper.

## 5. Bare String-Net Spaces and Cylinder Categories

### 5.1. A Vector Space for Surfaces

We have set the stage for constructing a $G$-equivariant string-net space for 1morphisms. The main idea is to start with an ideal $G$-triangulation T on $(\Sigma, \zeta)$ and define a string-net space relative to it by only allowing graphs transversal to T . We proceed by subsequently taking quotients implementing the rules of the graphical calculus for $\mathcal{C}$ inside disks having different relative positions to
T. This will yield string-net spaces $\operatorname{SN}_{\mathrm{T}}^{\mathcal{C}}(\Sigma, \zeta)$ for any $G$-triangulation T. In order to get a space which solely depends on $(\Sigma, \zeta)$, we define for any edge of $\mathcal{P}^{G}(\Sigma, \zeta)$ an isomorphism between string-nets spaces. These maps will satisfy the relations GP1-6, and we get a projective system of vector spaces. Taking the limit of the system will ultimately give the string-net space.

At first sight, this might appear as an overly cumbersome way of defining a string-net space on a $G$-surface. In the usual string-net approach one just considers all embedded $\mathcal{D}$-colored graphs, for $\mathcal{D}$ any spherical fusion category. However, there seems to be no way of deciding whether an arbitrary $\mathcal{C}$-colored, embedded graph on $\Sigma$ is compatible with the $G$-structure on $\Sigma$. This appears to be only possible if the graph is fine enough, meaning that it is at least isotopic to the 1 -skeleton of a CW-decomposition of $\Sigma$. In other words, in the equivariant case string-net graphs need to be sensitive to the global topology of $\Sigma$. Instead of arbitrary CW-decompositions, we work with ideal triangulations, because there we have full control over the combinatorics involved. Without further ado, we now spell out the details of the construction.

We choose an arbitrary finite, but fixed set of points $M \subset \Sigma \backslash \partial \Sigma$. Note that $M$ can be empty except for the disk.

Definition 5.1. Let $\Gamma$ be an arbitrary embedded finite oriented graph in $\Sigma$ and $\mathcal{C}$ a $G$-graded fusion category. A $\mathcal{C}$-coloring of $\Gamma$ comprises two functions. First

$$
\begin{equation*}
c: E^{\mathrm{or}}(\Gamma) \rightarrow \mathcal{C}_{0}^{\mathrm{hom}} \tag{20}
\end{equation*}
$$

assigns to each oriented edge ( $e$, or) an homogeneous object of $\mathcal{C}$, such that $c(\mathbf{e})=c(\overline{\mathbf{e}})^{*}$. The second function is a map

$$
\begin{align*}
\phi: V(\Gamma) & \rightarrow \mathcal{C}_{1} \\
v & \mapsto \phi_{v} \in \mathcal{C}\left(\bigotimes_{\mathbf{e} \in H(v)} c(\mathbf{e})^{\epsilon_{\mathbf{e}}}\right) \tag{21}
\end{align*}
$$

where $\epsilon_{\mathbf{e}}=1$ if $\mathbf{e}$ is oriented away from $v$ and $\epsilon_{\mathbf{e}}=-1$ if it is oriented toward $v$. A negative exponent for an object in $c \in \mathcal{C}$ indicates its dual.

In particular, via the grading map $p: \mathcal{C}_{0}^{\text {hom }} \rightarrow G$ we get a $G$-labeling $g_{\Gamma}$ of $\Gamma$ satisfying $g_{\Gamma}(\overline{\mathbf{e}})=g_{\Gamma}(\mathbf{e})^{-1}$. The graph and its $G$-labeling should be compatible with a $G$-triangulation in a sense we are about to define (cf. Definition 5.3).

Definition 5.2. Let $\mathrm{T} \in \mathcal{P}^{G}(\Sigma, \zeta, M)$. An embedded graph $\Gamma \hookrightarrow \Sigma$ is totally transversal to T if $V(\mathrm{~T}) \cap \Gamma=\mathrm{T} \cap V(\Gamma)=\emptyset$ and any edge of T that is not on the boundary intersects at least one edge of $\Gamma$. Furthermore, all intersections are transversal.

The obvious example of a graph transversal to an ideal triangulation is its dual trivalent fat graph. Let $\Gamma$ be a $\mathcal{C}$-colored graph, which is totally transversal to the underlying ideal triangulation of a $G$-triangulation T. Pick a representative in the isotopy class for $T$. For an edge $\mathbf{e} \in E^{\text {or }}(T)$, consider the edges $\left\{\mathbf{f}_{1}, \ldots, \mathbf{f}_{n}\right\}$ of $\Gamma$ intersecting with $\mathbf{e}$, as in Fig. 6. The edge $\mathbf{f}_{i}$ has


Figure 6. In the color version, the orange line represents an edge of T with $G$-label $g$. The black lines are edges $(\Gamma, c)$
$G$-color $g_{i}$. The orientation of $\mathbf{e}$ gives a linear order on the set of intersecting edges. The transversal intersection can be positive or negative, depending on the orientation of the $\mathbf{f}_{i}$, e.g., the intersection $\mathbf{e} \cap \mathbf{f}_{1}$ is positive, whereas $\mathbf{e} \cap \mathbf{f}_{n-1}$ is negative in Fig. 6. The linear order allows us to multiply the $G$-labels of the $\mathbf{f}_{i}$ 's taking orientations into account. For the graph in Fig. 6, this gives

$$
\begin{equation*}
G \ni m_{\mathbf{e}}=g_{1} \cdots g_{n-1}^{-1} g_{n} . \tag{22}
\end{equation*}
$$

It is easy to check that $m_{\mathbf{e}}$ is well defined, i.e., independent of the chosen representative in the isotopy class for T .

Definition 5.3. Let $\Gamma$ be a $\mathcal{C}$-colored graph which is totally transversal to the underlying ideal triangulation $G$-triangulation T. Then $\Gamma$ is $G$-transversal to T if

$$
\begin{equation*}
m_{\mathbf{e}}=g(\mathbf{e}) \tag{23}
\end{equation*}
$$

holds for any edge of T .
If the $G$-triangulation is clear from the context we sometimes just say that an embedded graph is $G$-transversal.

We are working on surfaces with boundary; therefore, string-net spaces will depend on a choice of boundary condition, or boundary value.

Definition 5.4. (1) A boundary value for a surface $(\Sigma, \zeta)$ is a disjoint union of finitely many points $B \in \partial \Sigma$ together with a map $\mathbf{B}: B \rightarrow \mathcal{C}_{0}^{h o m}$.
(2) The boundary value of a $G$-transversal graph $\Gamma$ on $\Sigma$ is defined as the disjoint union of intersection points $B_{\Gamma}$ of the graph with $\partial \Sigma$, together with the map $\mathbf{B}: B_{\Gamma} \rightarrow \mathcal{C}_{0}^{\text {hom }}$, mapping an intersection point to the $\mathcal{C}$-color of its corresponding edge.

Remark 5.5. Note that one can state the definition for a boundary value in case of surfaces without the map $\zeta$ in the exact same fashion. $\zeta$ enters the definition only in so far, as restricting the possible boundary values, as $G$-labels of edges have to compatible with $\zeta$.

Definition 5.6. Let $\operatorname{VGraph}_{G T}^{\mathcal{C}}(\Sigma, \zeta, M ; \mathbf{B})$ be the $\mathbb{K}$-vector space freely generated by all graphs which are $G$-transversal to T with boundary value $\mathbf{B}$.

Remark 5.7. The vector space $\operatorname{VGraph}_{G T}^{\mathcal{C}}(\Sigma, \zeta, M ; \mathbf{B})$ is huge, e.g., isotopic graphs with the same $\mathcal{C}$-coloring are considered as different graphs so far.


Figure 7. $\gamma$ is the subgraph inside $D$ of a $G$-transversal graph $\Gamma$

Isotopy invariance will follow only after requiring the graphical calculus for $\mathcal{C}$ to hold locally on disks.

In a first step, we consider embedded closed disks $D \hookrightarrow \Sigma \backslash T$. Given a $G$-transversal graph $\Gamma$ the boundary of the disk $D$ is required to intersect edges of $\Gamma$ transversally and must not intersect $\Gamma$ at vertices of $\Gamma$. As usual, we get a cyclic set of objects $\left\{c_{i}\right\}$ from the $\mathcal{C}$-color of the edges intersecting $\partial D$ and a linear evaluation map [21, Theorem 2.3]

$$
\begin{equation*}
\langle\bullet\rangle_{D}: \operatorname{VGraph}_{G \mathrm{~T}}^{\mathcal{C}}(\Sigma, \zeta, M ; \mathbf{B}) \rightarrow \mathcal{C}\left(\bigotimes_{i} c_{i}^{\epsilon_{i}}\right) \tag{24}
\end{equation*}
$$

which is defined on vectors meeting the transversality requirements with respect to $D$. The sign conventions are obvious from Fig. 7

Definition 5.8. Let $D \hookrightarrow \Sigma \backslash T$ be an embedded disk. An element $\Gamma:=\sum_{i} x_{i} \Gamma_{i} \in$ $\operatorname{VGraph}_{G \mathrm{~T}}^{\mathcal{C}}(\Sigma, \zeta, M ; \mathrm{T})$ is null with respect to $D$ if

- all $\Gamma_{i}$ 's are transversal to $D$.
- $\left.\Gamma_{i}\right|_{\Sigma \backslash D}=\left.\Gamma_{j}\right|_{\Sigma \backslash D}$ for all $i, j$
- $\langle\Gamma\rangle_{D}=0$.

The vector space of T -disk null graphs $\mathrm{NGraph}_{G \mathrm{~T}}^{\mathcal{C}}(\Sigma, \zeta, M ; \mathbf{B})$ is the subspace of $\operatorname{VGraph}_{G T}^{\mathcal{C}}(\Sigma, \zeta, M ; \mathbf{B})$ spanned by all null graphs for all possible disks $D \hookrightarrow \Sigma \backslash T$.

The quotient of $\operatorname{VGraph}_{G T}^{\mathcal{C}}(\Sigma, \zeta, M ; \mathbf{B})$ by the vector space of T-disk null graphs $\operatorname{NGraph}_{G \mathrm{~T}}^{\mathcal{C}}(\Sigma, \zeta, M ; \mathbf{B})$ is denoted

$$
\begin{equation*}
\mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime \prime}(\Sigma, M ; \mathbf{B}):=\frac{\operatorname{VGraph}_{G \mathrm{~T}}^{\mathcal{C}}(\Sigma, \zeta, M ; \mathbf{B})}{\operatorname{NGraph}_{G \mathrm{~T}}^{\mathcal{C}}(\Sigma, \zeta, M ; \mathbf{B})} \tag{25}
\end{equation*}
$$



Figure 8. Cyclicity condition for the color at a vertex


Figure 9. The two string-nets $\Gamma, \Gamma^{\prime}$ agree outside of the neighborhood shown here. The move consists of isotoping the string-net vertex trough the edge $\mathbf{f}$ of $\mathbf{T}$

Remark 5.9. Due to the $G$-grading of $\mathcal{C}$, nonzero vectors in $\mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime \prime}(\Sigma, M ; \mathbf{B})$ satisfy a cyclicity condition at all of their vertices. That is, given $v \in V(\Gamma)$, the $\mathcal{C}$-color of its incident edges $\left\{\mathbf{e}_{i}\right\}=E^{\text {or }}(v)$ has to satisfy

$$
\begin{equation*}
p\left(c\left(\mathbf{e}_{1}\right)^{\epsilon_{1}}\right) \cdots p\left(c\left(\mathbf{e}_{n}^{\epsilon_{n}}\right)\right)=e \tag{26}
\end{equation*}
$$

There is a sign convention involving the orientation of the incident edges, which can be easily deduced from Fig. 8 and is similar to the one used in Fig. 7.

The space $\mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime \prime}(\Sigma, M ; \mathbf{B})$ only partly achieves the goal of globalizing the graphical calculus for $\mathcal{C}$ to $\Sigma$. So far, we have imposed local relations inside 2 -faces of T . In order to define local relations everywhere on $\Sigma$, we take a further quotient. This will allow for relations inside disks intersecting edges of the ideal triangulation T . We introduce an equivalence relation implementing isotoping a vertex of $\Gamma$ through an edge of $T$.

Definition 5.10. Let $\Gamma, \Gamma^{\prime} \in \mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime \prime}(\Sigma, M ; \mathbf{B})$, then $\Gamma \sim_{\mathbf{f}} \Gamma^{\prime}$ if and only if the two graphs are related by the move shown in Fig. 9.

Note that due to the cyclicity condition (26), the relation $\sim_{\mathbf{f}}$ is well defined. We take the quotient of $\mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime \prime}(\Sigma, M ; \mathbf{B})$ by $\sim_{\mathbf{f}}$ for all edges $\mathbf{f}$ of T and denote it

$$
\begin{equation*}
\mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime}(\Sigma, M ; \mathbf{B}) \tag{27}
\end{equation*}
$$

In $\mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime}(\Sigma, M ; \mathbf{B})$, local relations hold inside all disks not meeting the images of the marked points. As a consequence, boundary points of string-nets cannot be moved.

Finally, the graphs still depend on the chosen $G$-triangulation T. We get rid of this dependence next. Let $\Delta$ be a 2 -face of T and we take its counterclockwise oriented boundary $\left\{\mathbf{e}_{1}, \mathbf{e}_{2}, \mathbf{e}_{3}\right\}$ with $G$-color $\left\{g_{1}, g_{2}, g_{3}\right\}$. The 2 -face $\Delta$ may contain a boundary edge, which we assume to be $\mathbf{e}_{1}$. We denote $\mathbf{B}_{\Delta}$ for the boundary value restricted to that boundary component. So for a 2-face with boundary edge, we consider the vector space

$$
\begin{equation*}
\operatorname{Hom}_{\Delta}(\mathcal{C}):=\bigoplus_{i \in I_{g_{2}}, j \in I_{g_{3}}} \operatorname{Hom}_{\mathcal{C}}\left(\mathbb{1}, \mathbf{B}_{\Delta} \otimes j \otimes k\right) \tag{28}
\end{equation*}
$$

For all other 2-faces, we set

$$
\begin{equation*}
\operatorname{Hom}_{\Delta}(\mathcal{C}):=\bigoplus_{i \in I_{g_{1}}, j \in I_{g_{2}}, k \in I_{g_{3}}} \operatorname{Hom}_{\mathcal{C}}(\mathbb{1}, i \otimes j \otimes k) \tag{29}
\end{equation*}
$$

and define a vector space

$$
\begin{equation*}
H_{\mathrm{T}}^{\mathcal{C}}(\Sigma ; \mathbf{B}):=\bigotimes_{\Delta \in \mathrm{T}^{[2]}} \operatorname{Hom}_{\Delta}(\mathcal{C}) \tag{30}
\end{equation*}
$$

where $T^{[2]}$ denotes the set of 2-faces of $T$.
Lemma 5.11. There is an isomorphism

$$
\begin{equation*}
\Psi: H_{\mathrm{T}}^{\mathcal{C}}(\Sigma ; \mathbf{B}) \rightarrow \mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime}(\Sigma, M ; \mathbf{B}) \tag{31}
\end{equation*}
$$

which maps an element $\bigotimes_{f \in \pi_{0}(\Delta)} \phi_{f} \in H_{\mathrm{T}}^{\mathcal{C}}(\Sigma ; \mathbf{B})$ to the equivalence class of the dual fat graph $\Gamma$ with boundary value $\mathbf{B}_{\Gamma}=\mathbf{B}$ and whose vertices are colored by the maps $\phi_{f}$.

Among other things Lemma 5.11 tells that $\mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime}(\Sigma, M ; \mathbf{B})$ has a basis in terms of equivalence classes of fat graphs with edges meeting $\partial \Sigma$ which are dual to T . Internal edges of the basis elements are colored with simple objects and its vertices are colored with basis elements in the corresponding hom-spaces $H_{\mathrm{T}}^{\mathcal{C}}(\Sigma ; \mathbf{B})$ via the map $\Psi$. We call this the fat graph basis for T . The proof of Lemma 5.11 is very similar to the proof of [21, Lemma 5.3], and we leave it to the reader to adapt it to the present situation.

We almost arrived at a sensible definition of a string-net space. A remaining issue is that edges of a string-net cannot be moved through the marked points $M$. This issue can be approached by cloaking marked points, which is defined by projecting to a subspace of $\mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime}(\Sigma ; \mathbf{B})$. Let

$$
\begin{equation*}
\Pi_{g}: \mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime}(\Sigma, M ; \mathbf{B}) \rightarrow \mathrm{SN}_{\mathrm{T}, \mathcal{C}}^{\prime}(\Sigma, M ; \mathbf{B}) \tag{32}
\end{equation*}
$$

be the map acting in a neighborhood of $M$ by Fig. 10.
For the neutral element $e \in G$, we simply write $\Pi:=\Pi_{e}$ and depict it with a purple circle without group label It is easy to see that $\Pi^{2}=\Pi$ and we set

$$
\begin{equation*}
\mathrm{SN}_{\mathrm{T}}^{\mathcal{C}}(\Sigma, M ; \mathbf{B}):=\operatorname{Im}(\Pi) . \tag{33}
\end{equation*}
$$



Figure 10. Action of the map $\Pi_{g}$ on a vertex $s$ of an ideal triangulation





Figure 11. We can use completeness in $\mathcal{C}$ to pull a string-net edge through a vertex of T. However, the color of the cloaking circle might change

Using the completeness relation of Fig. 1, the cloaking circle allows us to move edges of string-nets through marked points. If the edge has non-trivial $G$-color, though, e.g., $c \in \mathcal{C}_{g}$ in Fig. 11, the color of the cloaking circle might change. The resulting string-net is still in the image of $\Pi$, as can be easily checked using the completeness relation again.

The vector space $\mathrm{SN}_{T}^{\mathcal{C}}(\Sigma, M ; \mathbf{B})$ clearly depends on the choice of a $G$ triangulation. Let $(\mathrm{T}, g) \xrightarrow{F_{f}}\left(\mathrm{~T}^{\prime}, g^{\prime}\right)$ be a $G$-flip, where none of the edges of
the two triangles is a boundary edge. We define the associated linear map

$$
\begin{equation*}
F_{\mathrm{T}, \mathrm{~T}^{\prime}}: \mathrm{SN}_{\mathrm{T}}^{\mathcal{C}}(\Sigma, M ; \mathbf{B}) \rightarrow \mathrm{SN}_{\mathrm{T}^{\prime}}^{\mathcal{C}}(\Sigma, M ; \mathbf{B}) \tag{34}
\end{equation*}
$$

as the linear extension of the map defined on elements of the fat graph basis by

where the map is the identity outside the open neighborhood of the edge $\mathbf{f}$ shown in red in the picture. In case there is a boundary edge, we use first use the completeness relation to decompose the $\mathcal{C}$-colored edges ending on $\partial \Sigma$ into a sum of simple colored edges and then perform the same move as in the previous case, followed by the inverse of the completeness relation.

To a gauge transformation $\mathrm{T} \xrightarrow{\lambda_{g}(v)} \mathrm{T}^{\prime}$, we associate the map $G_{g}(v)$ : $\mathrm{SN}_{\mathrm{T}}^{\mathcal{C}}(\Sigma, M ; \mathbf{B}) \rightarrow \mathrm{SN}_{\mathrm{T}^{\prime}}^{\mathcal{C}}(\Sigma, M ; \mathbf{B})$ adding an $I_{g}$-colored cloaking circle around the internal vertex $v$ of the $G$-triangulation $T$.
Theorem 5.12. There is a functor

$$
\begin{align*}
\mathrm{SN}^{\mathcal{C}}: \Pi_{1}\left(\mathcal{P}^{G}(\Sigma, \zeta)\right) & \rightarrow \mathcal{V}_{e c t} \\
(\mathrm{~T}, g) & \mapsto \mathrm{SN}_{\mathrm{K}}^{\mathcal{C}}(\Sigma, M ; \mathbf{B}) \\
{\left[\mathrm{T} \xrightarrow{F_{e}} \mathrm{~T}^{\prime}\right] } & \mapsto F_{\mathrm{T}, \mathrm{~T}^{\prime}}  \tag{35}\\
{\left[\mathrm{T} \xrightarrow{\lambda_{g}(v)} \mathrm{T}^{\prime}\right] } & \mapsto G_{g}(v)
\end{align*}
$$

Proof. The linear maps $F_{\mathrm{T}, \mathrm{T}^{\prime}}, G_{g}(v)$ are indeed isomorphisms. To see the first statement, observe that $F_{\mathrm{T}, \mathrm{T}^{\prime}}$ is its own inverse. In case of the gauge map, this
is a consequence of the completeness relation in $\mathcal{C}$. Thus, we only have to check that relations GP1-GP6 are mapped to an equivalence of linear isomorphisms. Relation GP1 is exactly the statement that $F_{\mathrm{T}, \mathrm{T}^{\prime}}$ is its own inverse. The GP2 equivalence follows, since we are manipulating disjoint parts of the string-net, which obviously commutes. Next, GP3 gets mapped to the pentagon relation for the $6 j$-symbols. Relation GP4 is an easy computation using the graphical calculus in $\mathcal{C}$. Since adding cloaking circles at different internal vertices commutes, relation GP5 holds. The group law GP6 follows by applying the completeness relation twice.

Definition 5.13. Given a $G$-surface $(\Sigma, \zeta)$, a set of marked points $M$ and a $G$ graded spherical fusion category $\mathcal{C}$, the based bare string-net space is defined by

$$
\begin{equation*}
\mathrm{SN}^{\mathcal{C}}(\Sigma, M ; \mathbf{B}):=\lim _{\Pi_{1}\left(\overleftarrow{\left.\mathcal{P}^{G}(\Sigma, \zeta)\right)}\right.} \mathrm{SN}_{\mathrm{T}}^{\mathcal{C}}(\Sigma, M ; \mathbf{B}) \tag{36}
\end{equation*}
$$

Remark 5.14. Instead of taking the limit over the diagram, we could have instead taken its colimit. The resulting vector spaces are isomorphic, since the diagram involves only linear isomorphisms. For the same reason, we have a canonical isomorphism

$$
\begin{equation*}
\mathrm{SN}^{\mathcal{C}}(\Sigma, M ; \mathbf{B}) \cong \mathrm{SN}_{\mathrm{T}}^{\mathcal{C}}(\Sigma, M ; \mathbf{B}) \tag{37}
\end{equation*}
$$

for any $G$-triangulation T.

The bare string-net space based at $M$ only depends on the $G$-surface $(\Sigma, \zeta)$, the set of marked points $M$ and $\mathcal{C}$, but not on the choice of a $G$ triangulation based at $\delta \cup M$. The set $\delta$ will be uniquely determined by the input datum of a 1 -morphism. Thus, the only arbitrary choice left is the set $M$ of marked points. However, there is a distinguished isomorphism between based bare string-net spaces for different choices of marked points. We show this by showing that for all based string-net spaces, there is a specific isomorphism to the based bare string-net space with $M=\emptyset$, or $M$ a single point in case of a disk. For a given $M$, we pick a $G$-triangulation, and from Lemma 5.11, we get that $\mathrm{SN}^{\mathcal{C}}(\Sigma, M ; \mathbf{B})$ has a fat graph basis with additional cloaking circles around elements in $M$. To make arguments more palpable, we discuss the specific case of a genus 2 surface $\Sigma$ with 4 boundary components. We choose an arbitrary point $\star \in \Sigma \backslash M$ and a set of generators in $\pi_{1}(\Sigma \backslash M, \star)$


For the dual fat graph $\Gamma$ of the ideal triangulation $T$, we pick a maximal tree $t$ and denote $\ell:=E(\Gamma)-E(t)$. Since the fat graph $\Gamma$ generates $\pi_{1}(\Sigma \backslash M, \star)$, we can in fact choose $t$ in such a way that upon collapsing $t$ to $\star$, the homotopy classes of the remaining edges $\ell$ agree with the chosen generators of $\pi_{1}(\Sigma \backslash M, \star)$. This means that for any of the generators $\gamma$, there exists an element $l_{\gamma} \in \ell$, such that after collapsing $t$ to $\star$ it holds that $\left[l_{\gamma}\right]=\gamma$. The tree $t$ has a disk-shaped neighborhood. Using local relations on this disk, we can replace the elements of the fat graph basis with string-nets supported on the contracted graphs. These string-nets still constitute a basis for the based bare string-net space. We use the completeness relation, possibly followed by a gauge transformation isomorphism, to map these new basis elements to stringnets where the points of $M$ are solely surrounded by cloaking circles


After further contracting the graph, we obtain

where we can restrict to colorings by simple objects. Thus, for any two choices of marked points $M_{1}, M_{2}$, by identifying the basis elements, we have a distinguished isomorphism between string-net spaces defined in terms of $G$-triangulations based at $\delta \cup M_{1}$ and $\delta \cup M_{2}$. The isomorphisms give a projective system of vector spaces and we can take its projective limit, which we denote by $\mathrm{SN}^{\mathcal{C}}(\Sigma ; \mathbf{B})$.
Definition 5.15. The vector space $\mathrm{SN}^{\mathcal{C}}(\Sigma ; \mathbf{B})$ is called bare string-net space.
To summarize our construction, we first needed a combinatorial replacement for a surface $\Sigma$ with map $\zeta: \Sigma \rightarrow B G$, in order to give a sensible definition for a string-net graph in $\Sigma$ to respect the $G$-structure. This point was settled by using $G$-triangulations. Once we came up with a definition of a pre-string-net space depending on a given $G$-triangulation, the whole rest of the section was devoted to build up a definition depending only on $\Sigma$ and $\zeta$, but not on the combinatorial model. Taking limits over projective systems only involving isomorphisms has the advantage that, whenever we want to work with the bare string-net space, we can just pick our favorite $G$-triangulation and compute everything relative to it. We will make use of this fact in Sect. 6.2, when actually computing string-net spaces.

Remark 5.16. Allowing an additional set $M$ of vertices in an ideal triangulation will simplify the description of the behavior of string-net spaces under gluing of $G$-surfaces described in Sect. 6.1. The set of $G$-triangulations based on boundary points of surfaces is not preserved under gluing of surfaces since gluing gives a new internal vertex. Hence, we had to allow more general $G$-triangulations from the start.

### 5.2. Cylinder Category

Since we want to construct the $(2,1)$-part of a functor

$$
\begin{equation*}
Z: \operatorname{G\mathcal {B}}^{\operatorname{sord}}{ }_{\star}(3,2,1) \rightarrow \operatorname{BiMod}_{\mathbb{K}} \tag{38}
\end{equation*}
$$

a surface $\Sigma$ with $n$-boundary components should get mapped to a profunctor

$$
\begin{equation*}
Z_{\Sigma}: \underbrace{\operatorname{Cyl}(\mathcal{C})^{\epsilon_{1}} \boxtimes \cdots \boxtimes \operatorname{Cyl}(\mathcal{C})^{\epsilon_{n}}}_{n} \rightarrow \operatorname{Vect}_{\mathbb{K}} \tag{39}
\end{equation*}
$$

with $\operatorname{Cyl}(\mathcal{C}):=Z\left(S^{1}\right)$ or its dual, depending on the orientation of the boundary component. For example, a pair of pants with two incoming and one outgoing boundary endows $\operatorname{Cyl}\left(S^{1}\right)$ with a monoidal structure. For a discussion of how a once extended $G$-equivariant HTQFT gives $\operatorname{Cyl}\left(S^{1}\right)$ the structure of a $G$ braided category, see [26].

Thus, we need to get the value of the string-net functor on a circle with additional $G$-label $g$ (cf. Sect. 4). That is, we need to compute the cylinder category. Since in the non-equivariant case of the string-net construction, the circle category is equivalent to the Drinfeld center of the input category, it is reasonable to expect that equivariant string-nets on a $g$-labeled circle compute the $g$-homogeneous component $Z_{G}(\mathcal{C})_{g}$ of the $G$-center $Z_{G}(\mathcal{C})$.

Remark 5.17. In [32], the cylinder category is chosen to be $\mathrm{Z}_{G}(\mathcal{C})$ by choosing a $Z_{G}(\mathcal{C})$-coloring on surfaces with boundary. It is a nice feature of the stringnet construction that one does not need to make that choice, but one can actually compute the category. The only input needed is a $G$-graded spherical fusion category $\mathcal{C}$, which is the supposed to be the value of a fully extended HTQFT on a point.

At the same time, the following computation is a check that our construction is really sensible.

Let $S_{g}^{1}$ be the $g$-labeled circle. Recall that $S_{g}^{1}$ has a distinguished base point $\star \in S_{g}^{1}$. We define a category $\widetilde{\mathcal{C y L}_{g}}$ as follows. Objects are sets of $\mathcal{C}$-labeled, signed marked points $(\mathbf{m}, \mathbf{c}):=\left\{\left(m_{1}^{\epsilon_{1}}, c_{1}^{\epsilon_{1}}\right), \ldots,\left(m_{N}^{\epsilon_{N}}, c_{N}^{\epsilon_{N}}\right)\right\}$. The points are linearly ordered by starting at $\star$ and going around the circle in counterclockwise direction. The object $c_{i}^{\epsilon_{i}}$ has underlying $G$-color $g_{i}^{\epsilon_{i}}$ and we require

$$
\begin{equation*}
g_{1}^{\epsilon_{1}} \cdots g_{N}^{\epsilon_{N}}=g \tag{40}
\end{equation*}
$$

In the non-equivariant string-net construction, the morphism space of the cylinder category is defined to be the string-net space, with the appropriate boundary value, on the cylinder. Similarly, in our equivariant setting, $\operatorname{Hom}_{\widetilde{C y l}_{g}}((\mathbf{m}, \mathbf{c}),(\mathbf{n}, \mathbf{d}))$ is given is by the string-net spaces, where we have to additionally specify a map to the classifying space $B G$. On the circle, this map is given by $g$. It is appropriate for cylinder categories to require it to be given by the neutral element $e \in G$ for the "radial" direction, cf. Figure 12. Composition of morphisms is given by gluing cylinders and concatenating string-nets.

Proposition 5.18. Let $\mathcal{C y l}_{g}$ be the Karoubi envelope of $\widetilde{\mathcal{C y}}{ }_{g}$. There is an equivalence

$$
\begin{equation*}
\mathcal{C y l}{ }_{g} \simeq \mathrm{Z}_{G}(\mathcal{C})_{g} \tag{41}
\end{equation*}
$$



Figure 12. In red, we show an ideal triangulation of the cylinder with in- and outgoing boundary $S_{g}^{1}$. The vertical red lines are along the "radial direction" of the cylinder

Proof. An object in $\mathrm{Cyl}_{g}$ is a pair $((\mathbf{m}, \mathbf{c}), p)$ of an object in $\widetilde{\mathcal{C y l}}{ }_{g}$ and an idempotent $p$ in $\widetilde{\mathcal{C y L}_{g}}$. We pick a $G$-triangulation of the cylinder without internal marked points. It is not hard to see that $p$ can be represented by a graph shown in red in Fig. 12 with $c, d \in \mathcal{C}_{g}$ the signed tensor products of objects $\mathbf{c}$, $\mathbf{d}$ and $x \in \mathcal{C}_{e}$.
We define a functor $F: \mathcal{C y l}_{g} \rightarrow \mathrm{Z}_{G}(\mathcal{C})_{g}$, whose action on objects is given by

$$
\begin{equation*}
F:((\mathbf{m}, \mathbf{c}), p) \longmapsto(I(c), F(p)) \tag{42}
\end{equation*}
$$

where $I$ is the induction functor (cf. (11)) and $F(p)$ is the idempotent shown in the following figure:

where $f$ is the morphism in the representation of $p$ given above. Its action on morphisms is defined in the obvious way. In the other direction, we define a functor $K: \mathrm{Z}_{G}(\mathcal{C})_{g} \longrightarrow \mathcal{C y \mathcal { C } _ { g }}$ mapping $Z \in \mathrm{Z}_{G}(\mathcal{C})_{g}$ to the underlying object in $\mathcal{C}_{g}$ and the idempotent given by, i.e., the crossing is the half-braiding in $\mathrm{Z}_{G}(\mathcal{C})_{g}$. The proof that the functors $F, K$ give an equivalence of categories is equivalent to the one given in [21, Theorem 6.4].


Figure 13. Part of the functor $K$ from $G$-center to the cylinder category

## 6. G-Equivariant String-Nets

### 6.1. Construction of $G$-Equivariant String-Net Space

In this section, we demonstrate how to extract structure of the category $\mathrm{Z}_{G}(\mathcal{C})$ from $G$-equivariant string-nets. To this end, we have to be able to assign to each connected component of the boundary of $\Sigma$ an object $Z \in Z_{G}(\mathcal{C})_{g}$. Given $Z$, consider the $G$-equivariant string-net given by the cylinder

where the purple line is a cloaking circle. This is exactly the idempotent of Fig. 13. Gluing such a cylinder to a boundary component defines an idempotent on the corresponding bare string-net space.

Given a boundary value $Z_{i} \in \mathrm{Z}_{G}(\mathcal{C})$ for each connected component of the boundary, we define the string-net space $\operatorname{KSN}^{\mathcal{C}}(\Sigma ; \mathbf{Z}):=\operatorname{KSN}^{\mathcal{C}}\left(\Sigma ; Z_{1}, \ldots, Z_{n}\right)$ for these boundary values as the common image of all these idempotents.

The cylinder category together with the string-net space should be the $(2,1)$-part of functor

$$
\begin{equation*}
\operatorname{KSN}^{\mathcal{C}}: G \mathcal{B o r d} d_{\star}(3,2,1) \rightarrow \mathcal{B i M}^{\operatorname{Mod}} \mathbb{K}_{\mathbb{K}} \tag{43}
\end{equation*}
$$

Thus, we have to discuss how string-net spaces behave under gluing of surfaces. Let $\Sigma, \Sigma^{\prime}$ be composable 1-morphisms in $G \mathcal{B}$ ord ${ }_{\star}(3,2,1)$, then we need to show

$$
\begin{equation*}
\operatorname{KSN}^{\mathcal{C}}\left(\Sigma \circ \Sigma^{\prime} ; \mathbf{Z}, \mathbf{Z}^{\prime}\right) \simeq \int^{Z \in \mathrm{Z}_{G}(\mathcal{C})} \operatorname{KSN}^{\mathcal{C}}(\Sigma ; \mathbf{Z}, Z) \otimes \operatorname{KSN}^{\mathcal{C}}\left(\Sigma^{\prime} ; Z^{*}, \mathbf{Z}^{\prime}\right) \tag{44}
\end{equation*}
$$

where $Z, Z^{*} \in \mathrm{Z}_{G}(\mathcal{C})$ are the boundary values of the boundary components at which $\Sigma, \Sigma^{\prime}$ are glued. Let $\Sigma, \Sigma^{\prime}$ be equipped with $G$-triangulations T, $\mathrm{T}^{\prime}$. The glued surface $\Sigma \circ \Sigma^{\prime}$ inherits a $G$-triangulation, where the boundary vertices of the individual surfaces are mapped to new internal vertices. For $\Sigma$, we pick a $G$-triangulation which looks around the gluing boundary like the $G$-triangulation shown in Fig. 13. Using the chosen $G$-triangulation on an annular neighborhood around the gluing boundary, the proof of (44) is the same as in the non-equivariant case, c.f. [17, section 5.2].

The alert reader may have noticed that we have not defined a string-net space on surfaces without boundary components. However, we can consider $\left(\Sigma^{\prime}, \zeta^{\prime}\right)$, a surface with a single boundary component $b$ and $\left.\zeta\right|_{b}=e \in G$ the neutral element. Furthermore, $(\Sigma, \zeta)$ is the surface obtained from $\left(\Sigma^{\prime}, \zeta^{\prime}\right)$ by gluing a disk $D$ at $b$. Note that all closed $G$-surfaces can be obtained in this way and we simply define

$$
\begin{equation*}
\operatorname{KSN}^{\mathcal{C}}(\Sigma):=\int^{Z \in \mathrm{Z}_{G}(\mathcal{C})_{e}} \operatorname{KSN}^{\mathcal{C}}(D ; Z) \otimes \operatorname{KSN}^{\mathcal{C}}\left(\Sigma ; Z^{*}\right) \tag{45}
\end{equation*}
$$

We expect that, as in the non-equivariant case, the maps

$$
\begin{align*}
\text { Cyl }:\left(S^{1}, g\right) & \rightarrow \mathcal{C y l}_{g} \\
\operatorname{KSN}^{\mathcal{C}}:(\Sigma, \zeta) & \rightarrow \operatorname{KSN}^{\mathcal{C}}(\Sigma) \tag{46}
\end{align*}
$$

comprise the $(2,1)$-part of a symmetric monoidal functor
$\mathrm{KSN}^{\mathcal{C}}: G \mathcal{B o r d}_{\star}(3,2,1) \rightarrow \mathcal{B i M}^{(1)}{ }_{\mathbb{K}}$.

### 6.2. Computations of the $G$-String-Net Space

In this section, we compute the string-net space for different surfaces. We present the computation for a cylinder, a pair of pants and a genus two surface with three boundary components. The cylinder computation shows how to recover the $G$-crossing on $\mathrm{Z}_{G}(\mathcal{C})$ and the pair of pants gives the monoidal product. The higher genus computation connects our construction to the $G$ equivariant Turaev-Viro construction of [32]. However, as we are using a slightly different input datum, we cannot get a direct isomorphism between vector spaces for $G$-surfaces. Nevertheless, we get an expression for the stringnet space in terms of morphism space in the $G$-center, which is closely related to the result in [32]. Finally we note that computing the string-net space for all other surfaces is analogous to the procedures presented here.
6.2.1. Cylinder. We start by computing the string-net space on a cylinder. Since a cylinder exhibits a homotopy between the $G$-bundles on its boundary, the boundary circles have to carry conjugate $G$-labels. We pick an ideal triangulation and note that any string-net is a sum of string-nets of the following type where $j \in I_{h}(\mathcal{C})$ for any $h \in G$.


Figure 14. Basic string-net on a cylinder

We denote the cylinder with lower/upper boundary labeled by $g / h^{-1} g h$ by $g_{C h^{-1} g h}$. From now on, we will not show the ideal triangulation in the arguments. Pictures get overloaded fast and besides establishing a useful presentation for string-nets, an ideal triangulation does not serve any further purpose.

Proposition 6.1. For $X_{1} \in \mathrm{Z}_{G}(\mathcal{C})_{g}$ and $X_{2} \in \mathrm{Z}_{G}(\mathcal{C})_{h^{-1} g h}$, we have

$$
\begin{equation*}
\operatorname{KSN}^{\mathcal{C}}\left(g_{C h^{-1} g h} ; X_{1}, X_{2}\right) \simeq \operatorname{Hom}_{\mathrm{Z}(\mathcal{C})}\left(\phi_{h}\left(X_{1}\right), X_{2}\right) \tag{47}
\end{equation*}
$$

where $\phi_{h}$ is the $G$-crossing on $\mathrm{Z}_{G}(\mathcal{C})$ from Proposition 2.3. In this way, $G$ equivariant string-nets encode the $G$-action on $\mathrm{Z}_{G}(\mathcal{C})$.

Proof. We can use the lower cloaking circle in Fig. 14 and the completeness relation to bring the $j$-labeled line and the other cloaking circle to the front


This also changes the color of the lower cloaking circle. The lower three consecutive undercrossings in (48) constitute half-braidings $\gamma_{x_{1}, i \otimes n_{0} \otimes j^{*}}$, where $n_{0} \in I_{e}(\mathcal{C})$ is a simple object in the neutral component of $\mathcal{C}$ coming from the cloaking circle in the middle.

We define linear maps

$$
F: \operatorname{KSN}^{\mathcal{C}}\left({ }_{g} C_{h^{-1} g h} ; X_{1}, X_{2}\right) \quad \longrightarrow \quad \operatorname{Hom}_{\mathrm{Z}(\mathcal{C})}\left(\phi_{h}\left(X_{1}\right), X_{2}\right)
$$



$$
G: \operatorname{Hom}_{\mathrm{Z}(\mathcal{C})}\left(\phi_{h}\left(X_{1}\right), X_{2}\right) \quad \longrightarrow \quad \mathrm{KSN}^{\mathcal{C}}\left({ }_{g} C_{h^{-1} g h} ; X_{1}, X_{2}\right)
$$



We claim that the two maps are inverse to each other. Applying $G \circ F$ to the string-net (48) yields


In (1), we use the definition of the idempotent $\pi\left(X_{2}\right)$ and then the completeness relation to get rid of the cloaking circle around $f$. Step (2) is again using the completeness relation, this time with the $I_{h}$-colored cloaking circle.

For the other composition, let $f^{\prime} \in \operatorname{Hom}_{\mathcal{Z}(\mathcal{C})}\left(\phi_{h}\left(X_{1}\right), X_{2}\right)$, then


Equality (1) is the definition of the half-braiding $\gamma_{\phi_{h}\left(X_{1}\right), \bullet}$ (2) holds, because $f^{\prime}$ is already a morphism in the $G$-center.
6.2.2. Pair of Pants. Similar to the cylinder, a pair of pants should equip the cylinder category with additional structure which in this case is the functor underlying a monoidal product. As a second consistency check, we prove that the equivariant string-net construction really induces the monoidal product in $\mathrm{Z}_{G}(\mathcal{C})$. We use the symbol $P_{g, h}^{g h}$ for the following surface with the shown boundary holonomies


Proposition 6.2. For $X_{1} \in \mathrm{Z}_{G}(\mathcal{C})_{g}, X_{2} \in \mathrm{Z}_{G}(\mathcal{C})_{h}$ and $X_{3} \in \mathrm{Z}_{G}(\mathcal{C})_{g h}$. Then it holds

$$
\begin{equation*}
\operatorname{KSN}^{\mathcal{C}}\left(P_{g, h}^{h} ; X_{1}, X_{2}, X_{3}\right) \simeq \operatorname{Hom}_{Z_{G}(\mathcal{C})}\left(X_{1} \otimes X_{2}, X_{3}\right) \tag{49}
\end{equation*}
$$

In this way, $G$-equivariant string-nets encode the monoidal structure of $\mathrm{Z}_{G}(\mathcal{C})$.

Proof. The proof is very similar as the previous one. We start with an ideal $G$-triangulation on a pair of pants


In black, we also included the dual uni-trivalent fat graph. We pick a maximal subtree in the fat graph and collapse the corresponding string-net to get


Note that due to the $G$-coloring in the ideal triangulation it holds $i, j \in I_{e}$. Using the completeness relation in $\mathcal{C}$, we get a string-net

which has an expansion in terms of string-nets


From that presentation of the string-net, the isomorphism is obvious.
6.2.3. Higher Genus Surfaces. As an example, we consider $\Sigma$ a surface of genus 2 with three boundary components. We pick a $G$-triangulation for $\Sigma$ with no internal vertices.


In the figure, the $G$-triangulation is shown in red. Its bounding octagon gives a polygonal decomposition of a genus 2 surface, and the edges of the polygon are labeled with group elements $\alpha, \beta, \gamma$ and $\delta$. Boundary edges have $G$-label $b_{1}, b_{2}$ and $b_{3}$. In black, we show a string-net on the surface, where we already collapsed all planar graphs inside faces of the $G$-triangulation to one-vertex graphs with labels $\phi_{i}$. The edges labeled by $\alpha, \beta, \gamma, \delta$ and $b_{1}, b_{2}, b_{3}$ generate the first homology group of $\Sigma$. By Lemma 5.11, we have a basis for $\operatorname{KSN}^{\mathcal{C}}(\Sigma ; \mathbf{B})$ spanned by equivalence classes of above graphs, with edges labeled by compatible simple objects.

We pick an arbitrary but fixed maximal tree $T$ in the above fat graph underlying the string-net. Such a tree has to lie inside a disk and we can collapse all its edges using the local relations. This yields a basis of $\operatorname{KSN}^{\mathcal{C}}(\Sigma ; \mathbf{B})$ in terms of (equivalence classes of) graphs of the form

where the $i_{n}$ 's are simple objects. We will not show the $G$-triangulation anymore from now on. We can isotope the cloaking circle around the first boundary component to run parallel to the fat graph. The resulting string-net, drawn on $\Sigma$, then looks as follows


We can use the completeness relation several times, to bring the fat graph in the form

$$
\sum_{\substack{i \in I_{e} \\ m_{1}, \ldots, m_{6}}} \frac{d_{i}}{D} d_{m_{1}} d_{m_{2}} d_{m_{3}} d_{m_{4}} d_{m_{5}} d_{m_{6}}
$$



Note that $\zeta$ fixes the $G$-color of a $G$-triangulation with no internal marked points uniquely. By collapsing a maximal tree, we get a fixed $G$-color of the
edges in the one-vertex graph. Since the cyclicity condition on vertices is preserved by the local moves, this coloring satisfies

$$
\begin{equation*}
h b_{3} h^{-1} b_{1} g b_{2} g^{-1} \gamma \delta \gamma^{-1} \delta^{-1} \alpha \beta \alpha^{-1} \beta^{-1}=e \tag{50}
\end{equation*}
$$

where $h$ is the $G$-color of $i_{8}$ and $g$ the one of $i_{5}$. By contracting the maximal tree $T$ of the fat graph, the non-contracted edges constitute generators for $\pi_{1}(\Sigma)$ based at the single vertex. Choosing another maximal tree, we get a different set of free generators. However, since the $G$-color of the fat graph was uniquely determined, the $G$-color for the generators is also uniquely determined in both cases. It solely depends on the map $\zeta$.

For the tree chosen here, we define an object

$$
\begin{equation*}
c_{X_{3}, X_{1}, X_{2}, \mathbf{m}}:=\phi_{h}\left(X_{3}\right) \otimes X_{1} \otimes \phi_{g}\left(X_{2}\right) \otimes m_{5} \otimes m_{4} \otimes m_{5}^{*} \otimes m_{4}^{*} \otimes m_{3} \otimes m_{2} \otimes m_{3}^{*} \otimes m_{2}^{*}, \tag{51}
\end{equation*}
$$

where $m_{5} \in I_{\beta}, m_{4} \in I_{\alpha}, m_{3} \in I_{\delta}, m_{2} \in I_{\gamma}$ and $X_{i} \in \mathrm{Z}_{G}(\mathcal{C})_{b_{i}}$ are the boundary values. The object

$$
\begin{equation*}
c:=\bigoplus_{\mathbf{m}} c_{X_{3}, X_{1}, X_{2}, \mathbf{m}} \tag{52}
\end{equation*}
$$

has a natural half-braiding

$$
\sum_{\substack{m_{2}, m_{3}, m_{4}, m_{5}}} d_{m_{2}} d_{m_{3}} d_{m_{4}} d_{m_{5}}
$$


and thus can be seen as an object $(c, \sigma) \in \mathrm{Z}_{G}(\mathcal{C})([32$, section 10.4]).
The computation is now similar to the one for the cylinder. We define a linear map

$$
\begin{gathered}
\operatorname{Hom}_{Z_{G}(\mathcal{C})}(\mathbb{1}, c) \longrightarrow \operatorname{KSN}^{\mathcal{C}}\left(\Sigma ; X_{1}, X_{2}, X_{3}\right) \\
\text { which acts on } \Upsilon \in \operatorname{Hom}_{Z_{G}(\mathcal{C})}(\mathbb{1}, c) \text { by }
\end{gathered}
$$


and a linear map in the other direction

$$
\operatorname{KSN}^{\mathcal{C}}\left(\Sigma ; X_{1}, X_{2}, X_{3}\right) \longrightarrow \operatorname{Hom}_{Z_{G}(\mathcal{C})}(\mathbb{1}, c)
$$



Using the same arguments as before, it is now easy to show that the two maps are inverse to each other.

The above computation can be summarized in the following proposition.

Proposition 6.3. For $\Sigma$ a genus two surface with three boundary components and $c \in \mathrm{Z}_{G}(\mathcal{C})$ as in (52), it holds

$$
\begin{equation*}
\operatorname{KSN}^{\mathcal{C}}\left(\Sigma ; X_{1}, X_{2}, X_{3}\right) \simeq \operatorname{Hom}_{\mathrm{Z}_{G}(\mathcal{C})}(\mathbb{1}, c) \tag{53}
\end{equation*}
$$

Remark 6.4. The string-net space $\mathrm{KSN}^{\mathcal{C}}\left(\Sigma ; X_{1}, X_{2}, X_{3}\right)$ has the same structure as the vector spaces obtained in a different construction in [32]. The different geometric input data in the two constructions make a detailed comparison of the vector spaces uninstructive.

## Acknowledgements

The authors thank Yang Yang and Theo Johnson-Freyd for useful discussions. CS and MT are supported by the Deutsche Forschungsgemeinschaft (DFG, German Research Foundation) under SCHW1162/6-1; CS is also supported by

DFG under Germany's Excellence Strategy - EXC 2121 "Quantum Universe" - 390833306.

Funding Open Access funding enabled and organized by Projekt DEAL.

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## Appendix Proof of Theorem 4.9

The theorem will follow from [6, Proposition 6.2], which we quickly state for the reader's convenience.

Theorem A.1. Given two two-dimensional CW-complexes $F, B$ and a map $\pi: F^{[1]} \rightarrow B^{[1]}$ of their 1-skeletons, such that: (i) For any vertex $v$ and edge $e$ in $B$, there exists a vertex $v^{\prime}$ and an edge $e^{\prime}$ with $\pi\left(v^{\prime}\right)=v$ and $\pi\left(e^{\prime}\right)=e$. Furthermore, for $v_{1} \xrightarrow{e} v_{2}$ an edge in $B$, and any $v_{1}^{\prime} \in \pi^{-1}\left(v_{1}\right)$, there exists an edge $v_{1}^{\prime} \xrightarrow{e^{\prime}} v_{2}^{\prime}$ such that $\pi\left(e^{\prime}\right)=e$.
(ii) $B$ is connected and simply connected.
(iii) There exist a vertex $v \in B$, such that $\pi^{-1}(v)$ is connected and simply connected.
(iv) Given an edge $v_{1} \xrightarrow{e} v_{2}$ in $B$, an edge $v_{1}^{\prime} \xrightarrow{f_{1}} v_{1}^{\prime \prime}$ in $\pi^{-1}\left(v_{1}\right)$ and two lifts $v_{1}^{\prime} \xrightarrow{e^{\prime}} v_{2}^{\prime}, v_{1}^{\prime \prime} \xrightarrow{e^{\prime \prime}} v_{2}^{\prime \prime}$ of $e$, there exists an edge $v_{2}^{\prime} \xrightarrow{f_{2}} v_{2}^{\prime \prime}$ in $\pi^{-1}\left(v_{2}\right)$ such that the square

is contractible in $F$.
(v) The boundary $\partial C$ of any 2-cell in $B$ has a contractible lift in $B$.

Then the complex $F$ is connected and simply connected.
Proof of Theorem 4.9. The proof consists of checking points (i)-(v) of Theorem A. 1 for the forgetful map $\pi: \mathcal{P}^{G}(\Sigma, \zeta, M)^{[1]} \rightarrow \mathcal{P}(\Sigma, M)^{[1]}$, which forgets
the $G$-labels of the edges, maps $G$-flips to flips of the underlying ideal triangulation and maps gauge transformations to the identity. Note that the fiber of $\pi$ over any ideal triangulation T is the set of all possible $G$-labels of $\Delta$ induced by $\zeta: \Sigma \rightarrow B G$.
(i) The forgetful map is clearly surjective on vertices and edges. In addition, given a flip $\Delta \xrightarrow{F_{e}} \Delta^{\prime}$ and any $G$-triangulation $(\Delta, g) \in \pi^{-1}(\Delta)$, there is a $G$-flip $(\Delta, g) \xrightarrow{F_{e}^{G}}\left(\Delta^{\prime}, g^{\prime}\right)$ covering $e$.
(ii) The Ptolemy complex is connected and simply connected by theorem 4.4.
(iii) Let $\Delta$ be any ideal triangulation of $\Sigma$. The fiber $\pi^{-1}(\Delta)$ is connected, as the gauge group $\mathcal{G}_{\Delta}$ relates any two $G$-labelings induced by $\zeta$. Due to relation GP5, any sequence of gauge transformations in $\pi^{-1}(\Delta)$ is homotopic to a sequence of the form $\lambda_{1}(v) \cdots \lambda_{n_{v}}(v) \lambda_{1}(w) \cdots \lambda_{n_{w}}(w) \lambda_{1}(u) \cdots$ running trough all vertices of $\Delta$ in arbitrary order. Furthermore, any two paths having the same start and endpoint in $\pi^{-1}(\Delta)$ and consisting entirely of different gauge transformations at the same vertex are homotopic. This follows from the observation that due to GP6 the subcomplex spanned by gauge transformations at a single vertex of a $G$-triangulation is the 2 -skeleton of a $|G|-1$-dimensional simplex and therefore is simply connected. It follows that $\pi^{-1}(\Delta)$ is simply connected.
(iv) This directly follows from the mixed relation GP4.
(v) Similar to the previous point, this holds, since the boundaries of the 2cells GP1, GP2 and GP3 project to the boundaries of the 2-cells P1, P2 and P3.

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Communicated by Denis Bernard.
Received: October 7, 2022.
Accepted: April 2, 2023.


[^0]:    ${ }^{1}$ See [31, section 4.1] for the definition of a non-singular category.

