

ERRATA

MAXIMUM WEIGHT INDEPENDENT SET IN TREES

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An algorithm of Frank [1] when applied to trees requires $O(n)$ time. Although this is better by a factor of $\log n$, this algorithm is highly sequential. The algorithm in [2] is a divide and conquer and can be parallelized. Also, the notion of alternating trees will be important in doing so.

1. A. Frank, *Some polynomial algorithms for certain graphs and hypergraphs*, Proc. 5th British Combinatorics Conference, 1975, pp. 211–226.
2. S. Pawagi, *Maximum weight independent set in trees*, BIT 27, 1987, 170–180.

ON THE UNIFORM CONVERGENCE OF A COLLOCATION METHOD FOR A CLASS OF SINGULAR INTEGRAL EQUATIONS

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The proof of theorem 1 on page 196 in my paper published in BIT 27, 1987, 190–202, only holds for $x = 1$. For $x = -1$ it is not valid, and for $x = 0$ it holds (with minor modifications) if the following additional constraints are introduced in the statement of that theorem:

$$\Gamma_2 \equiv \frac{1}{2} \min\{p + \mu, q + \sigma\} > \max\{0, \alpha, \beta\} \equiv \Gamma_1$$

with $r = \min\{p + \mu, q + \sigma\} - 2\xi$, $\xi \in (\Gamma_1, \Gamma_2)$.