

Disjoint Edges in Topological Graphs and the Tangled-Thrackle Conjecture^{*}

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Abstract. It is shown that for a constant $t \in \mathbb{N}$, every simple topological graph on n vertices has $O(n)$ edges if the graph has no two sets of t edges such that every edge in one set is disjoint from all edges of the other set (i.e., the complement of the intersection graph of the edges is $K_{t,t}$ -free). As an application, we settle the *tangled-thrackle* conjecture formulated by Pach, Radoičić, and Tóth: Every n -vertex graph drawn in the plane such that every pair of edges have precisely one point in common, where this point is either a common endpoint, a crossing, or a point of tangency, has at most $O(n)$ edges.

1 Introduction

A *topological graph* is a graph drawn in the plane such that its vertices are represented by distinct points and its edges are represented by Jordan arcs between the corresponding points satisfying the following (nondegeneracy) conditions: (a) no edge intersects any vertex other than its endpoints, (b) any two edges have only a finite number of interior points in common, (c) no three edges have a common interior point, and (d) if two edges share an interior point, then they properly cross at that point [7]. A topological graph is *simple* if every pair of edges intersect in at most one point. Two edges of a topological graph *cross* if their interiors share a point, and are *disjoint* if they neither share a common vertex nor cross.

In 2005, Pach and Tóth [10] conjectured that for every constant $t \geq 3$, an n -vertex simple topological graph has $O(n)$ edges if no t edges are pairwise disjoint. They gave an upper bound of $|E(G)| \leq O(n \log^{4t-8} n)$ for all such graphs. Despite much attention over the last 10 years (see related results in [3,11,15,16]), the conjecture is still open.

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The condition that *no t edges are pairwise disjoint* means that the intersection graph of the edges (Jordan arcs) contains no anti-clique of size t , or equivalently the complement of the intersection graph of the edges is K_t -free. In this paper, we consider a weaker condition that the complement of the intersection graph of the edges is $K_{t,t}$ -free, where $t \in \mathbb{N}$ is a constant. This means that graph G has no set of t edges that are all disjoint from another set of t edges. Since no such graph G contains $2t$ pairwise disjoint edges, [10] implies $|E(G)| \leq O(n \log^{8t-8} n)$. Our main result improves this upper bound to $O(n)$.

Theorem 1. *Let $t \in \mathbb{N}$ be a constant. The maximum number of edges in a simple topological graph with n vertices that does not contain t edges all disjoint from another set of t edges is $O(n)$.*

Application to thrackles. More than 50 years ago, Conway asked what is the maximum number of edges in an n -vertex *thrackle*, that is, a simple topological graph G in which every two edges intersect, either at a common endpoint or at a proper crossing [1]. He conjectured that every n -vertex thrackle has at most n edges. The first linear upper bound was obtained by Lovász, Pach, and Szegedy [4], who showed that all such graphs have at most $2n$ edges. This upper bound was successively improved, and the current record is $|E(G)| \leq \frac{167}{117}n < 1.43n$ due to Fulek and Pach [2].

As an application of Theorem 1, we prove the tangled-thrackle conjecture recently raised by Pach, Radoičić, and Tóth [9]. A drawing of a graph G is a *tangled-thrackle* if it satisfies conditions (a)-(c) of topological graphs and every pair of edges have precisely one point in common: either a common endpoint, or a proper crossing, or a point of tangency. Note that such a drawing is not a topological graph due to tangencies. Pach, Radoičić, and Tóth [9] showed that every n -vertex tangled-thrackle has at most $O(n \log^{12} n)$ edges, and described a construction with at least $\lfloor 7n/6 \rfloor$ edges. They conjectured that the upper bound can be improved to $O(n)$. Here, we settle this conjecture in the affirmative.

Theorem 2. *Every tangled-thrackle on n vertices has $O(n)$ edges.*

2 Disjoint Edges in Topological Graphs

In this section, we prove Theorem 1. We start with reviewing a few graph theoretic results used in our argument. The following is a classic result in extremal graph theory due to Kővári, Sós, and Turán.

Theorem 3 (see [8]). *Let $G = (V, E)$ be a graph that does not contain $K_{t,t}$ as a subgraph. Then $|E(G)| \leq c_1 |V(G)|^{2-1/t}$, where c_1 is an absolute constant.*

Two edges in a graph are called *independent* if they do not share an endpoint. We define the *odd-crossing number* $\text{odd-cr}(G)$ of a graph G to be the minimum number of unordered pairs of edges that are independent and cross an odd number of times over all

topological drawings of G . The *bisection width* of a graph G , denoted by $b(G)$, is the smallest nonnegative integer such that there is a partition of the vertex set $V = V_1 \cup V_2$ with $\frac{1}{3}|V| \leq |V_i| \leq \frac{2}{3}|V|$ for $i = 1, 2$, and $|E(V_1, V_2)| = b(G)$. The following result, due to Pach and Tóth, relates the odd-crossing number of a graph to its bisection width.¹

Theorem 4 ([10]). *There is an absolute constant c_2 such that if G is a graph with n vertices of vertex degrees d_1, \dots, d_n , then*

$$b(G) \leq c_2 \log n \sqrt{\text{odd-cr}(G) + \sum_{i=1}^n d_i^2}.$$

We also rely on the result due to Pach and Tóth [10] stated in the introduction.

Theorem 5 ([10]). *Let $G = (V, E)$ be an n -vertex simple topological graph, such that G does not contain t pairwise disjoint edges. Then $|E(G)| \leq c_3 n \log^{4t-8} n$, where c_3 is an absolute constant.*

From disjoint edges to odd crossings. Using a combination of Theorems 3–5, we establish the following lemma.

Lemma 1. *Let $G = (V, E)$ be simple topological bipartite graph on n vertices with vertex degrees d_1, \dots, d_n , such that G does not contain a set of t edges all disjoint from another set of t edges. Then*

$$b(G) \leq c_4 n^{1-\frac{1}{2t}} \log^{8t-3} n + c_4 \log n \sqrt{\sum_{i=1}^n d_i^2}, \tag{1}$$

where c_4 is an absolute constant.

Proof: Since G does not contain $2t$ pairwise disjoint edges, Theorem 5 yields

$$|E(G)| \leq c_3 n \log^{8t-8} n. \tag{2}$$

Let V_a and V_b be the vertex classes of the bipartite graph G . Consider a simple curve γ that decomposes the plane into two parts, containing all points in V_a and V_b , respectively. By applying a suitable homeomorphism to the plane that maps γ to a horizontal line, G is deformed into a topological graph G' such that (refer to Fig. 1)

1. The vertices in V_a are above the line $y = 1$, the vertices in V_b are below the line $y = 0$,

¹ Pach and Tóth [10] defined the odd-crossing number of a graph G to be the minimum number of pairs of edges that cross an odd number of times (over all drawings of G), including pairs of edges with a common endpoint. However, since the number of pairs of edges with a common endpoint is at most $\sum_{i=1}^n d_i^2$, this effects Theorem 4 only by a constant factor.

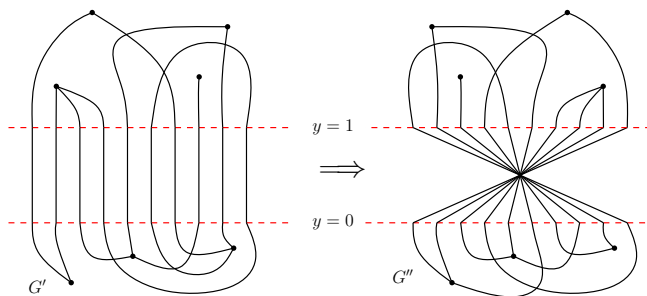


Fig. 1. Redrawing procedure

2. The part of any edge lying in the horizontal strip $0 \leq y \leq 1$ consists of vertical segments.

Since a homeomorphism neither creates nor removes intersections between edges, G and G' are isomorphic and their edges have the same intersection pattern.

Next we transform G' into a topological graph G'' by the following operations. Reflect the part of G' that lies above the $y = 1$ line about the y -axis. Replace the vertical segments in the horizontal strip $0 \leq y \leq 1$ by straight line segments that reconnect the corresponding pairs on the line $y = 0$ and $y = 1$, and perturb the segments if necessary to avoid triple intersections. Notice that if any two edges cross in G (and G'), then they must cross an even number of times in G'' . Indeed, suppose the edges e_1 and e_2 cross in G . Since G is simple, they share exactly one point in common. Let k_i denote the number of times edge e_i crosses the horizontal strip for $i \in 1, 2$, and note that k_i must be odd since the graph is bipartite. These $k_1 + k_2$ segments within the strip pairwise cross in G'' , creating $\binom{k_1+k_2}{2}$ crossings. Since edge e_i now crosses itself $\binom{k_i}{2}$ times in G'' , there are

$$\binom{k_1 + k_2}{2} - \binom{k_1}{2} - \binom{k_2}{2} = k_1 k_2 \tag{3}$$

crossings between edges e_1 and e_2 within the strip, which is odd when k_1 and k_2 are odd. Since e_1 and e_2 had one point in common outside the strip in both G and G'' , then e_1 and e_2 cross each other an even number of times in G'' . (Note that one can easily eliminate self-intersections by local modifications around these crossings.)

Hence, the number of pairs of edges that are independent and cross an odd number of times in G'' is at most the number of disjoint pairs of edges in G , which is in turn at most $c_1 |E(G)|^{2-1/t}$ by Theorem 3. Combined with (2), we have

$$\begin{aligned} \text{odd-cr}(G) &\leq c_1 (c_3 n \log^{8t-8} n)^{2-\frac{1}{t}} \\ &\leq cn^{2-\frac{1}{t}} \log^{16t-8} n, \end{aligned}$$

where c is an absolute constant. Together with Theorem 4, we have

$$\begin{aligned}
 b(G) &\leq c_2 \log n \sqrt{\left(cn^{2-\frac{1}{t}} \log^{16t-8} n \right) + \sum_{i=1}^n d_i^2} \\
 &\leq c_2 \sqrt{cn}^{1-\frac{1}{2t}} \log^{8t-3} n + c_2 \log n \sqrt{\sum_{i=1}^n d_i^2} \\
 &\leq c_4 n^{1-\frac{1}{2t}} \log^{8t-3} n + c_4 \log n \sqrt{\sum_{i=1}^n d_i^2},
 \end{aligned}$$

where c_4 is an absolute constant, as required. □

If the maximum degree of G is relatively small, we obtain a sublinear bound on the bisection width.

Corollary 1. *Let $G = (V, E)$ be simple topological bipartite graph on n vertices with vertex degrees $d_1, \dots, d_n \leq n^{1/5}$ such that G does not contain a set of t edges all disjoint from another set of t edges. Then*

$$b(G) \leq c_5 n^{1-\frac{1}{4t}}, \tag{4}$$

where c_5 is an absolute constant.

Proof: Substituting $d_i \leq n^{1/5}$ into (1), we have

$$\begin{aligned}
 b(G) &\leq c_4 n^{1-\frac{1}{2t}} \log^{8t-3} n + c_4 \log n \sqrt{n \cdot n^{2/5}} \\
 &\leq c_4 n^{1-\frac{1}{2t}} \log^{8t-3} n + c_4 n^{7/10} \log n \\
 &\leq c_5 n^{1-\frac{1}{4t}},
 \end{aligned} \tag{5}$$

for a sufficiently large constant c_5 . □

Vertex splitting for topological graphs. Given a simple topological graph with n vertices, we reduce the maximum degree below $n^{1/5}$ by a standard vertex splitting operation. Importantly, this operation can be performed such that it preserves the intersection pattern of the edges.

Lemma 2. *Let G be a simple topological graph with n vertices and m edges; and let $\Delta \geq 2m/n$. Then there is a simple topological graph G' with maximum degree at most Δ , at most $n + 2m/\Delta$ vertices, and precisely m edges such that the intersection graph of its edges is isomorphic to that of G .*

Proof: We successively split every vertex in G whose degree exceeds Δ as follows. Refer to Fig. 2. Let v be a vertex of degree $d(v) = d > \Delta$, and let vu_1, vu_2, \dots, vu_d be the edges incident to v in counterclockwise order. In a small neighborhood around v , replace v by $\lceil d/\Delta \rceil$ new vertices, $v_1, \dots, v_{\lceil d/\Delta \rceil}$ placed in counterclockwise order on a

circle of small radius centered at v . Without introducing any new crossings, connect u_j to v_i if and only if $\Delta(i - 1) < j \leq \Delta i$ for $j \in \{1, \dots, d\}$ and $i \in \{1, \dots, \Delta\}$. Finally, we do a local change within the small circle by moving each vertex v_i across the circle, so that every edge incident to v_i crosses all edges incident to $v_{i'}$, for all $i \neq i'$. As a result, any two edges incident to some vertex $\{v_1, \dots, v_{\lceil d/\Delta \rceil}\}$ intersect precisely once: either at a common endpoint or at a crossing within the small circle centered at v .

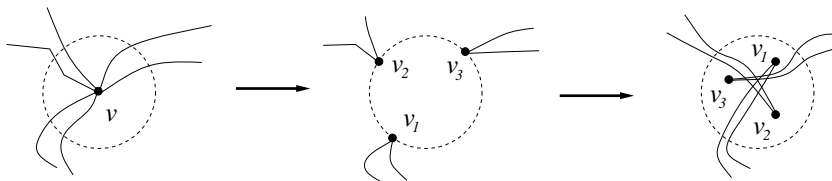


Fig. 2. Splitting a vertex v into new vertices v_1, v_2, v_3 , such that each v_i has degree at most Δ . Moreover, we do not introduce any disjoint pairs of edges and our new graph remains simple.

After applying this procedure to all vertices G , we obtain a simple topological graph G' of maximum degree at most Δ . By construction, G' has m edges, and the intersection pattern of the edges is the same as in G . The number of vertices in G' is

$$|V(G')| \leq \sum_{v \in V} \left\lceil \frac{d(v)}{\Delta} \right\rceil \leq n + \sum_{v \in V} \frac{d(v)}{\Delta} \leq n + \frac{2m}{\Delta},$$

as claimed. □

Putting things together. Since all graphs have a bipartite subgraph with at least half of its edges, Theorem 1 immediately follows from the following.

Theorem 6. *Let G be an n -vertex simple topological bipartite graph such that G does not contain t edges all disjoint from another set of t edges. Then*

$$|E(G)| \leq c_6(n - n^{1-\frac{1}{4t}}), \tag{6}$$

where $c_6 = c_6(t)$ depends only on t .

Proof: Let $t \in \mathbb{N}$ be a constant. We proceed by induction on n . Let $n_0 = n_0(t)$ be a sufficiently large constant (specified in (8) and (9) below) that depends only on t , and on the constants c_3 and c_5 defined in Theorem 5 and Corollary 1, respectively. Let c_6 be a sufficiently large constant such that $c_6 \geq 2c_5$ and for every positive integer $n \leq n_0$, we have

$$c_3 n \log^{8t-8} n \leq c_6(n - n^{1-\frac{1}{4t}}), \tag{7}$$

The choice of n_0 and c_6 ensures that (6) holds for all graphs with at most n_0 vertices. Now consider an integer $n > n_0$, and assume that (6) holds for all graphs with fewer than n vertices. Let G be a simple topological bipartite graph with n vertices such that G does not contain t edges all disjoint from another set of t edges.

By Theorem 5, G has $m \leq c_3 n \log^{8t-8} n$ edges. By Lemma 2, there is a simple topological graph G' of maximum degree at most $\Delta = n^{1/5}$, $n' \leq n + 2m/n^{1/5}$ vertices, and $m' = m$ edges, such that the intersection graph of its edges is isomorphic to that of G . Theorem 5 implies that $n' \leq n + 2m/n^{1/5} \leq n + 2c_3 n^{4/5} \log^{8t-8} n$. If $n \geq n_0$ for a sufficiently large constant n_0 , then

$$n' \leq n + 2c_3 n^{4/5} \log^{8t-8} n \leq n + n^{5/6}. \tag{8}$$

Since G and G' have the same number of edges, it is now enough to estimate $|E(G')|$. Note that $\Delta = n^{1/5} \leq (n')^{1/5}$, and by Corollary 1, the bisection width of G' is bounded by

$$b(G') \leq c_5 (n')^{1-\frac{1}{4t}} \leq c_5 (n + n^{5/6})^{1-\frac{1}{4t}} \leq 2c_5 n^{1-\frac{1}{4t}}.$$

Partition the vertex set of G' as $V' = V_1 \cup V_2$ with $\frac{1}{3}|V'| \leq V_i \leq \frac{2}{3}|V'|$ for $i = 1, 2$, such that G' has $b(G')$ edges between V_1 and V_2 . Denote by G_1 and G_2 the subgraphs induced by V_1 and V_2 , respectively. Put $n_1 = |V_1|$ and $n_2 = |V_2|$, where $n_1 + n_2 = n' \leq n + n^{5/6}$.

Note that both G_1 and G_2 are simple topological graphs that do not contain t edges all disjoint from another set of t edges. By the induction hypothesis, $|E(G_i)| \leq c_6 (n_i - n_i^{1-1/7t})$ for $i = 1, 2$. The total number of edges in G_1 and G_2 is

$$\begin{aligned} |E(G_1)| + |E(G_2)| &\leq c_6 (n_1 - n_1^{1-\frac{1}{7t}}) + c_6 (n_2 - n_2^{1-\frac{1}{7t}}) \\ &\leq c_6 (n_1 + n_2) - c_6 (n_1^{1-\frac{1}{7t}} + n_2^{1-\frac{1}{7t}}) \\ &\leq c_6 (n') - c_6 \left(\left(\frac{n_1}{n'} \right)^{1-\frac{1}{7t}} + \left(\frac{n_2}{n'} \right)^{1-\frac{1}{7t}} \right) (n')^{1-\frac{1}{7t}} \\ &\leq c_6 (n + n^{5/6}) - c_6 \left(\left(\frac{1}{3} \right)^{1-\frac{1}{7t}} + \left(\frac{2}{3} \right)^{1-\frac{1}{7t}} \right) n^{1-\frac{1}{7t}} \\ &= c_6 (n + n^{5/6}) - c_6 \alpha n^{1-\frac{1}{7t}} \\ &\leq c_6 (n - n^{1-\frac{1}{7t}}) + c_6 \left(n^{5/6} - (\alpha - 1)n^{1-\frac{1}{7t}} \right), \end{aligned}$$

where we write $\alpha = (1/3)^{1-1/7t} + (2/3)^{1-1/7t}$ for short. Note that for every $t \in \mathbb{N}$, we have $\alpha > 1$. Taking into account the edges between V_1 and V_2 , the total number of edges in G' (and hence G) is

$$\begin{aligned} |E(G')| &= |E(G_1)| + |E(G_2)| + b(G') \\ &\leq c_6 (n - n^{1-\frac{1}{7t}}) + c_6 \left(n^{5/6} - (\alpha - 1)n^{1-\frac{1}{7t}} \right) + 2c_5 n^{1-\frac{1}{4t}} \\ &\leq c_6 (n - n^{1-\frac{1}{7t}}) + c_6 \left(n^{5/6} + n^{1-\frac{1}{4t}} - (\alpha - 1)n^{1-\frac{1}{7t}} \right) \\ &\leq c_6 (n - n^{1-\frac{1}{7t}}), \end{aligned} \tag{9}$$

where the last inequality holds for $n \geq n_0$ if n_0 is sufficiently large (independent of c_6). This completes the induction step, hence the proof of Theorem 6. \square

3 Application: The Tangled-Thrackle Conjecture

Let G be tangled-thrackle with n vertices. By slightly modifying the edges (i.e., Jordan arcs) near the points of tangencies, we obtain a simple topological graph \tilde{G} with the same number of vertices and edges such that every pair of tangent edges in G become disjoint in \tilde{G} and all other intersection points between edges remain the same. In order to show that $|E(G)| \leq O(n)$, invoking Theorem 1, it suffices to prove the following.

Lemma 3. *For every tangled-thrackle G , the simple topological graph \tilde{G} does not contain a set of 200 edges all disjoint from another set of 200 edges.*

Before proving the lemma, we briefly review the concept of Deavenport-Schinzel sequences and arrangements of pseudo-segments.

A finite sequence $U = (u_1, \dots, u_t)$ of symbols over a finite alphabet is a *Davenport-Schinzel sequence of order s* if it satisfies the following two properties:

- no two consecutive symbols in the sequence are equal to each other;
- for any two distinct letters of the alphabet, a and b , the sequence does not contain a (not necessarily consecutive) subsequence (a, b, a, \dots, b, a) consisting of $s + 2$ symbols alternating between a and b .

The maximum length of a Davenport-Schinzel sequence of order s over an alphabet of size n is denoted $\lambda_s(n)$. Sharp asymptotic bounds for $\lambda_s(n)$ were obtained by Nivasch [6] and Pettie [13]. However, to avoid the constants hidden in the big-Oh notation, we use simpler explicit bounds. Specifically, we use the following upper bound for $\lambda_3(n)$.

Theorem 7 (see Proposition 7.1.1 in [5]). $\lambda_3(n) < 2n \ln n + 3n$.

A set \mathcal{L} of m Jordan arcs in the plane is called an *arrangement of pseudo-segment* if each pair of arcs intersects in at most one point (at an endpoint, a crossing, or a point of tangency), and no three arcs have a common interior point. An arrangement of pseudo-segments naturally defines a plane graph: The *vertices* of the arrangement are the endpoints and the intersection points of the Jordan arcs, and the *edges* are the portions of the Jordan arcs between consecutive vertices. The *faces* of the arrangement are the connected components of the complement of the union of the Jordan arcs. The vertices and edges are said to be *incident to a face* if they are contained in the (topological) closure of that face. The following theorem is a particular case of Theorem 5.3 of [14].

Theorem 8 (see [14]). *Let \mathcal{L} be an arrangement of m pseudo-segments and F be a face of \mathcal{L} . Then number of edges incident to F is at most $\lambda_3(2m)$.*

Lemma 4. *Let $\mathcal{L}_1 \cup \mathcal{L}_2$ be an arrangement of pseudo-segments such that every arc in \mathcal{L}_1 is tangent to all arcs in \mathcal{L}_2 ; and \mathcal{L}_1 and \mathcal{L}_2 each form a connected arrangement. Then \mathcal{L}_1 or \mathcal{L}_2 contains at most 200 arcs.*

Proof: Suppose to the contrary, that both \mathcal{L}_1 and \mathcal{L}_2 contain at least 200 arcs. Without loss of generality, we may assume $|\mathcal{L}_1| = |\mathcal{L}_2| = 200$. Since no arc in \mathcal{L}_1 crosses any

arc in \mathcal{L}_2 , the arrangement \mathcal{L}_1 lies in the closure of a single face F_2 of the arrangement \mathcal{L}_2 , and vice versa \mathcal{L}_2 lies in the closure of a single face F_1 of the arrangement \mathcal{L}_1 . We construct a plane graph H representing the tangencies between the edges of the two arrangements: place a vertex on the relative interior of each edge of \mathcal{L}_1 incident to F_1 and each edge of \mathcal{L}_2 incident to F_2 . Join two vertices, v and u , by an edge iff their corresponding edges of the arrangements, e_u and e_v , are tangent to each other. To see that H is indeed planar, note that each edge uv can be drawn closely following the arcs e_u and e_v to their intersection point in such a way that H has no crossings. As every arc in \mathcal{L}_1 is tangent to all arcs in \mathcal{L}_2 , the graph H has exactly 200^2 edges. By Theorem 8, H has at most $2\lambda_3(400)$ vertices. However,

$$\frac{|E(H)|}{|V(H)|} \geq \frac{200^2}{2\lambda_3(400)} > \frac{200^2}{4 \cdot 400 \ln 400 + 6 \cdot 400} > 3.3,$$

which contradicts Euler’s formula. □

Note that Lemma 4 easily generalizes to the case when the arrangements \mathcal{L}_1 and \mathcal{L}_2 are not necessarily connected but they have the property that every pseudo-segment in \mathcal{L}_1 lies in the same face of \mathcal{L}_2 , and vice versa.

It is now easy to see that Lemma 3 follows directly from Lemma 4: if 200 edges of the simple topological graph \tilde{G} are disjoint from another set of 200 edges of \tilde{G} , then a set of the corresponding 200 edges of the tangled-thrackle G are tangent to corresponding other set of 200 edges of G .

Proof of Theorem 2: The statement follows by combining Theorem 1 and Lemma 3. □

We now show an analogue of Lemma 4 where we drop the condition that the arrangements \mathcal{L}_1 and \mathcal{L}_2 are connected. We find this interesting for its own sake.

Proposition 1. *Let $\mathcal{L}_1 \cup \mathcal{L}_2$ be an arrangements of pseudo-segments such that every arc in \mathcal{L}_1 is tangent to all arcs in \mathcal{L}_2 . Then \mathcal{L}_1 or \mathcal{L}_2 contains at most 400 arcs.*

Proof: Suppose to the contrary, that both \mathcal{L}_1 and \mathcal{L}_2 contain at least 400 arcs. Without loss of generality, we may assume $|\mathcal{L}_1| = |\mathcal{L}_2| = 400$. The difference from Lemma 3 is that the arrangement \mathcal{L}_1 or \mathcal{L}_2 may not be connected, and so the arcs in \mathcal{L}_1 could be distributed in several faces of the arrangement \mathcal{L}_2 .

Consider an arc $\ell \in \mathcal{L}_1$, and denote by s the number of other arcs in \mathcal{L}_1 that intersect ℓ , where $0 \leq s \leq 399$. Then ℓ contains precisely $s + 1$ edges of the arrangement \mathcal{L}_1 . Partition ℓ into two Jordan arcs: $\ell_a \subset \ell$ consists of the first $\lfloor (s + 1)/2 \rfloor$ edges along ℓ , and $\ell_b = \ell \setminus \ell_a$. Recall that ℓ is tangent to all 400 arcs in \mathcal{L}_2 . By the pigeonhole principle, we may assume w.l.o.g. that ℓ_a is tangent to at least 200 arcs in \mathcal{L}_2 . Let $\mathcal{L}'_2 \subset \mathcal{L}_2$ be a set of 200 arcs in \mathcal{L}_2 that are tangent to ℓ_a . By construction, ℓ_a intersects at most $\lfloor s/2 \rfloor \leq 200$ arcs in \mathcal{L}_1 . Consequently, there is a set $\mathcal{L}'_1 \subset \mathcal{L}_1$ of 200 arcs in \mathcal{L}_1 that do not intersect ℓ_a . Observe that ℓ_a lies in a single face of the arrangement \mathcal{L}'_1 . Since every arc in \mathcal{L}'_2 intersects ℓ_a , all arcs in \mathcal{L}'_2 lie in the same face of the arrangement \mathcal{L}'_1 .

We have found subsets $\mathcal{L}'_1 \subset \mathcal{L}_1$ and $\mathcal{L}'_2 \subset \mathcal{L}_2$ of size $|\mathcal{L}'_1| = |\mathcal{L}'_2| = 200$ such that every arc in \mathcal{L}'_1 is tangent to all arcs in \mathcal{L}'_2 ; and all arcs of \mathcal{L}'_1 lie in the same face of \mathcal{L}'_2 and vice versa. This contradicts Lemma 4 and the remark following its proof. □

4 Concluding Remarks

1. We showed that for every integer t , the maximum number of edges in a simple topological graph with n vertices that does not contain t edges all disjoint from another set of t edges is cn , where $c = c(t)$. A careful analysis of the proof shows that $c = 2^{O(t \log t)}$. It would be interesting to see if one could improve the upper bound on c to $O(t)$.

2. We suspect that the bounds of 200 and 400 in Lemma 4 and Proposition 1 are not optimal. Since any constant bound yields a linear upper bound for the number of edges in tangled-thrackles, we have not optimized these values. However, finding the best possible constants, or shorter proofs for some arbitrary constant bounds, would be of interest.

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