

Tolerance and Control Solutions of Two-Sided Interval Linear System and Their Applications

Worrawate Leela-apiradee^{1(⊠)}, Phantipa Thipwiwatpotjana², and Artur Gorka^{3,4},

- Department of Mathematics and Statistics, Faculty of Science and Technology, Thammasat University, Pathum Thani 12121, Thailand worrawateleela@gmail.com
 - Department of Mathematics and Computer Science, Faculty of Science, Chulalongkorn University, Bangkok 10330, Thailand phantipa.t@chula.ac.th
 - Department of Mathematics, Erskine College, Due West, SC 29639, USA Kamnoetvidya Science Academy, Wangchan, Rayong 21210, Thailand art.gorka@erskine.edu

Abstract. This work investigates tolerance and control solutions to a two-sided interval linear system. Their semantics are different, even though, we would be able to interchange the role of the interval information algebraically. We present necessary and sufficient conditions of their solvabilities as the inequalities depending on center and radius of coefficient interval matrices on both sides of the system. In a situation when the vector of variables is nonnegative, the conditions can simply be modified as the inequalities depending on boundaries of the interval matrices. This result helps to find out the feasible solutions of a quadratic programming problem with two-sided interval linear equation constraints.

Keywords: Interval linear system \cdot Tolerance solution \cdot Control solution

1 Introduction

An interval linear system of equations is normally referred to as a system Ax = b, where A is an interval matrix and b is an interval vector, while x is a vector of variables. As the matrix and the right hand side vector information of the system is not precise, it is impossible to provide a solution x to the system without any appropriate meaning.

Some literature presented different types of solutions of the system Ax = b depending on the purposes of the solutions. These solution types include weak, strong, tolerance and control solutions whose names reflect well on their mathematical definitions. For example, "x is a weak solution to Ax = b" means

© Springer Nature Switzerland AG 2020 M.-J. Lesot et al. (Eds.): IPMU 2020, CCIS 1239, pp. 436–448, 2020. $\frac{1}{1000} \frac{1}{1000} \frac{1}{1000}$ that

$$\exists A \in \mathbf{A}, \exists b \in \mathbf{b} \text{ such that } Ax = b,$$

while "x is a strong solution to Ax = b" means that

$$\forall A \in \mathbf{A}, \forall b \in \mathbf{b} \text{ such that } Ax = b.$$

Beaumont presented in [1] an efficient method derived from the simplex algorithm to compute inner and outer inclusion of the weak (united) solution set. A full analysis of the solvability and the conditions for checking whether x is a particular solution type to the system $\mathbf{A}x = \mathbf{b}$ was also provided in the literature [2-4,8-12].

The background on a tolerance solution arose from the crane construction problem in [8] and the input-output planning problem with inexact data in [11]. The characteristic of a tolerance solution x is to make Ax stay in the boundary of b. Shary [12] first motivated an idea of a control solution, which is the opposite case of a tolerance solution. In addition, Tian et al. developed in [13] a tolerance-control solution for the case when each row index of the system Ax = b performs either tolerance or control. Recently, Leela-apiradee, [6], have provided its solution set in terms of level set.

Instead of the system $\mathbf{A}x = \mathbf{b}$, the goal of this paper is to deal with a two-sided interval linear system. "Two-sided" means that the right hand side interval vector \mathbf{b} is substituted by the term $\mathbf{B}y$. The paper then presents the tolerance and control solutions of a two-sided interval linear system together with their solvability conditions.

To lead to the main idea of the paper, let us first introduce some basic notation of an interval matrix and an interval vector that can be seen as a matrix and a vector of interval components as follows.

- An $m \times n$ interval matrix **A** is defined by

$$\boldsymbol{A} = \begin{pmatrix} [\underline{a}_{11}, \overline{a}_{11}] & [\underline{a}_{12}, \overline{a}_{12}] & \cdots & [\underline{a}_{1n}, \overline{a}_{1n}] \\ [\underline{a}_{21}, \overline{a}_{21}] & [\underline{a}_{22}, \overline{a}_{22}] & \cdots & [\underline{a}_{2n}, \overline{a}_{2n}] \\ \vdots & \vdots & \ddots & \vdots \\ [\underline{a}_{m1}, \overline{a}_{m1}] & [\underline{a}_{m2}, \overline{a}_{m2}] & \cdots & [\underline{a}_{mn}, \overline{a}_{mn}] \end{pmatrix},$$

where \underline{a}_{ij} and \overline{a}_{ij} are real numbers such that $\underline{a}_{ij} \leq \overline{a}_{ij}$ for each $i \in \{1, 2, ..., m\}$ and $j \in \{1, 2, ..., n\}$.

- The interpretation of the interval matrix \boldsymbol{A} can be written as the set of matrices, that is,

$$\mathbf{A} = \left[\underline{A}, \overline{A}\right] = \left\{A \in \mathbb{R}^{m \times n} : \underline{A} \le A \le \overline{A}\right\},\,$$

where

$$\underline{A} = \begin{pmatrix} \underline{a}_{11} & \underline{a}_{12} & \cdots & \underline{a}_{1n} \\ \underline{a}_{21} & \underline{a}_{22} & \cdots & \underline{a}_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ \underline{a}_{m1} & \underline{a}_{m2} & \cdots & \underline{a}_{mn} \end{pmatrix} \quad and \quad \overline{A} = \begin{pmatrix} \overline{a}_{11} & \overline{a}_{12} & \cdots & \overline{a}_{1n} \\ \overline{a}_{21} & \overline{a}_{22} & \cdots & \overline{a}_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ \overline{a}_{m1} & \overline{a}_{m2} & \cdots & \overline{a}_{mn} \end{pmatrix}.$$

The ordering ' \leq ' between two matrices \underline{A} and \overline{A} is referred to componentwise inequality, i.e.,

$$\underline{A} \leq \overline{A} \ \text{ if and only if } \ \underline{a}_{ij} \leq \overline{a}_{ij} \ \forall i \in \{1,2,\ldots,m\}, \forall j \in \{1,2,\ldots,n\}.$$

Therefore, we note here that $A \in \mathbf{A}$ means $\underline{A} \leq A \leq \overline{A}$.

Moreover, we can use the following notation to represent the interval matrix A.

$$\mathbf{A} = \left[\underline{A}, \overline{A} \right] = \left[A_c - \Delta_A, A_c + \Delta_A \right], \tag{1}$$

where $A_c = \frac{1}{2}(\overline{A} + \underline{A})$ and $\Delta_A = \frac{1}{2}(\overline{A} - \underline{A})$.

Fiedler et al. proved in [2] that the lower and upper bounds of the interval vector $\mathbf{A}x$, denoted by $\underline{\mathbf{A}x}$ and $\overline{\mathbf{A}x}$ for any $x = (x_1, x_2, \dots, x_n)^T \in \mathbb{R}^n$, can be viewed by

$$\mathbf{A}x = \left[\underline{\mathbf{A}x}, \overline{\mathbf{A}x}\right] = \left[A_c x - \Delta_A |x|, \ A_c x + \Delta_A |x|\right],\tag{2}$$

where |x| is defined as the absolute of vector x, i.e., $|x| = (|x_1|, |x_2|, \dots, |x_n|)^T$.

Given B be another interval matrix and y be another vector of variables. In the situation when we have the term By on the right hand side instead of the vector b, the system "Ax = b" would become "Ax = By", which is called a two-sided interval linear system. The dimensions of x and y do not need to be the same but the number of rows of interval matrices A and B does. The definitions of weak, strong, tolerance control and tolerance-control solutions of Ax = By would mathematically be defined in the same fashion as the case of Ax = b.

In this paper, we focus on the tolerance, control and tolerance-control solutions of the system Ax = By. In Sect. 2 we give the definitions and their characterizations by equivalent conditions. Usually, two sets of quantities are equal when the left and the right quantities are the same. However, there is often a situation with imprecise information that one set of quantities is being controlled by the other. This means that if the two sets of quantities are not precise with having interval information, then one set of quantities must be subset of the other. Moreover, one set of the interval information could be more important to the system than the other. The other set of the interval information must follow the semantics of the context. This leads to the interpretation of tolerance and control solutions of a two-sided interval linear system discussed in Sect. 3. A couple of application examples are demonstrated in Sect. 4, which can be modeled by a two-sided interval linear system and a quadratic programming problem with two-sided interval linear constraints. The conclusion is addressed in the last section.

2 Tolerance and Control Solutions of Two-Sided Interval System of Linear Equations

To see how the two-sided interval systems are motivated, consider two systems of standard linear equations below:

$$Ax = By, (3)$$

$$Ax - By = 0, (4)$$

where $A \in \mathbb{R}^{m \times n}$, $B \in \mathbb{R}^{m \times p}$, and m, n and p are positive integers. Any (x, y) in $\mathbb{R}^n \times \mathbb{R}^p$ satisfying System (3) (or (4)) is called a **solution** of the system. It is clear that Systems (3) and (4) are algebraically equivalent. The set of solutions of (3) is the same as the set of solutions of (4). However, if the entries of coefficient matrices A and B in (3) and (4) are interval data, Systems (3) and (4) turn into

$$\mathbf{A}x = \mathbf{B}y,\tag{5}$$

$$\mathbf{A}x - \mathbf{B}y = 0, (6)$$

respectively, where coefficient terms A and B in (5) and (6) are interval matrices as defined in the introduction section. System (6) could not be well-defined in general since the left-hand side may be represented as interval vector with nonzero width as a result of Moore's standard interval arithmetic [7], while the right-hand side is a real zero vector with zero width. Therefore, we would not be able to move By to the same side as Ax of the equality as usual, when we deal with interval data. However, System (5) called a **two-sided interval system** of linear equations is well-defined since both sides of the equation are interval vectors. A solution $(x,y) \in \mathbb{R}^n \times \mathbb{R}^p$ to System (5) is not as simple as the case of standard matrices A and B, but it comes with its semantics.

In [2], Fiedler et al. defined tolerance and control solutions to an interval linear system $\mathbf{A}x = \mathbf{b}$. Based on the concepts of these solutions, Tian et al. later proposed a tolerance-control solution in [13].

Definition 1 (see [2] and [13]). A vector $x \in \mathbb{R}^n$ is called

- 1. a tolerance solution of Ax = b if for each $A \in A$ there exists $b \in b$ such that Ax = b,
- 2. a control solution of Ax = b if for each $b \in b$ there exists $A \in A$ such that Ax = b,
- 3. a tolerance-control solution of Ax = b if each row index of the system is either tolerance or control.

As we expand an interval linear system Ax = b to a two-sided interval linear system Ax = By, the types of solutions of Ax = By presented in the following definition are developed in similar fashion as the solution concepts in Ax = b.

Definition 2. A vector $(x,y) \in \mathbb{R}^n \times \mathbb{R}^p$ is called

1. a tolerance solution of Ax = By if for each $A \in A$ there exists $B \in B$ such that Ax = By,

- 2. a control solution of Ax = By if for each $B \in B$ there exists $A \in A$ such that Ax = By,
- 3. a tolerance-control solution of Ax = By if each row index of the system is either tolerance or control.

Without semantics, the mathematical expression of tolerance and control solutions are the same. In Sect. 3, we will mention more about these two solutions and separately redefine them according to their semantics.

We now consecutively establish Theorems 1–3 based on Definition 2 to find necessary and sufficient conditions for checking tolerance and tolerance-control solvabilities of $\mathbf{A}x = \mathbf{B}y$. The conditions presented in Theorems 2 and 3 are in a form of inequalities depending on center and radius of the coefficient interval matrices \mathbf{A} and \mathbf{B} with absolute terms |x| and |y|.

Theorem 1. A vector (x, y) is a tolerance solution of $\mathbf{A}x = \mathbf{B}y$ if and only if it satisfies $\mathbf{A}x \subseteq \mathbf{B}y$.

Proof. Assume that (x, y) is a tolerance solution of $\mathbf{A}x = \mathbf{B}y$. Let $A \in \mathbf{A}$. Then, there exists $B \in \mathbf{B}$ such that Ax = By. Thus,

$$Ax = By \in \mathbf{B}y = \left[\mathbf{B}y, \overline{\mathbf{B}y}\right],$$

that is,

$$By \le Ax \le \overline{By}$$
 for any $A \in A$,

which concludes $\mathbf{A}x \subseteq \mathbf{B}y$. Conversely, we suppose that (x,y) satisfies $\mathbf{A}x \subseteq \mathbf{B}y$. Then, $Ax \in \mathbf{B}y$ for all $A \in \mathbf{A}$. Therefore, Ax = By for all $A \in \mathbf{A}$ and for some $B \in \mathbf{B}$. Hence, (x,y) is a tolerance solution.

Theorem 2. A vector (x, y) is a tolerance solution of $\mathbf{A}x = \mathbf{B}y$ if and only if it satisfies

$$|A_c x - B_c y| \le \Delta_B |y| - \Delta_A |x|. \tag{7}$$

Proof. Let (x, y) be a tolerance solution of Ax = By. By Theorem 1 and (2),

$$[A_c x - \Delta_A |x|, A_c x + \Delta_A |x|] = \mathbf{A} x \subseteq \mathbf{B} y = [B_c y - \Delta_B |y|, B_c y + \Delta_B |y|].$$

Thus,

$$B_c y - \Delta_B |y| \le A_c x - \Delta_A |x| \le A_c x + \Delta_A |x| \le B_c y + \Delta_B |y|,$$

which implies

$$-(\Delta_B|y| - \Delta_A|x|) \le A_c x - B_c y \le \Delta_B|y| - \Delta_A|x|. \tag{8}$$

Conversely, let (x, y) satisfy condition (7) Then, it gives (8), which means

$$\mathbf{B}y = B_c y - \Delta_B |y| \le A_c x - \Delta_A |x| = \underline{\mathbf{A}}x$$

and

$$\overline{Ax} = A_c x + \Delta_A |x| \le B_c y + \Delta_B |y| = \overline{By}.$$

Therefore, $\mathbf{A}x \subseteq \mathbf{B}y$ and (x,y) becomes a tolerance solution.

We can use inequality (7) to verify whether a given vector (x, y) is a tolerance solution to our two-sided interval linear system. The system of inequality (7) has the absolute terms |x| and |y|, which means it is not the system of linear inequalities, in general. It depends on the signs of x and y components. However, when we consider the nonnegative domain of vector variables x and y, the inequality becomes a simple form as

$$\Delta_A x - \Delta_B y \le A_c x - B_c y \le \Delta_B y - \Delta_A x$$

which is obtained by substituting |x| = x and |y| = y. This turns into

$$(B_c - \Delta_B)y \le (A_c - \Delta_A)x$$
 and $(A_c + \Delta_A)x \le (\Delta_B + B_c)y$. (9)

According to (1), the inequalities (9) can be concluded as the corollary below.

Corollary 1. Let x and y be nonnegative vector variables. A vector (x, y) is a tolerance solution of $\mathbf{A}x = \mathbf{B}y$ if and only if it satisfies

$$\underline{B}y \le \underline{A}x \text{ and } \overline{A}x \le \overline{B}y.$$

The similar statements of Theorems 1-2 and Corollary 1 for a control solution can be done easily by interchanging the roles of "A and B" and "x and y".

Theorem 3. A vector (x, y) is a tolerance-control solution of $\mathbf{A}x = \mathbf{B}y$ if and only if it satisfies

$$|A_c x - B_c y| \le \Delta_A |x| + \Delta_B |y| - 2\delta, \tag{10}$$

where δ is a vector in \mathbb{R}^m with the following components

$$\delta_i = \begin{cases} (\Delta_A |x|)_i, & \text{if } i \in P; \\ (\Delta_B |y|)_i, & \text{if } i \in M \setminus P, \end{cases}$$

and
$$P = \{i \in M : (\mathbf{A}x)_i \subseteq (\mathbf{B}y)_i\}, M = \{1, 2, \dots, m\}.$$

Proof. Assume that (x, y) is a tolerance-control solution of $\mathbf{A}x = \mathbf{B}y$. Let P be a subset of M such that row $i \in P$ of the system is tolerance. Then, the other $i \in M \setminus P$ of the system is control. Using Theorem 2, we have

$$|(A_c x)_i - (B_c y)_i| \le (\Delta_B |y|)_i - (\Delta_A |x|)_i \text{ for each } i \in P$$
(11)

and

$$|(A_c x)_i - (B_c y)_i| \le (\Delta_A |x|)_i - (\Delta_B |y|)_i \text{ for each } i \in M \setminus P$$
 (12)

By putting Inequalities (11) and (12) together,

$$|A_c x - B_c y| \le \left(\Delta_A |x| - 2\sum_{i \in P} (\Delta_A |x|)_i e_i\right) + \left(\Delta_B |y| - 2\sum_{i \in M \setminus P} (\Delta_B |y|)_i e_i\right)$$

$$= \Delta_A |x| + \Delta_B |y| - 2\left(\sum_{i \in P} (\Delta_A |x|)_i e_i + \sum_{i \in M \setminus P} \Delta_B |y|)_i e_i\right)$$

$$= \Delta_A |x| + \Delta_B |y| - 2\delta,$$

where e_i is a m-column vector containing 1 at the i^{th} row and 0 elsewhere. Conversely, let (x, y) satisfy condition (10), which implies Inequalities (11) and (12). Therefore, (x, y) turns into a tolerance-control solution.

3 Semantics of Tolerance and Control Solutions

Looking at the mathematical definitions in the previous section, it may seem that tolerance and control solutions are algebraically the same. So, why would we need to define them both? This is because one interval information could be more important than another. There are some missing details in Definition 2 about the priority of the interval information \boldsymbol{A} and \boldsymbol{B} , which may not be represented clearly by just the mathematical quantification; "for all" and "for some".

To be able to achieve a control solution of the system Ax = By and justify the semantics of the word "control", the boundary matrices \underline{A} and \overline{A} of A must be more important than the ones of B. It could be the interval information that is given by an expert so that any quantity on the right must be controlled in the boundary of quantities on the left.

Similarly, to get the semantics of the word "tolerance" in a tolerance solution, the boundary matrices \underline{A} and \overline{A} of A must be more important. They provide the range of left hand side quantities. Moreover, the situation to come up with a tolerance solution is such that we need every element in the range to be under the control of the range of the right hand side quantities. In other words, the left hand side quantities tolerate themselves within the range of the right hand side quantities.

For those reasons, we cannot simply substitute \boldsymbol{A} with \boldsymbol{B} and \boldsymbol{x} with \boldsymbol{y} , and infer that tolerance and control solutions are the same. Therefore, the priority of the information \boldsymbol{A} and \boldsymbol{B} need to be stated in their definitions as rewritten below.

Definition 3. Let an interval matrix A play more important role to the twosided interval system of linear equations Ax = By. A vector $(x, y) \in \mathbb{R}^n \times \mathbb{R}^p$ is called

- 1. a tolerance solution of Ax = By if for each $A \in A$ there exists $B \in B$ such that Ax = By. It is in the sense that the range of Ax tolerates within the range of By,
- 2. a control solution of Ax = By if for each $B \in B$ there exists $A \in A$ such that Ax = By. It is in the sense that the range of Ax controls the range of By.

4 Applications on Tolerance and Control Solutions of a Two-Sided Interval Linear System

In this section, we illustrate two small examples to show the difference between tolerance and control solutions. These examples could be parts of any relevant application systems.

 Problem Statement 1 is formulated as a tolerance solution to a two-sided interval linear system, which can be solved by system of inequalities. Problem Statement 2 is formulated as a quadratic programming with control solutions to two-sided interval linear constraints, whose numerical example is demonstrated in Example 1.

Problem Statement 1. An animal food manufacturing company has its own quality control for its three formulas of chicken food: I, II, and III. These formulas must be checked for the level of carbohydrate, fiber, protein, and vitamins per kilogram per bag. The qualified bags must have the nutrients within the given boundaries as shown in the table below. The company also has a nutritionist who can advise customers about using these three formulas of chicken food for meat and egg chickens. To raise chickens for healthy meat and chickens for healthy eggs, those chickens should have nutrients within the range represented also in the table.

Type of nutrients	Interval amount of nutrient (kg)						
	Formula I	Formula II	Formula III	Meat chicken	Egg chicken		
1. Carbohydrate	$[\underline{a}_{11},\overline{a}_{11}]$	$[\underline{a}_{12},\overline{a}_{12}]$	$[\underline{a}_{13}, \overline{a}_{13}]$	$\left[\underline{b}_{11},\overline{b}_{11} ight]$	$\left[\underline{b}_{12},\overline{b}_{12} ight]$		
2. Fiber	$[\underline{a}_{21},\overline{a}_{21}]$	$[\underline{a}_{22},\overline{a}_{22}]$	$[\underline{a}_{23}, \overline{a}_{23}]$	$\left[\underline{b}_{21}, \overline{b}_{21}\right]$	$\left[\underline{b}_{22},\overline{b}_{22}\right]$		
3. Protein	$[\underline{a}_{31}, \overline{a}_{31}]$	$[\underline{a}_{32},\overline{a}_{32}]$	$[\underline{a}_{33}, \overline{a}_{33}]$	$\left[\underline{b}_{31},\overline{b}_{31}\right]$	$\left[\underline{b}_{32},\overline{b}_{32}\right]$		
4. Vitamins	$[\underline{a}_{41}, \overline{a}_{41}]$	$[\underline{a}_{42},\overline{a}_{42}]$	$[\underline{a}_{43}, \overline{a}_{43}]$	$\left[\underline{b}_{41},\overline{b}_{41} ight]$	$\left[\underline{b}_{42},\overline{b}_{42} ight]$		

The nutritionist suggests customers to mix three formulas together before feeding. The total amount of nutrients in the mixed chicken food must be within the range of the total amount of all needed nutrients to guarantee that all chickens provide healthy products. This relationship can be represented as a two-sided interval linear system (13)

where x_j is the amount of bags of animal food formula I, II and III that should be mixed, when j = 1, 2, 3, respectively, y_k is the amount of meat and egg chickens that the customer should raise, when k = 1, 2, respectively.

From the company's point of view, it is important to control the nutrients in each formula of chicken foods. It turns out that the set of solutions to System (13) is the set of tolerance solutions as the total of the nutrients created by the mixed food must be within the range of healthy nutrients. On the other hand, when considering the customer's side, the interval information about the nutrients needed for each chicken is the priority for the customer. In this case, the customer must control the nutrients in the mixed food by the range of the total amount of all needed nutrients

for healthy chicken meat and eggs. Then, the set of solutions to System (13) can also be viewed as the tolerance solution set. Obviously, vector variables $x = (x_1, x_2, x_3)^T$ and $y = (y_1, y_2)^T$ are nonnegative. From Corollary 1, the tolerance solution (x, y) is obtained by solving the system of inequalities as follows:

$$\underline{b}_{i1}y_1 + \underline{b}_{i2}y_2 - \underline{a}_{i1}x_1 - \underline{a}_{i2}x_2 - \underline{a}_{i3}x_3 \le 0$$

and

$$\overline{a}_{i1}x_1 + \overline{a}_2x_2 + \overline{a}_{i3}x_3 - \overline{b}_{i1}y_1 - \overline{b}_{i2}y_2 \le 0,$$

for each $i \in \{1, 2, 3, 4\}$.

Problem Statement 2. A famous family bakery shop sells homemade fruit cakes and fruit tarts. The owner uses three grades of mixed dried berries: A, B, and C, by mixing them together to get the best dessert according to the family recipe. Suppose the owner determines the initial amount of fruit cakes and fruit tarts he/she wanted to make as α_1 and α_2 , respectively, and sets up the initial amount of mixed dried berries grades A, B and C that he/she aims to buy as β_1 , β_2 and β_3 , respectively. One kilogram of each mixed dried fruit grade contains dried blueberries, dried cranberries, dried raspberries, and dried strawberries in the different uncertain quantities. The table below provides the interval quantities per kilogram of each mixed dried berries grade.

Type of dried berries	Interval amount of berry (kg)			
	Grade A	Grade B	Grade C	
1. Blueberries	$\left[\underline{b}_{11},\overline{b}_{11} ight]$	$\left[\underline{b}_{12},\overline{b}_{12} ight]$	$\left[\underline{b}_{13}, \overline{b}_{13}\right]$	
2. Cranberries	$\left[\underline{b}_{21}, \overline{b}_{21}\right]$	$\left[\underline{b}_{22}, \overline{b}_{22}\right]$	$\left[\underline{b}_{23}, \overline{b}_{23}\right]$	
3. Raspberries	$\left[\underline{b}_{31},\overline{b}_{31}\right]$	$\left[\underline{b}_{32},\overline{b}_{32}\right]$	$\left[\underline{b}_{33}, \overline{b}_{33}\right]$	
4. Strawberries	$\left[\underline{b}_{41},\overline{b}_{41} ight]$	$\left[\underline{b}_{42},\overline{b}_{42} ight]$	$\left[\underline{b}_{43}, \overline{b}_{43}\right]$	

To control the quality of the desserts, the recipe says that the dessert must contain each type of dried berries in a certain level as shown in the following table.

Type of dried berries	Interval amount of berry (kg)		
	Fruit cake	Fruit tart	
1. Blueberries	$[\underline{a}_{11}, \overline{a}_{11}]$	$[\underline{a}_{12},\overline{a}_{12}]$	
2. Cranberries	$[\underline{a}_{21}, \overline{a}_{21}]$	$[\underline{a}_{22},\overline{a}_{22}]$	
3. Raspberries	$[\underline{a}_{31}, \overline{a}_{31}]$	$[\underline{a}_{32},\overline{a}_{32}]$	
4. Strawberries	$[\underline{a}_{41}, \overline{a}_{41}]$	$[\underline{a}_{42},\overline{a}_{42}]$	

The relationship of how many kilograms of the mixed dried berries grades A, B, and C that the shop should have and how many fruit cakes and fruit tarts

that the shop should make to guarantee the overall quality of the dessert becomes a two-sided interval linear system as follows:

where x_j is the amount of fruit cakes and fruit tarts that should be made, when j = 1, 2, respectively,

 y_k is the amount of mixed dried berries grades A, B and C, when k = 1, 2, 3, respectively.

The quality of desserts is very important to the bakery shop. The owner wants to re-evaluate the amount x_1, x_2 and y_1, y_2, y_3 to make sure that the mixed dried berries the shop has would be covered by the range of the total amount of the mixed fruits in all dessert items. This way the shop would be able to guarantee the quality of the desserts. The set of solutions satisfies System (14) is considered to be the control solution set. In a situation when the owner wants the amount x_1, x_2 and y_1, y_2, y_3 as close as possible to the given values α_1, α_2 and $\beta_1, \beta_2, \beta_3$, respectively, it is the same as to minimize the function below:

$$\|(x,y)^{T} - (\alpha,\beta)^{T}\|^{2} = \sum_{j=1}^{2} (x_{j} - \alpha_{j})^{2} + \sum_{k=1}^{3} (y_{k} - \beta_{k})^{2}$$
$$= \sum_{j=1}^{2} (x_{j}^{2} - 2\alpha_{j}x_{j} + \alpha_{j}^{2}) + \sum_{k=1}^{3} (y_{k}^{2} - 2\beta_{k}y_{k} + \beta_{k}^{2}). \quad (15)$$

It is sufficient to remove the constant terms α_j^2 and β_k^2 from Eq. (15) for every $j \in \{1,2\}$ and $k \in \{1,2,3\}$. Therefore, we can model this problem as a quadratic program P_1 with a two-sided interval linear constraint (14), whose constraints are obtained by using the statement of Corollary 1 for a control solution in the following way:

$$\begin{split} P_1: & \textit{minimize } f(x,y) = \sum_{j=1}^2 (x_j^2 - 2\alpha_j x_j) + \sum_{k=1}^3 (y_k^2 - 2\beta_k y_k) \\ & \textit{subject to } \sum_{j=1}^2 \underline{a}_{ij} x_j - \sum_{k=1}^3 \underline{b}_{ik} y_k \leq 0, \quad \forall i \in \{1,2,3,4\} \\ & - \sum_{j=1}^2 \overline{a}_{ij} x_j + \sum_{k=1}^3 \overline{b}_{ik} y_k \leq 0, \quad \forall i \in \{1,2,3,4\} \\ & x_j, y_k \geq 0. \quad \forall j \in \{1,2\} \ \forall k \in \{1,2,3\} \end{split}$$

The Lagrangian function for Problem P_1 can be written by

$$L(z,\mu) = c^T z + \frac{1}{2} z^T Q z + \mu (A'z - b'), \tag{16}$$

446

where

$$c = \begin{pmatrix} -2\alpha_1 \\ -2\alpha_2 \\ -2\beta_1 \\ -2\beta_2 \\ -2\beta_3 \end{pmatrix}, \quad Q = \begin{pmatrix} 2 & 0 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 \\ 0 & 0 & 0 & 2 & 0 \\ 0 & 0 & 0 & 0 & 2 \end{pmatrix}, \quad z = \begin{pmatrix} x_1 \\ x_2 \\ y_1 \\ y_2 \\ y_3 \end{pmatrix}$$

and

and $\mu = (\mu_1, \mu_2, \dots, \mu_8) \ge 0$ is the Lagrange multiplier with 5-dimensional row vector. Let u and v be surplus and slack variables to the inequalities

$$c + Qz + (A')^T \mu^T \ge 0$$
 and $A'z - b' \le 0$,

respectively. As shown in [5], we can represent (16) in the following linear constraints form:

$$Qz + (A')^T \mu^T - u = -c, (17)$$

$$A'z + v = b', (18)$$

$$z \ge 0, \ \mu \ge 0, \ u \ge 0, \ v \ge 0,$$
 (19)

$$u^T z = 0, \quad \mu v = 0,$$
 (20)

where the equations shown in (20) prescribe complementary slackness. Since it can clearly be seen that the matrix Q is positive definite, the conditions (17)–(20) are necessary and sufficient for a global optimum. To create the appropriate linear program, we add thirteen artificial variables a_1, a_2, \ldots, a_{13} to each constraint of (17) and (18) together with minimizing their sum, that is,

P₂: minimize
$$a_1 + a_2 + \dots + a_{13}$$

subject to $Qz + (A')^T \mu^T - u + a' = -c$,
 $A'z + v + a'' = b'$,
 $z \ge 0, \quad \mu \ge 0, \quad u \ge 0, \quad v \ge 0$,
 $u^T z = 0, \quad \mu v = 0$,

where $a' = (a_1, a_2, \dots, a_5)^T$ and $a'' = (a_6, a_7, \dots, a_{13})^T$. Hence, the optimal solution to the quadratic program P_1 is found out by solving the linear program P_2 .

Type of dried berries	Interval amount of berry (kg)					
	Fruit cake	Fruit tart	Grade A	Grade B	Grade C	
1. Blueberries	[0.25, 0.32]	[0.20, 0.41]	[0.15, 0.30]	[0.60, 0.70]	[0.22, 0.25]	
2. Cranberries	[0.05, 0.18]	[0.26, 0.35]	[0.20, 0.20]	[0.23, 0.35]	[0.08, 0.12]	
3. Raspberries	[0.03, 0.15]	[0.24, 0.64]	[0.10, 0.16]	[0.04, 0.10]	[0.45, 0.55]	
4. Strawberries	[0.04, 0.36]	[0.34, 0.45]	[0.32, 0.48]	[0.14, 0.28]	[0.18, 0.20]	

Example 1. According to Problem statement 2, we provide the numerical information in the following table.

- The initial amount of fruit cakes and fruit tarts is $\alpha_1 = 16$ and $\alpha_2 = 24$, respectively.
- The initial amount of mixed dried berries grade A, B and C is $\beta_1 = 15$, $\beta_2 = 12$ and $\beta_3 = 10$, respectively.

As explained in Problem statement 2 with the above information, our problem can be represented by Problem P_2 , that is,

minimize $a_1 + a_2 + a_3 + a_4 + a_5 + a_6 + a_7 + a_8 + a_9 + a_{10} + a_{11} + a_{12} + a_{13}$ subject to

$$\begin{aligned} 2x_1 + 0.25\mu_1 + 0.05\mu_2 + 0.03\mu_3 + 0.04\mu_4 - 0.32\mu_5 - 0.18\mu_6 - 0.15\mu_7 - 0.36\mu_8 - u_1 + a_1 &= 32, \\ 2x_2 + 0.20\mu_1 + 0.26\mu_2 + 0.24\mu_3 + 0.34\mu_4 - 0.41\mu_5 - 0.35\mu_6 - 0.64\mu_7 - 0.45\mu_8 - u_2 + a_2 &= 48, \\ 2y_1 - 0.15\mu_1 - 0.20\mu_2 - 0.10\mu_3 - 0.32\mu_4 + 0.30\mu_5 + 0.20\mu_6 + 0.16\mu_7 + 0.48\mu_8 - u_3 + a_3 &= 30, \\ 2y_2 - 0.60\mu_1 - 0.23\mu_2 - 0.04\mu_3 - 0.14\mu_4 + 0.70\mu_5 + 0.35\mu_6 + 0.10\mu_7 + 0.28\mu_8 - u_4 + a_4 &= 24, \\ 2y_3 - 0.22\mu_1 - 0.08\mu_2 - 0.45\mu_3 - 0.18\mu_4 + 0.25\mu_5 + 0.12\mu_6 + 0.55\mu_7 + 0.20\mu_8 - u_5 + a_5 &= 20, \\ 0.25x_1 + 0.20x_2 - 0.15y_1 - 0.60y_2 - 0.22y_3 + v_1 + a_6 &= 0, \\ 0.05x_1 + 0.26x_2 - 0.20y_1 - 0.23y_2 - 0.08y_3 + v_2 + a_7 &= 0, \\ 0.03x_1 + 0.24x_2 - 0.10y_1 - 0.04y_2 - 0.45y_3 + v_3 + a_8 &= 0, \\ 0.04x_1 + 0.34x_2 - 0.32y_1 - 0.14y_2 - 0.18y_3 + v_4 + a_9 &= 0, \\ -0.32x_1 - 0.41x_2 + 0.30y_1 + 0.70y_2 + 0.25y_3 + v_5 + a_{10} &= 0, \\ -0.18x_1 - 0.35x_2 + 0.20y_1 + 0.35y_2 + 0.12y_3 + v_6 + a_{11} &= 0, \\ -0.15x_1 - 0.64x_2 + 0.16y_1 + 0.10y_2 + 0.55y_3 + v_7 + a_{12} &= 0, \\ -0.36x_1 - 0.45x_2 + 0.48y_1 + 0.28y_2 + 0.20y_3 + v_8 + a_{13} &= 0, \end{aligned}$$

where all variables are nonnegative and complementary conditions are satisfied. By applying the simplex method, the optimal solution to this problem is eventually displayed as

$$\begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 18.49 \\ 22.16 \end{pmatrix}$$
 and $\begin{pmatrix} y_1 \\ y_2 \\ y_3 \end{pmatrix} = \begin{pmatrix} 17.49 \\ 11.01 \\ 8.19 \end{pmatrix}$.

In this section we have presented two particular situations to show the difference between tolerance and control solutions together with the processes for solving them. The tolerance solution as Problem Statement 1 can directly be accomplished using system of inequalities. The quadratic programming problem constrained with control solutions as Problem Statement 2 was transformed to become an appropriate linear program. Then, the simplex method enables us to obtain the optimal solution of the problem as displayed in the above example.

5 Conclusion

This paper presents the concepts of tolerance and control solutions of a two-sided interval linear system together with their semantics to the system $\mathbf{A}x = \mathbf{B}y$. The conditions to verify that (x,y) is a tolerance or a control or a tolerance-control solution are also achieved in Theorems 1–3. In application problems, the vectors x and y are normally specified by nonnegative variables. Therefore, we simplify Theorem 2 to be Corollary 1 and applied it to implement a two-sided interval linear system and a quadratic programming with two-sided interval linear constraints. This work should be beneficial to any applications that have the restrictions in the format of two-sided interval linear system.

References

- 1. Beaumont, O.: Solving interval linear systems with linear programming techniques. Linear Algebra Appl. **281**(1–3), 293–309 (1998)
- Fiedler, M., Nedoma, J., Ramik, J., Rohn, J., Zimmermann, K.: Linear Optimization Problems with Inexact Data. Springer, New York (2006). https://doi.org/10.1007/0-387-32698-7
- 3. Gerlach, W.: Zur lösung linearer ungleichungssysteme bei störimg der rechten seite und der koeffizientenmatrix. Optimization 12(1), 41–43 (1981)
- Hladík, M.: Weak and strong solvability of interval linear systems of equations and inequalities. Linear Algebra Appl. 438(11), 4156–4165 (2013)
- Jensen, P.A., Bard, J.F.: Operations Research Models and Methods, vol. 1. Wiley, Hoboken (2003)
- Leela-apiradee, W.: New characterizations of tolerance-control and localized solutions to interval system of linear equations. J. Comput. Appl. Math. 355, 11–22 (2019)
- 7. Moore, R.E.: Interval Analysis. Prince-Hall, Englewood Cliffs (1969)
- 8. Nuding, E., Wilhelm, J.: Über gleichungen und über lösungen. Zeitschrift für Angewandte Mathematik und Mechanik **52**. T188–T190 (1972)
- Oettli, W., Prager, W.: Compatibility of approximate solution of linear equations with given error bounds for coeficients and right-hand sides. Numer. Math. 6, 405–409 (1964). https://doi.org/10.1007/BF01386090
- Prokopyev, O.A., Butenko, S., Trapp, A.: Checking solvability of systems of interval linear equations and inequalities via mixed integer programming. Eur. J. Oper. Res. 199(1), 117–121 (2009)
- 11. Rohn, J.: Input-Output Planning with Inexact Data, vol. 9. Albert-Ludwigs-Universität, Freiburg (1978). Freiburger Intervall-Berichte
- Shary, S.P.: On controlled solution set of interval algebraic systems. Interval Comput. 6, 66–75 (1992)
- Tian, L., Li, W., Wang, Q.: Tolerance-control solutions to interval linear equations.
 In: International Conference on Artificial Intelligence and Software Engineering,
 ICAISE 2013, pp. 18–22. Atlantis Press (2013)