Complete intersection toric ideals of oriented graphs and chorded-theta subgraphs

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Abstract Let $G = (V, E)$ be a finite, simple graph. We consider for each oriented graph $G_{\mathcal{O}}$ associated to an orientation $\mathcal{O}$ of the edges of G , the toric ideal $P_{G_{\mathcal{O}}}$. In this paper we study those graphs with the property that $P_{G_{\mathcal{O}}}$ is a binomial complete intersection, for all $\mathcal{O}$. These graphs are called CIO graphs. We prove that these graphs can be constructed recursively as clique-sums of cycles and/or complete graphs. We introduce chorded-theta subgraphs and some of their properties. Also we establish that the CIO graphs are determined by the property that each chorded-theta has a transversal triangle. Finally we explicitly give the minimal forbidden induced subgraphs that characterize these graphs, these families of forbidden graphs are: prisms, pyramids, thetas and a particular family of wheels that we call θ -partial wheels.

Keywords Toric ideals · Oriented graphs · Complete intersections · Forbidden induced subgraphs

1 Introduction

Let $G = (V, E)$ be a graph with $|E(G)| = q$ and whose vertex set is given by $V(G) = \{x_1, \dots, x_n\}$. An orientation $\mathcal{O}$ of the edges of G is an assignment of a direction to each edge of G . Let $D = G_{\mathcal{O}}$ denote the oriented graph associated to an orientation $\mathcal{O}$ of the edges of G . In particular if G is a complete graph, $G_{\mathcal{O}}$ is called a *tournament*.

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To each oriented edge $e = (x_i, x_j)$ of D , we associate the vector $v_e \in \{0, 1, -1\}^n$ defined as follows: the i th entry is -1 , the j th entry is 1 , and the remaining entries are zero. The *incidence matrix* A_D of D is the $n \times q$ matrix whose columns are the vectors of the form v_e , with e an edge of D . The set of column vectors of A_D will be denoted by $\mathcal{A} = \{v_1, \dots, v_q\}$.

Consider the *edge subring* $k[D] := k[x^{v_1}, \dots, x^{v_q}] \subset k[x_1^{\pm 1}, \dots, x_n^{\pm 1}]$ of D , where $v_i = (v_i^1, \dots, v_i^n)$ and $x^{v_i} = x_1^{v_i^1} \cdots x_n^{v_i^n}$. Let $E(D) = \{t_1, \dots, t_q\}$ be the edge set of D . There is an epimorphism of k -algebras given by

$$\varphi: k[t_1, \dots, t_q] \longrightarrow k[D], \quad \text{where } t_i \longmapsto x^{v_i},$$

and $k[t_1, \dots, t_q]$ is a polynomial ring. The kernel of φ , denoted by P_D , is called the *toric ideal* of D . This ideal has been thoroughly studied in [5, 6, 12]. The toric ideal P_D is a prime ideal of height $q - n + r$, where r is the number of connected components of D . P_D is generated by binomials and $k[D]$ is a normal domain (see [6, 18]). Thus any minimal generating set of P_D must have at least $q - n + r$ elements, by the Principal Ideal Theorem (see [1]). If P_D can be generated exactly by $q - n + r$ polynomials it is called a *complete intersection*. If these polynomials are binomials then P_D is called a *binomial complete intersection*. In Sect. 2 we study the binomials of P_D . In particular, we prove that the primitive binomials of P_D are the binomials associated to cycles. We also recover the result given in [6] that P_D is generated by binomials associated to cycles without chords. If two graphs G and H each contain cliques of equal size, the *clique-sum* of G and H is formed from their disjoint union by identifying pairs of vertices in these two cliques to form a single shared clique.

It is known [6, 12] that any graph has at least one acyclic orientation such that the corresponding toric ideal is a binomial complete intersection. It is natural to ask what is the class of graphs with the property that $P_{G_{\mathcal{O}}}$ is generated by exactly $q - n + r$ binomials for every orientation $\mathcal{O}$ of G . For ease of exposition these graphs are called *CIO* graphs. In the bipartite case a graph G is in this class if and only if G is a ring graph (see [6]). Ring graphs are introduced in [5], they are those graphs for which their non trivial blocks (blocks that are not vertices or bridges) can be constructed as 2-clique-sums of cycles. In Sect. 3 we show that the *CIO* property is closed under induced subgraphs and we prove that G is a *CIO* graph if and only if its blocks are *CIO* graphs.

In Sect. 4 we introduce the family of *chorded-theta* subgraphs and their transversal triangles. In this section we study the graphs G that satisfy: for every chorded-theta of G there exists a transversal triangle ($\forall \theta \exists \Delta$ -property). We call this type of graphs *theta-ring* graphs. In Theorem 4 we prove that these graphs can be constructed as clique-sums of cycles and complete graphs. Furthermore we show that they can be obtained as 0, 1, 2-clique-sums of chordal graphs and cycles. *Theta-ring* graphs are closed under induced subgraphs and we obtain the corresponding set of minimal forbidden induced subgraphs that characterize this family. These minimal forbidden induced subgraphs are: prisms, pyramids, thetas and θ -partial wheels (partial wheels that are chorded-theta graphs).

Our main result in Sect. 5 is Theorem 6: $P_{G_{\mathcal{O}}}$ is a binomial complete intersection for every orientation $\mathcal{O}$ of G if and only if G is a *theta-ring* graph, i.e., the *CIO*

property is equivalent to the $\forall\theta\exists\Delta$ -property. As a consequence of this result, we see that $P_{G_{\mathcal{O}}}$ is a binomial complete intersection if $G_{\mathcal{O}}$ is a tournament. Using the results given in Sect. 4 and Sect. 5 we obtain the following equivalences:

$$G \text{ is a } \text{CIO graph} \Leftrightarrow (G \text{ is a theta-ring graph} \wedge (G \text{ has the } \forall\theta\exists\Delta\text{-property})) \Leftrightarrow G \text{ is obtained by clique-sums of complete graphs and cycles}$$

In particular, we recover that ring graphs are CIO graphs and that the converse also holds in the bipartite case.

The paper is essentially self contained. For unexplained terminology and notation on graph theory and toric ideals we refer to [3, 6, 7, 18]. Some references for toric ideals associated to graphs (without orientation) are [5, 6, 8, 9, 13, 15–17].

2 Toric ideals of oriented graphs

Let G be a graph and consider the oriented graph $D = G_{\mathcal{O}}$ associated to an orientation $\mathcal{O}$ of the edges of G . Let $\alpha = (\alpha_1, \dots, \alpha_q) \in \mathbb{R}^q$, the support of α is the set $\text{supp}(\alpha) = \{i \mid \alpha_i \neq 0\}$. Furthermore $\alpha = \alpha_+ - \alpha_-$, where α_+ and α_- are two non-negative vectors with disjoint support. Recall that $\{v_1, \dots, v_q\}$ is the set of column vectors of the incidence matrix A_D . We denote by P_D the kernel of the epimorphism of k -algebras

$$\varphi: k[t_1, \dots, t_q] \longrightarrow k[D], \quad \text{where } t_i \longmapsto x^{v_i}.$$

On the other hand we have the following linear map:

$$\psi: \mathbb{Z}^q \longrightarrow \mathbb{Z}^n, \quad \text{where } e_i \longmapsto v_i,$$

i.e., ψ is the linear map associated to the incidence matrix A_D of D . If $0 \neq \alpha \in \mathbb{Z}^q$ we associate to it the binomial $t_{\alpha} = t^{\alpha_+} - t^{\alpha_-}$. Notice that $\alpha \in \ker(\psi)$ if and only if $t_{\alpha} \in P_D$. Given a cycle $C = (x_{j_1}, x_{j_2}, \dots, x_{j_k}, x_{j_{k+1}} = x_{j_1})$ of D , we split the edge set of C in two disjoint sets C_+ and C_- , where C_+ is the set of clockwise oriented edges, i.e., $C_+ = \{t_s \in E(C) \mid \text{there exist } i \text{ such that } t_s = (x_{j_i}, x_{j_{i+1}})\}$ and $C_- = C \setminus C_+$. Associated to C we have the binomial

$$t_C = \prod_{t_i \in C_+} t_i - \prod_{t_i \in C_-} t_i.$$

This binomial belongs to P_D . If $C_+ = \emptyset$ or $C_- = \emptyset$ we set $\prod_{t_i \in C_+} t_i = 1$ or $\prod_{t_i \in C_-} t_i = 1$.

In this section we study some properties of the binomials in P_D . In Theorem 1 we prove that if b is a binomial in P_D , then there exists a binomial b' associated to a cycle whose monomials divide the monomials of b . Using this result we prove that the primitive binomials in P_D are the ones associated with cycles and we recover the result in [6] that P_D is generated by the binomials corresponding to cycles without chords.

Lemma 1 Let $t^\alpha - t^\beta$ be a binomial in P_D with $\gcd(t^\alpha, t^\beta) = 1$. If $i \in \text{supp}(\alpha)$ and $t_i = (x_{j_1}, x_{j_2})$ then there exist edges t_{i_1} and t_{i_2} both different from t_i such that $x_{j_1} \in V(t_{i_1})$ and $x_{j_2} \in V(t_{i_2})$. Furthermore if $t_{i_2} = (x_{j_2}, x_{j_3})$, then $t_{i_2} \mid t^\alpha$ and if $t_{i_2} = (x_{j_3}, x_{j_2})$ then $t_{i_2} \mid t^\beta$.

Proof Since $t^\alpha - t^\beta \in P_D$ then

$$\varphi(t_1^{\alpha_1}) \cdots \varphi(t_q^{\alpha_q}) = \varphi(t^\alpha) = \frac{m_1}{m_2} = \varphi(t^\beta) = \varphi(t_1^{\beta_1}) \cdots \varphi(t_q^{\beta_q})$$

where $\alpha = (\alpha_1, \dots, \alpha_q)$ and $\beta = (\beta_1, \dots, \beta_q)$. We can assume that m_1 and m_2 are monomials with $\gcd(m_1, m_2) = 1$. Since $t_i = (x_{j_1}, x_{j_2})$, we have $\varphi(t_i^{\alpha_i}) = \frac{x_{j_2}^{\alpha_i}}{x_{j_1}^{\alpha_i}}$ and $\alpha_i \neq 0$. If $x_{j_2} \mid m_1$ then there exists an edge t_{i_2} such that $\beta_{i_2} \neq 0$ and $\varphi(t_{i_2}) = \frac{x_{j_2}}{x_{j_3}}$ for some x_{j_3} . Thus $t_{i_2} = (x_{j_3}, x_{j_2})$ and $t_{i_2} \mid t^\beta$. Furthermore $t_i \neq t_{i_2}$ because $\gcd(t^\alpha, t^\beta) = 1$. Now, if x_{j_2} does not divide m_1 then there exists an edge t_{i_2} such that $\varphi(t_{i_2}) = \frac{x_{j_4}}{x_{j_2}}$ with $\alpha_{i_2} \neq 0$ and $i \neq i_2$. Hence, $t_{i_2} = (x_{j_2}, x_{j_4})$ and $t_{i_2} \mid t^\alpha$, so $x_{j_2} \in V(t_{i_2})$ and $t_{i_2} \neq t_i$.

In a similar way we prove that there exists an edge t_{i_1} different from t_i such that $x_{j_1} \in V(t_{i_1})$. $\square$

Theorem 1 If $0 \neq t^\alpha - t^\beta \in P_D$ then there exists a cycle C such that $t_C = t^{\alpha'} - t^{\beta'}$, $t^{\alpha'} \mid t^\alpha$ and $t^{\beta'} \mid t^\beta$.

Proof We can suppose that $\gcd(t^\alpha, t^\beta) = 1$, otherwise we would have $\gcd(t^\alpha, t^\beta) = m$. Thus, $t^\alpha = mt^\gamma$, $t^\beta = mt^w$, $\gcd(t^\gamma, t^w) = 1$ and $t^\gamma - t^w \in P_D$. We set $\alpha = (\alpha_1, \dots, \alpha_q)$ and $\beta = (\beta_1, \dots, \beta_q)$. Since $t^\alpha - t^\beta \neq 0$, then there exists i such that $\alpha_i \neq 0$ or $\beta_i \neq 0$. We can suppose that $\alpha_i \neq 0$. By Lemma 1 if $x_{j_1} \in V(t_i)$ there exists an edge t_{i_1} different from t_i such that $x_{j_1} \in V(t_{i_1})$ and $t_{i_1} \mid t^\alpha t^\beta$. Now if $V(t_{i_1}) = \{x_{j_1}, x_{j_2}\}$ then by Lemma 1 there exists an edge t_{i_2} different from t_{i_1} such that $x_{j_2} \in V(t_{i_2})$ and $t_{i_2} \mid t^\alpha t^\beta$. We continue with this process and we obtain vertices $x_{j_1}, x_{j_2}, \dots$ such that $\{x_{j_k}, x_{j_{k+1}}\} \in E(G)$. We take $l = \max\{i \mid x_{j_i} \neq x_{j_k} \text{ for } k < i\}$, then $x_{j_{l+1}} = x_{j_s}$ for some $s \leq l$. Therefore $C = (x_{j_s}, x_{j_{s+1}}, \dots, x_{j_l}, x_{j_{l+1}} = x_{j_s})$ is a cycle in G and by construction if $t_C = t^{\alpha'} - t^{\beta'}$ then $t^{\alpha'} t^{\beta'} \mid t^\alpha t^\beta$. But, by Lemma 1 we have $t^{\alpha'} \mid t^\alpha$ and $t^{\beta'} \mid t^\beta$. $\square$

Definition 1 Let $C = (y_1, y_2, \dots, y_k, y_{k+1} = y_1)$ be a cycle of G . C is called an oriented cycle in $D = G_\odot$ if $(y_i, y_{i+1}) \in E(D)$ for $i = 1, \dots, k$ (i.e. $C_+ = \emptyset$ or $C_- = \emptyset$). Similarly a path $L = (y_1, \dots, y_s)$ of G is called an oriented path in D if $(y_i, y_{i+1}) \in E(D)$ for $i = 1, \dots, s-1$. If L is a path, we denote by L^o the path $(y_2, \dots, y_{s-1})$, if $s = 2$ then $L^o = \emptyset$.

Corollary 1 If $0 \neq 1 - t^\beta \in P_D$ then there exists an oriented cycle C in D .

Proof By Theorem 1, there is a cycle C such that $t_C = t^{\alpha'} - t^{\beta'}$, $t^{\alpha'} \mid 1$ and $t^{\beta'} \mid t^\beta$. Then $t^{\alpha'} = 1$ and C is an oriented cycle of D . $\square$

Proposition 1 Let f, g be binomials of P_D , such that $f = m_1 - m_2$ and $g = m_1 - m'_2$. If there is no oriented cycle in D and $m'_2 \mid m_2$ then $f = g$.

Proof There exists a monomial l such that $m_2 = lm'_2$ since $m'_2 \mid m_2$. On the other hand we have $0 = \varphi(f) = \varphi(m_1) - \varphi(l)\varphi(m'_2)$, then $\varphi(m_1) = \varphi(l)\varphi(m'_2)$. But $\varphi(m_1) = \varphi(m'_2)$ because $g \in P_D$. Hence, $\varphi(1) = \varphi(l)$ and $(1 - l) \in P_D$. Using Corollary 1 and since there is no oriented cycle in D we have $l = 1$ and $f = g$. $\square$

Corollary 2 If $0 \neq t^\alpha - t^\beta = f$ in P_D with $\gcd(t^\alpha, t^\beta) = 1$, then there exist cycles $C_1, \dots, C_s$ such that $t^\alpha = t^{\alpha_1} \dots t^{\alpha_s}$ and $t^\beta = t^{\beta_1} \dots t^{\beta_s}$ where $t_{C_i} = t^{\alpha_i} - t^{\beta_i}$ for $i = 1, \dots, s$.

Proof By Theorem 1 there exists a cycle C_1 with $t_{C_1} = t^{\alpha_1} - t^{\beta_1}$ such that $t^\alpha = t^{\alpha_1} t^{\alpha'}$ and $t^\beta = t^{\beta_1} t^{\beta'}$. Since $f = t^\alpha - t^\beta$ in P_D then $\varphi(f) = 0$ and

$$\varphi(t^{\alpha_1})\varphi(t^{\alpha'}) = \varphi(t^\alpha) = \varphi(t^\beta) = \varphi(t^{\beta_1})\varphi(t^{\beta'}).$$

But $t_{C_1} \in P_D$ then $\varphi(t^{\alpha_1}) = \varphi(t^{\beta_1})$. Hence $\varphi(t^{\alpha'}) = \varphi(t^{\beta'})$ and $f_1 = t^{\alpha'} - t^{\beta'} \in P_D$. Since $\gcd(t^{\alpha'}, t^{\beta'}) = 1$, if $f_1 = 0$ then $t^{\alpha'} = t^{\beta'} = 1$ and $f = t_{C_1}$. Now, if $f_1 \neq 0$ by Theorem 1 there exists a cycle C_2 with $t_{C_2} = t^{\alpha_2} - t^{\beta_2}$ such that $t^{\alpha'} = t^{\alpha_2} t^{\alpha''}$ and $t^{\beta'} = t^{\beta_2} t^{\beta''}$. Furthermore $t^\alpha = t^{\alpha_1} t^{\alpha_2} t^{\alpha''}$, $t^\beta = t^{\beta_1} t^{\beta_2} t^{\beta''}$ and $f_2 = t^{\alpha''} - t^{\beta''} \in P_D$. We continue this process to obtain the result. $\square$

Definition 2 A binomial $t^\alpha - t^\beta$ in P_D is called primitive if there exists no other binomial $t^{\alpha'} - t^{\beta'} \in P_D$ such that $t^{\alpha'}$ divides t^α and $t^{\beta'}$ divides t^β .

Proposition 2 The primitive binomials of P_D are the binomials of P_D associated to cycles.

Proof Let $g = t^\alpha - t^\beta$ be a primitive binomial. By Theorem 1 there exists a cycle C such that $t_C = t^{\alpha'} - t^{\beta'}$, $t^{\alpha'} \mid t^\alpha$ and $t^{\beta'} \mid t^\beta$. But g is primitive then $g = t_C$, i.e., g is associated to C . Therefore to obtain the result it is only necessary to prove that the binomials associated to cycles are primitive. So, let $t_C = t^\gamma - t^w$ be the binomial associated to C , then $\gcd(t^\gamma, t^w) = 1$. Suppose that there exists a binomial $f = t^{\gamma'} - t^{w'} \in P_D$ such that $t^{\gamma'} \mid t^\gamma$ and $t^{w'} \mid t^w$. Then, by Theorem 1 there exists a cycle C' such that $t_{C'} = t^{\gamma''} - t^{w''}$, $t^{\gamma''} \mid t^{\gamma'}$ and $t^{w''} \mid t^{w'}$. Hence, $t^{\gamma''} \mid t^\gamma$ and $t^{w''} \mid t^w$, this implies $E(C') \subset E(C)$. Thus, $C' = C$, $t_{C'} = t_C$ and $t_C = f$. Therefore t_C is primitive. $\square$

Proposition 3 The toric ideal P_D is generated by the binomials associated to cycles: $P_D = (t_C \mid C \text{ is cycle of } D)$.

Proof First we observe that

$$t^{\alpha_1} t^{\alpha_2} - t^{\beta_1} t^{\beta_2} = t^{\alpha_1} (t^{\alpha_2} - t^{\beta_2}) + t^{\beta_2} (t^{\alpha_1} - t^{\beta_1})$$

then $t^{\alpha_1} t^{\alpha_2} - t^{\beta_1} t^{\beta_2} \in (t^{\alpha_1} - t^{\beta_1}, t^{\alpha_2} - t^{\beta_2})$. Now, if $g = t^\alpha - t^\beta \in P_D$ then by Corollary 2, there exist cycles $C_1, \dots, C_s$ in D with $t_{C_i} = t^{\alpha_i} - t^{\beta_i}$ such that $t^\alpha = m t^{\alpha'} = m t^{\alpha_1} \dots t^{\alpha_s}$ and $t^\beta = m t^{\beta'} = m t^{\beta_1} \dots t^{\beta_s}$ where $m = \gcd(t^\alpha, t^\beta)$. By the first observation we have $g = m(t^{\alpha'} - t^{\beta'}) \in (t_{C_1}, \dots, t_{C_s})$. Therefore $P_D = (t_C \mid C \text{ is cycle of } D)$. $\square$

The last result also follows from the fact that the Graver basis is a generating set for the ideal.

Lemma 2 *Let C_1, C_2 be two cycles of D whose intersection is an oriented path P . Then $C_3 = (C_1 \cup C_2) \setminus P^o$ is a cycle and $t_{C_3} \in (t_{C_1}, t_{C_2})$.*

Proof We can assume that $C_1 = (x = x_0, x_1, \dots, x_{r_1} = y, y_1, \dots, y_{r_2}, x)$, $C_2 = (x = x_0, x_1, \dots, x_{r_1} = y, z_1, \dots, z_{r_3}, x)$ and $C_1 \cap C_2 = P = (x, x_1, \dots, x_{r_1-1}, y)$ where $(x_i, x_{i+1}) \in E(D)$ for $i = 0, \dots, r_1 - 1$, then $C_3 = (C_1 \cup C_2) \setminus P^o = (y, y_1, \dots, y_{r_2}, x, z_{r_3}, \dots, z_1, y)$. Thus, C_3 is a cycle. Without loss of generality we can suppose that $t_{i+1} = (x_i, x_{i+1})$ for $i = 0, \dots, r_1 - 1$. Hence, $t_{C_1} = t^\alpha t^{\alpha_1} - t^{\beta_1}$ and $t_{C_2} = t^\alpha t^{\alpha_2} - t^{\beta_2}$ where $\alpha = \sum_{i=1}^{r_1} e_i$. Furthermore $t_{C_3} = t^{\alpha_1} t^{\beta_2} - t^{\alpha_2} t^{\beta_1}$. But $t^{\alpha_1} t^{\beta_2} - t^{\alpha_2} t^{\beta_1} = t^{\alpha_2} t_{C_1} - t^{\alpha_1} t_{C_2}$ therefore, $t_{C_3} \in (t_{C_1}, t_{C_2})$. $\square$

Proposition 4 ([6]) P_D is generated by the set of binomials corresponding to cycles without chords.

Proof By Proposition 3 it is only necessary to prove that if

$$C = (x_1, x_2, \dots, x_s, x_{s+1} = x_1)$$

is a cycle of G and $e = \{x_i, x_j\}$ is a chord of C , then $t_C \in (t_{C_1}, t_{C_2})$ where $C_1 = (x_i, x_{i+1}, \dots, x_j, x_i)$ and $C_2 = (x_1, \dots, x_i, x_j, x_{j+1}, \dots, x_s, x_1)$. But the intersection of C_1 and C_2 is the edge e , then we obtain the result by Lemma 2. $\square$

3 CIO graphs and toric ideals

Let G be a graph, with q edges and r connected components such that $V(G) = \{x_1, \dots, x_n\}$. For ease of exposition we will say that G is CIO if the toric ideal $P_{G_{\mathcal{O}}}$ is a binomial complete intersection for each orientation $\mathcal{O}$. Recall that this means that $P_{G_{\mathcal{O}}}$ is generated by $q - n + r$ binomials, for every orientation $\mathcal{O}$ of G . The main technical result in this section is Theorem 2, using this theorem and Lemma 6 we prove that the CIO property is closed under induced subgraphs. Also we prove that a graph is CIO if and only if its blocks are CIO.

Recall that if $v \in V(G)$, the neighborhood of v is the set $N_G(v) = \{w \in V(G) \mid \{v, w\} \in E(G)\}$. We take $x \in V(G)$, $G' = G \setminus x$ and $\mathcal{O}'$ an orientation of G' . If $N_G(x) = \{x_1, \dots, x_s\}$, then we define the orientation $\mathcal{O}_x$ of G as follows: set $E(G_{\mathcal{O}_x}) = E(G'_{\mathcal{O}'}) \cup \{(x, x_1), \dots, (x, x_s)\}$. In the following three lemmas assume that $E(G_{\mathcal{O}_x}) = \{t_1, \dots, t_q\}$, where $t_i = (x, x_i)$ for $i = 1, \dots, s$. We denote the oriented graph $G_{\mathcal{O}_x}$ by D_x (i.e. $D_x = G_{\mathcal{O}_x}$).

Lemma 3 Let $t^\alpha - t^\beta$ be a binomial in P_{D_x} . If $\alpha = (\alpha_1, \dots, \alpha_q)$ and $\beta = (\beta_1, \dots, \beta_q)$, then $\alpha \cdot \underbrace{(1, \dots, 1, 0, \dots, 0)}_s = \beta \cdot \underbrace{(1, \dots, 1, 0, \dots, 0)}_s$, i.e., $\sum_{i=1}^s \alpha_i = \sum_{i=1}^s \beta_i$.

Proof We have $t_i = (x, x_i)$ for $i = 1, \dots, s$ thus, $\varphi(t_i) = x_i x^{-1}$. Hence, $\varphi(t^\alpha) = \frac{m_1}{x^{k_1 m_2}}$ and $\varphi(t^\beta) = \frac{n_1}{x^{k_2 n_2}}$, where m_1, m_2, n_1 and n_2 are monomials. Also set $\gcd(x, m_1) = \gcd(x, m_2) = \gcd(x, n_1) = \gcd(x, n_2) = 1$. Furthermore $k_1 = \sum_{i=1}^s \alpha_i$ and $k_2 = \sum_{i=1}^s \beta_i$. On the other hand $t^\alpha - t^\beta \in P_{D_x}$ then $\varphi(t^\beta) = \varphi(t^\alpha)$. Therefore, $\frac{m_1}{x^{k_1 m_2}} = \frac{n_1}{x^{k_2 n_2}}$ and $k_1 = \sum_{i=1}^s \alpha_i = \sum_{i=1}^s \beta_i = k_2$. $\square$

Corollary 3 Let $t^\alpha - t^\beta$ be a binomial in P_{D_x} . If there exists $a \in \{1, \dots, s\}$ such that $t_a \mid t^\alpha$ then there exists $b \in \{1, \dots, s\}$ such that $t_b \mid t^\beta$.

Lemma 4 Let $f = t^\alpha - t^\beta$ be a binomial in $P_{D_x} \cap k[t_{s+1}, \dots, t_q]$. If B is a generating set of binomials of P_{D_x} and $B' = B \cap k[t_{s+1}, \dots, t_q]$ then $f \in (B') \subset k[t_{s+1}, \dots, t_q]$.

Proof Set $B = \{g_\lambda\}_{\lambda \in \Omega}$. Since $P_{D_x} = (B)$, then $f = t^\alpha - t^\beta = \sum_{\lambda \in \Omega} f_\lambda g_\lambda$, where $f_\lambda \in k[t_1, \dots, t_q]$. We set in the equation above $t_1 = t_2 = \dots = t_s = 0$. By Corollary 3

$$g_\lambda|_{t_1=t_2=\dots=t_s=0} = \begin{cases} g_\lambda & \text{if } g_\lambda \in B' \\ 0 & \text{if } g_\lambda \notin B'. \end{cases}$$

Hence, we obtain $f = f|_{t_1=t_2=\dots=t_s=0} = \sum_{g_\lambda \in B'} f'_\lambda g_\lambda$, where $f'_\lambda = f_\lambda|_{t_1=t_2=\dots=t_s=0} \in k[t_{s+1}, \dots, t_q]$. Therefore $f \in (B')$. $\square$

Lemma 5 If $G' = G \setminus x$ is a connected graph then for every $j_1, j_2 \in \{1, \dots, s\}$ there exists a binomial $t_{j_1} m_{j_1} - t_{j_2} m_{j_2}$ in P_{D_x} with $m_{j_1}, m_{j_2} \in k[t_{s+1}, \dots, t_q]$.

Proof We know that $t_{j_1} = (x, x_{j_1})$ and $t_{j_2} = (x, x_{j_2})$. Since G' is connected, there exists a path $\mathcal{L}_{j_1, j_2}$ in $D' = G'_{\mathcal{O}'}$ between x_{j_1} and x_{j_2} . Thus, $C_{j_1, j_2} = (x, x_{j_1}, \mathcal{L}_{j_1, j_2}, x_{j_2}, x)$ is a cycle in D_x and $t_{C_{j_1, j_2}} = t_{j_1} m - t_{j_2} m'$ with $m, m' \in k[t_{s+1}, \dots, t_q]$. Therefore $t_{C_{j_1, j_2}} \in P_{D_x}$. $\square$

Theorem 2 Let G be a connected graph. If there exists a vertex x which is not a cut vertex of G such that $G' = G \setminus x$ is not a CIO graph then G is not a CIO graph.

Proof Set $|V(G)| = n$, $|E(G)| = q$ and $|E(G) \setminus E(G')| = s$. Since G' is not a CIO graph there exists an oriented graph $D' = G'_{\mathcal{O}'}$ associated to an orientation $\mathcal{O}'$ of G' such that $P_{D'}$ is not generated by $\text{ht}(P_{D'})$ binomials. We take the orientation $\mathcal{O}_x$ of G given by $E(G_{\mathcal{O}_x}) = E(G'_{\mathcal{O}'}) \cup \{(x, x_1), \dots, (x, x_s)\}$, where $N_G(x) = \{x_1, \dots, x_s\}$ and $D_x = G_{\mathcal{O}_x}$. We can suppose that $t_i = (x, x_i)$ for $1 \leq i \leq s$, i.e., $E(D_x) \setminus E(D') = \{t_1, \dots, t_s\}$. We take $\mathcal{G} = \{f_1, \dots, f_k\}$ a minimum generating set of binomials of P_{D_x} and we define $\mathcal{G}' = \mathcal{G} \cap k[t_{s+1}, \dots, t_q]$. By Lemma 4, $P_{D'} = (g \mid g \in \mathcal{G}')$ and

$\text{ht}(P_{D'}) = (q - s) - (n - 1) + 1$. Hence $|\mathcal{G}'| > (q - s) - (n - 1) + 1$. Now, we define the auxiliary graph $\mathcal{H}$ with $V(\mathcal{H}) = \{t_1, \dots, t_s\}$ and

$$E(\mathcal{H}) = \left\{ \{t_i, t_j\} \mid \begin{array}{l} \text{there exist monomials } m_i, m_j \in k[t_{s+1}, \dots, t_q] \text{ such that} \\ t_i m_i - t_j m_j \in \mathcal{G} \text{ or } t_j m_j - t_i m_i \in \mathcal{G} \text{ with } i, j \in \{1, \dots, s\} \end{array} \right\}.$$

If $A = \{t_i m_i - t_j m_j \in \mathcal{G} \mid \{t_i, t_j\} \in E(\mathcal{H})\}$ then $|A| = |E(\mathcal{H})|$. Also $A \cup \mathcal{G}' \subset \mathcal{G}$ and $A \cap \mathcal{G}' = \emptyset$. Hence we have $|E(\mathcal{H})| + |\mathcal{G}'| \leq |\mathcal{G}|$. Now, if $\mathcal{H}$ is connected, then $|E(\mathcal{H})| \geq |V(\mathcal{H})| - 1 = s - 1$. Furthermore, since $|\mathcal{G}'| > (q - s) - (n - 1) + 1$, then $|\mathcal{G}| > (s - 1) + (q - s) - (n - 1) + 1 = q - n + 1 = \text{ht}(P_{D_x})$. Thus, P_{D_x} is not generated by $q - n + 1$ binomials.

Thus, to obtain the result it is only necessary to prove that $\mathcal{H}$ is connected. Suppose, by way of contradiction, that the connected components of $\mathcal{H}$ are $\mathcal{H}_1, \dots, \mathcal{H}_p$ with $p > 1$. By Lemma 5 there exists a binomial $t_1 m_1 - t_j m_j \in P_{D_x}$ such that $t_1 \in V(\mathcal{H}_1)$ and $t_j \in V(\mathcal{H}_2)$, where m_1, m_j are monomials in $k[t_{s+1}, \dots, t_q]$ and $j \in \{2, \dots, s\}$. Recall that each $f_i \in \mathcal{G}$ has the form $f_i = t^{\alpha_i} - t^{\beta_i}$, where $t^{\alpha_i}, t^{\beta_i}$ are monomials in $k[t_1, \dots, t_q]$, then

$$t_1 m_1 - t_j m_j = \sum_{i=1}^k g_i (t^{\alpha_i} - t^{\beta_i}), \quad \text{with } g_i \in k[t_1, \dots, t_q].$$

If $\alpha_i = (\alpha_1^i, \dots, \alpha_q^i)$ and $\beta_i = (\beta_1^i, \dots, \beta_q^i)$, we take $\alpha'_i = (\alpha_1^i, \dots, \alpha_s^i, 0, \dots, 0)$ and $\beta'_i = (\beta_1^i, \dots, \beta_s^i, 0, \dots, 0)$; observe by Lemma 3, $\deg(t^{\alpha'_i}) = \deg(t^{\beta'_i})$ for $i \in \{1, \dots, k\}$. Moreover, $\alpha'_i = 0 \Leftrightarrow \beta'_i = 0$, and in this case, we have $(t^{\alpha'_i} - t^{\beta'_i})|_{t_{s+1}=\dots=t_q=1} = 0$. Hence, if we set in the equation above $t_{s+1} = \dots = t_q = 1$, we obtain

$$t_1 - t_j = \sum_{\deg(t^{\alpha'_i})=1} g'_i (t^{\alpha'_i} - t^{\beta'_i}) + \sum_{\deg(t^{\alpha'_i})>1} g'_i (t^{\alpha'_i} - t^{\beta'_i}),$$

where $g'_i = g_i|_{t_{s+1}=\dots=t_q=1}$. Furthermore, we can write $g'_i = a_i + \lambda_i$ with $a_i \in k[t_1, \dots, t_s]$ and $\lambda_i \in k$, such that if $a_i \neq 0$ then $\deg(a_i) > 0$. Thus,

$$t_1 - t_j = \sum_{\deg(t^{\alpha'_i})=1} \lambda_i (t^{\alpha'_i} - t^{\beta'_i}) + g',$$

where $g' = \sum_{\deg(t^{\alpha'_i})=1} a_i (t^{\alpha'_i} - t^{\beta'_i}) + \sum_{\deg(t^{\alpha'_i})>1} g'_i (t^{\alpha'_i} - t^{\beta'_i})$. Hence we have either $\deg(g') > 1$ or $g' = 0$. If we suppose $\deg(g') > 1$, then

$$t_1 - t_j - \sum_{\deg(t^{\alpha'_i})=1} \lambda_i (t^{\alpha'_i} - t^{\beta'_i}) = g',$$

however, the left side has degree at most one, but $\deg(g') > 1$, which is a contradiction. Therefore $g' = 0$, so

$$\begin{aligned} t_1 - t_j &= \sum_{\deg(t^{\alpha'_i})=1} \lambda_i (t^{\alpha'_i} - t^{\beta'_i}) \\ &= \sum_{\{t_{i_1}, t_{i_2}\} \in E(\mathcal{H}_1)} \lambda_{i_1, i_2}^1 (t_{i_1} - t_{i_2}) + \cdots + \sum_{\{t_{i_1}, t_{i_2}\} \in E(\mathcal{H}_p)} \lambda_{i_1, i_2}^p (t_{i_1} - t_{i_2}), \end{aligned}$$

where $\lambda_i = \lambda_{i_1, i_2}^u$ if $t^{\alpha'_i} = t_{i_1}$, $t^{\beta'_i} = t_{i_2}$ and $\{t_{i_1}, t_{i_2}\} \in E(\mathcal{H}_u)$. In the last equation we set $t_\ell = 1$ if $t_\ell \in V(\mathcal{H}_1)$ and $t_\ell = 0$ otherwise, and we obtain

$$\begin{aligned} 1 - 0 &= \sum_{\{t_{i_1}, t_{i_2}\} \in E(\mathcal{H}_1)} \lambda_{i_1, i_2}^1 (1 - 1) + \sum_{\{t_{i_1}, t_{i_2}\} \in E(\mathcal{H}_2)} \lambda_{i_1, i_2}^2 (0 - 0) + \cdots \\ &+ \sum_{\{t_{i_1}, t_{i_2}\} \in E(\mathcal{H}_p)} \lambda_{i_1, i_2}^p (0 - 0) = 0 \end{aligned}$$

since $t_1 \in V(\mathcal{H}_1)$ and $t_j \in V(\mathcal{H}_2)$. This is a contradiction. Therefore $\mathcal{H}$ is connected. $\square$

Definition 3 Let G be a graph. A subgraph H of G is called an induced subgraph of G if $E(H) = \{e \in E(G) \mid e \subseteq V(H)\}$. Furthermore if $X \subseteq V(G)$ then we denote by $[X]_G$ the induced subgraph H of G such that $V(H) = X$, and in this case we say that H is induced by X .

Lemma 6 Let G be a connected graph. If G' is a proper induced connected subgraph of G , then there exists $x \in V(G) \setminus V(G')$ such that $G \setminus x$ is connected.

Proof Set $|V(G)| = n$. If $v \in V(G) \setminus V(G')$, there exists G_v , a connected component of $G \setminus v$ such that $G' \subset G_v$. We define $r_v = |V(G_v)|$. Let x be the vertex in $V(G) \setminus V(G')$ with r_x maximal. Let $G_1, \dots, G_k$ be the connected components of $G \setminus x$ with $G' \subset G_1$, i.e., $G_v = G_1$. Hence, $|V(G_1)| = r_x$. Suppose that $r_x < n - 1$, then $k > 1$. Thus, there is $y \in V(G_2) \subset V(G) \setminus V(G')$. Since G is connected, we have x is adjacent to a vertex of G_1 . Hence $[V(G_1) \cup x]_G$ is contained in a connected component of $G \setminus y$. Then $r_y > r_x$, this is a contradiction. Therefore $r_x = n - 1$, $k = 1$ and $G \setminus x$ is connected. $\square$

Lemma 7 Let G_1, G_2 be CIO subgraphs of G , such that $G = G_1 \cup G_2$. If $|V(G_1) \cap V(G_2)| \leq 1$ then G is a CIO graph.

Proof Let $D = G_{\mathcal{O}}$ be the oriented graph associated to an orientation $\mathcal{O}$ of G and $\mathcal{O}_i$ the orientation of G_i induced by $\mathcal{O}$. We set $D_i = (G_i)_{\mathcal{O}_i}$. Then for $i = 1, 2$ there exists a binomial generating set $\mathcal{G}_i$ of P_{D_i} such that $|\mathcal{G}_i| = q_i - n_i + r_i$ where $q_i = |E(G_i)|$, $n_i = |V(G_i)|$ and r_i is the number of the connected components of G_i . Since $|V(D_1) \cap V(D_2)| \leq 1$, we see that if C is a cycle of D then either $C \subset D_1$ or

$C \subset D_2$. Hence, $t_C \in (\mathcal{G}_1 \cup \mathcal{G}_2)$. So $\mathcal{G} = \mathcal{G}_1 \cup \mathcal{G}_2$ is a binomial generating set of P_D . Since $E(G_1) \cap E(G_2) = \emptyset$ then $\mathcal{G}_1 \cap \mathcal{G}_2 = \emptyset$ and

$$|\mathcal{G}| = |\mathcal{G}_1| + |\mathcal{G}_2| = (q_1 + q_2) - (n_1 + n_2) + (r_1 + r_2).$$

On the other hand $|E(G)| = q_1 + q_2$ and the cardinality of $V(G)$ is

$$|V(G)| = \begin{cases} n_1 + n_2 & \text{if } V(D_1) \cap V(D_2) = \emptyset \\ n_1 + n_2 - 1 & \text{if } |V(D_1) \cap V(D_2)| = 1 \end{cases}$$

but if r is the number of components of G then $r = r_1 + r_2$ in the first case and $r = r_1 + r_2 - 1$ in the second case. Therefore $|\mathcal{G}| = |E(G)| - |V(G)| + r = \text{ht}(P_D)$ and G is a CIO graph. $\square$

Proposition 5 *Let G be a graph. G is a CIO graph if and only if its connected components are CIO graphs.*

Proof ($\Rightarrow$) By induction on $|V(G)|$. Let $G_1, \dots, G_k$ be the connected components of G . We can suppose that $k > 1$. By Lemma 6 there exists $x_i \in V(G_i)$ such that $G_i \setminus x_i$ is connected, for $i = 1, \dots, k$. Hence by the induction hypothesis the connected components of $G \setminus x_i$ are CIO graphs. Since this is true for every $i = 1, \dots, k$ then $G_1, \dots, G_k$ are CIO graphs.

($\Leftarrow$) It is a consequence of Lemma 7. $\square$

Corollary 4 *If G is a CIO graph then every induced subgraph H of G is a CIO graph.*

Proof Assume H is not a CIO graph. By Proposition 5, H has a connected component H_1 such that H_1 is not a CIO graph. Let G_1 be the connected component of G such that $H_1 \subset G_1$. If $G_1 \neq H_1$, then by Lemma 6 there exists $x_1 \in V(G_1) \setminus V(H_1)$ such that $G_1^1 = G_1 \setminus x_1$ is connected. Furthermore H_1 is an induced subgraph of G_1^1 thus, if $G_1^1 \neq H_1$, then there exists $x_2 \in V(G_1^1) \setminus V(H_1)$, such that $G_1^2 = G_1^1 \setminus x_2$ is a connected subgraph and $H_1 \subset G_1^2$. We continue with this process until we obtain the connected subgraphs $G_1 = G_1^0, G_1^1, \dots, G_1^k = H_1$ such that $G_1^{i+1} = G_1^i \setminus x_{i+1}$. Since H_1 is not a CIO graph by Theorem 2 $H_1 = G_1^k, \dots, G_1^1, G_1^0 = G_1$ are not CIO graphs. Therefore by Proposition 5, G is not a CIO graph. This is a contradiction, therefore H is a CIO graph. $\square$

Proposition 6 *G is a CIO graph if and only if every block B of G is a CIO graph.*

Proof ($\Rightarrow$) If B is a block of G then B is an induced subgraph of G . Hence by Theorem 4, B is a CIO graph.

($\Leftarrow$) It is a consequence of Lemma 7. $\square$

4 Chorded-theta graphs

In this section we introduce chorded-theta graphs, the notion of transversal triangles in these graphs and theta-ring graphs. In Theorem 3 we describe the minimal chorded-theta graphs without transversal triangles. In this section we prove that theta-ring graphs are closed under 0, 1, 2-clique-sums and that chordal graphs are theta-ring graphs. With these results we obtain the main result of this section: theta-ring graphs can be constructed by clique-sums of complete graphs and/or cycles, or equivalently by 0, 1, 2-clique-sums of chordal graphs and/or cycles. Finally we prove that the minimal forbidden induced subgraphs for the characterization of theta-ring graphs are: prisms, pyramids, thetas and partial wheels that are chorded-thetas. This section is independent of the rest of the paper and its techniques are purely combinatorial.

Definition 4 A chorded-theta subgraph T of G is a subgraph induced by three paths $\mathcal{L}_1, \mathcal{L}_2, \mathcal{L}_3$ each between the non adjacent vertices x and y such that $V(\mathcal{L}_i) \cap V(\mathcal{L}_j) = \{x, y\}$ for $1 \leq i < j \leq 3$. The edges that do not belong to any of the sets $E(\mathcal{L}_1), E(\mathcal{L}_2)$ and $E(\mathcal{L}_3)$ are called the chords of T . Furthermore if each chord of T has its end vertices in different paths (of $\mathcal{L}_1, \mathcal{L}_2, \mathcal{L}_3$), then T is called a simple chorded-theta subgraph.

To emphasize that $\mathcal{L}_1, \mathcal{L}_2$ and $\mathcal{L}_3$ are the three paths associated to T we denote them by $\mathcal{L}_i(T)$ for $i = 1, 2, 3$. On the other hand we denote by $\mathcal{L}_i^o$ the interior of $\mathcal{L}_i$, i.e., $\mathcal{L}_i^o = \mathcal{L}_i \setminus \{x, y\}$.

Definition 5 Let T be a chorded-theta of G , with $\mathcal{L}_1(T) = (x, x_1, \dots, x_{r_1}, y)$, $\mathcal{L}_2(T) = (x, y_1, \dots, y_{r_2}, y)$ and $\mathcal{L}_3(T) = (x, z_1, \dots, z_{r_3}, y)$. A transversal triangle H of T is a triangle in G such that $V(H) = \{x_i, y_j, z_k\}$ for some i, j, k .

We define the following sets: $\Theta(G) = \{T \mid T \text{ is a chorded-theta of } G\}$, $\Theta_\Delta(G) = \{T \in \Theta(G) \mid T \text{ has a transversal triangle}\}$ and $\Theta_\Delta^c(G) = \Theta(G) \setminus \Theta_\Delta(G)$. Moreover, a minimum chorded-theta subgraph T without transversal triangles is a subgraph in $\Theta_\Delta^c(G)$ such that if $T' \in \Theta_\Delta^c(G)$ then $|V(T)| \leq |V(T')|$.

Remark 1 If T is a minimum chorded-theta subgraph without transversal triangles, then T is a simple chorded-theta subgraph.

In the following two lemmas we assume that T is a chorded-theta subgraph of G with $\mathcal{L}_1(T) = (x, x_1, \dots, x_{r_1}, y)$, $\mathcal{L}_2(T) = (x, y_1, \dots, y_{r_2}, y)$ and $\mathcal{L}_3(T) = (x, z_1, \dots, z_{r_3}, y)$, where $|V(T)| = r_1 + r_2 + r_3 + 2$.

Lemma 8 Let T be a minimum chorded-theta subgraph without transversal triangles with $r_1 \leq r_2$. If $\{x_i, y_j\}$ is a chord of T , then $(i, j) \in \{(1, 1), (r_1, r_2)\}$ or $r_1 = 1$.

Proof By Remark 1, T is a simple chorded-theta. Let $\{x_i, y_j\}$ be a chord of T . We can assume that $r_1 > 1$. Hence, $i < r_1$ or $1 < i$.

Now, suppose $i < r_1$ and take the chorded-theta T_1 with $\mathcal{L}_1(T_1) = (x_i, \dots, x_{r_1}, y)$, $\mathcal{L}_2(T_1) = (x_i, y_j, \dots, y_{r_2}, y)$ and $\mathcal{L}_3(T_1) = (x_i, x_{i-1}, \dots, x_1, x, z_1, \dots, z_{r_3}, y)$. Since

T has no chords whose both end vertices are in $\mathcal{L}_1(T)$ and $T \in \Theta_{\Delta}^c(G)$, then $T_1 \in \Theta_{\Delta}^c(G)$. Observe that $|V(T_1)| = |V(T)| - (j - 1)$, so by the minimality of $|V(T)|$, we have $j = 1$. We will prove that $i = 1$, by contradiction we will assume $i > 1$. Thus, we can take the chorded-theta T_2 with $\mathcal{L}_1(T_2) = (x, y_1, x_i)$, $\mathcal{L}_2(T_2) = (x, x_1, \dots, x_i)$ and $\mathcal{L}_3(T_2) = (x, z_1, \dots, z_{r_3}, y, x_{r_1}, \dots, x_i)$. Again, since T has no chords whose both end vertices are in $\mathcal{L}_1(T)$ and $T \in \Theta_{\Delta}^c(G)$, we have $T_2 \in \Theta_{\Delta}^c(G)$. But $|V(T_2)| = |V(T)| - (r_2 - 1)$. Furthermore $1 < r_1 \leq r_2$ so $|V(T)| > |V(T_2)|$ which is a contradiction. Therefore $i = 1$ and $(i, j) = (1, 1)$.

Similarly, if $1 < i$ we obtain $(i, j) = (r_1, r_2)$. $\square$

Lemma 9 *Let T be a minimum chorded-theta without transversal triangles such that $r_1 = 1$ and $r_2 \leq r_3$. If there exist chords e_1, e_2 of T with $e_1 = \{x_1, z_{s_1}\}$ and $e_2 = \{y_j, z_{s_2}\}$ then $\{x_1, y_j\} \notin E(G)$. Moreover $(s_1, j, s_2) = (r_3, 1, 1)$ or $(s_1, j, s_2) = (1, r_2, r_3)$.*

Proof By Remark 1, T is a simple chorded-theta.

Case (1) $s_1 \leq s_2$. Suppose that $\{x_1, y_j\} \in E(G)$. Since $T \in \Theta_{\Delta}^c(G)$ then $\{x_1, z_{s_2}\} \notin E(G)$ and $z_{s_1} \neq z_{s_2}$. Thus, we take a chorded-theta T' given by $\mathcal{L}_1(T') = (z_{s_2}, y_j, x_1)$, $\mathcal{L}_2(T') = (z_{s_2}, \dots, z_{r_3}, y, x_1)$ and $\mathcal{L}_3(T') = (z_{s_2}, z_{s_2-1}, \dots, z_{s_1}, x_1)$. Furthermore there is no chord whose both end vertices are in $\mathcal{L}_3(T)$, hence $T' \in \Theta_{\Delta}^c(G)$. But $|V(T')| = |V(T)| - s_1 - j + 1 < |V(T)|$ which is a contradiction. Therefore $\{x_1, y_j\} \notin E(G)$. Now, we consider the chorded-theta T'' given by $\mathcal{L}_1(T'') = (x_1, z_{s_1}, z_{s_1+1}, \dots, z_{s_2}, y_j)$, $\mathcal{L}_2(T'') = (x_1, x, y_1, \dots, y_j)$ and $\mathcal{L}_3(T'') = (x_1, y, y_{r_2}, y_{r_2-1}, \dots, y_j)$. We have $T'' \in \Theta_{\Delta}^c(G)$ because there is no chord with its both end vertices in $\mathcal{L}_2(T)$. But $|V(T'')| = |V(T)| - (s_1 - 1) - (r_3 - s_2)$, then $s_1 = 1$ and $r_3 = s_2$ because T is minimum. If $j \neq r_2$ we can take a chorded-theta subgraph T''' , where $\mathcal{L}_1(T''') = (y_j, \dots, y_{r_2}, y)$, $\mathcal{L}_2(T''') = (y_j, \dots, y_1, x, x_1, y)$ and $\mathcal{L}_3(T''') = (y_j, z_{r_3}, y)$. Thus $T''' \in \Theta_{\Delta}^c(G)$ because T is a simple chorded-theta together with the fact that $T \in \Theta_{\Delta}^c(G)$. But $|V(T''')| = |V(T)| - (r_3 - 1)$ hence $r_3 = 1$. Since $r_2 \leq r_3$ then $r_2 = 1 = j$, which is a contradiction. Therefore $j = r_2$ and $(s_1, j, s_2) = (1, r_2, r_3)$.

Case (2) $s_2 \leq s_1$. Similarly to the case (1). We obtain $\{x_1, y_j\} \notin E(G)$ and $(s_1, j, s_2) = (r_3, 1, 1)$. $\square$

Definition 6 ($\forall\theta\exists\Delta$ -property) A graph G is called a theta-ring graph if every chorded-theta of G has a transversal triangle, in this case we say that G has the $\forall\theta\exists\Delta$ -property.

Remark 2 G is a theta-ring graph if and only if $\Theta_{\Delta}^c(G) = \emptyset$.

Theorem 3 *Let G be a graph. If G is not a theta-ring graph then there exists a chorded-theta T in G where $\mathcal{L}_1(T) = (x, x_1, \dots, x_{r_1}, y)$, $\mathcal{L}_2(T) = (x, y_1, \dots, y_{r_2}, y)$, $\mathcal{L}_3(T) = (x, z_1, \dots, z_{r_3}, y)$, with $1 \leq r_1 \leq r_2 \leq r_3$ such that T satisfies at least one of the following conditions:*

- (1) *The chords of T are contained in $\{\{y_1, z_1\}, \{y_{r_2}, z_{r_3}\}\}$.*
- (2) *$r_1 = 1$, and the chords of T are contained in $\{\{x_1, z_{r_3}\}, \{y_1, z_1\}\}$.*

- (3) $r_1 = 1$, and the chords of T have the form $\{x_1, z_j\}$ for some j .
 (4) $r_1 = 1$, $r_2 = 1$ and the chords of T have the form $\{x_1, y_1\}$ or $\{x_1, z_j\}$ for some j .

Proof Since $\Theta_\Delta^c(G) \neq \emptyset$. We can take a minimum chorded-theta T without transversal triangles. Thus, T is a simple chorded-theta. Let $\mathcal{L}_1(T) = (x, x_1, \dots, x_{r_1}, y)$, $\mathcal{L}_2(T) = (x, y_1, \dots, y_{r_2}, y)$ and $\mathcal{L}_3(T) = (x, z_1, \dots, z_{r_3}, y)$ with $1 \leq r_1 \leq r_2 \leq r_3$. We can suppose that $|\mathcal{L}_1(T)| \leq |\mathcal{L}_1(T')|$ for all $T' \in \Theta_\Delta^c(G)$ such that $|V(T')|$ is minimum. Also we assume that if $|\mathcal{L}_1(T)| = |\mathcal{L}_1(T')|$ then $|\mathcal{L}_2(T)| \leq |\mathcal{L}_2(T')|$.

Now, we suppose that T does not satisfy property (1). Hence, there exists a chord of T in one of the possible forms: $\{x_i, y_j\}$ or $\{x_i, z_j\}$ or $\{y_i, z_k\}$, where $(i, k) \notin \{(1, 1), (r_2, r_3)\}$.

In the first case we suppose that there exists $e = \{y_i, z_k\}$, a chord of T , where $(i, k) \notin \{(1, 1), (r_2, r_3)\}$. By Lemma 8 we have $r_2 = 1$ and $i = 1$. Now, since $r_1 \leq r_2$ then $r_1 = 1$. So, by Lemma 9, if there exists $\{x_1, z_j\} \in E(G)$ then $\{x_1, y_1\} \notin E(G)$ and $(j, i, k) = (r_3, 1, 1)$ or $(j, i, k) = (1, 1, r_3)$. Thus, we would be in case (2). Now, if $\{x_1, z_j\} \notin E(G)$ for every j then the chords of T have the form $\{x_1, y_1\}$ or $\{y_1, z_j\}$, and we are in case (4).

In the second case suppose there exists a chord e of T such that $e = \{x_i, y_j\}$ or $e = \{x_i, z_j\}$. We will prove that $r_1 = 1$. We argue by contradiction, assume that $r_1 > 1$ and $e = \{x_i, y_j\}$. By Lemma 8, we have $(i, j) \in \{(1, 1), (r_1, r_2)\}$. We can suppose that $(i, j) = (r_1, r_2)$ and take the chorded-theta graph T_1 where $\mathcal{L}_1(T_1) = (x, x_1, \dots, x_{r_1})$, $\mathcal{L}_2(T_1) = (x, y_1, \dots, y_{r_2}, x_{r_1})$ and $\mathcal{L}_3(T_1) = (x, z_1, \dots, z_{r_3}, y, x_{r_1})$. Since T is a simple chorded-theta and $T \in \Theta_\Delta^c(G)$ then $T_1 \in \Theta_\Delta^c(G)$. But $|V(T)| = |V(T_1)|$ and $|\mathcal{L}_1(T_1)| = |\mathcal{L}_1(T)| - 1$, which is a contradiction. The same result is obtained if $e = \{x_i, z_j\}$. Hence, we obtain $r_1 = 1$.

Now, we will prove that if there exists a chord e' of T such that $e' = \{x_1, y_u\}$ then T satisfies (4). First we will prove that $r_2 = 1$. Again we argue by contradiction, assume $r_2 > 1$. Hence, $u > 1$ or $u < r_2$. Without loss of generality, take $u > 1$ and consider the chorded-theta T_2 where $\mathcal{L}_1(T_2) = (x, x_1, y_u)$, $\mathcal{L}_2(T_2) = (x, y_1, \dots, y_u)$ and $\mathcal{L}_3(T_2) = (x, z_1, \dots, z_{r_3}, y, y_{r_2}, \dots, y_u)$. Since $T \in \Theta_\Delta^c(G)$ and there are no chords in $\mathcal{L}_2(T_2)$, then $T_2 \in \Theta_\Delta^c(G)$. But $|V(T)| = |V(T_2)|$, $|\mathcal{L}_1(T)| = |\mathcal{L}_1(T_2)|$ and $|\mathcal{L}_2(T)| > |\mathcal{L}_2(T_2)|$, which is a contradiction. Thus, $r_1 = r_2 = 1$ and $e' = \{x_1, y_1\}$. Suppose T does not satisfy (4) then there must exist $e_1, e_2 \in E(G)$ such that $e_1 = \{x_1, z_{j_1}\}$ and $e_2 = \{y_1, z_{j_2}\}$. But, by Lemma 9, $\{x_1, y_1\} \notin E(G)$. This is a contradiction. Hence, T satisfies (4).

Now, we can assume that T does not have a chord of the form $e' = \{x_1, y_u\}$ then $e = \{x_1, z_j\}$. Suppose T does not satisfy (3), then there must exist $e_3 = \{y_i, z_k\}$. By Lemma 9, $(j, i, k) \in \{(r_3, 1, 1), (1, r_2, r_3)\}$. Therefore T satisfies (2). $\square$

Definition 7 A partial wheel W is a graph where $V(W) = \{z, z_1, \dots, z_k\}$ such that $C = (z_1, \dots, z_k, z_1)$ is a cycle in W and the edges of W are the edges of C and some edges between z and vertices of C . C is called the rim of W and z is called the center of W .

A prism is a graph consisting of two vertex-disjoint triangles $C_1 = (x_1, x_2, x_3, x_1)$ and $C_2 = (y_1, y_2, y_3, y_1)$, and three paths L_1, L_2, L_3 pairwise vertex-disjoint, such that each L_i is a path between x_i and y_i for $i = 1, 2, 3$ and the subgraph induced by $V(L_i) \cup V(L_j)$ is a cycle for $1 \leq i < j \leq 3$.

A pyramid is a graph consisting of a vertex w , a triangle $C = (z_1, z_2, z_3, z_1)$, and three paths P_1, P_2, P_3 , such that: P_i is between w and z_i for $i = 1, 2, 3$; $V(P_i) \cap V(P_j) = \{w\}$ and the subgraph induced by $V(P_i) \cup V(P_j)$ is a cycle for $1 \leq i < j \leq 3$; and at most one of the P_1, P_2, P_3 has only one edge.

A theta is a graph consisting of two non adjacent vertices x and y , and three paths P, Q, R with ends x and y , such that the union of every two of P, Q, R is an induced cycle.

Detecting prisms, pyramids and thetas has been widely studied, for example see [2, 10] and their references.

Definition 8 A partial wheel T with rim C and center z is called a θ -partial wheel if $|V(C)| \geq 4$ and there exist two non adjacent vertices in $V(C) \cap N_T(z)$.

Remark 3 A theta is a chorded-theta without chords. Furthermore, if W is a partial wheel then W is a θ -partial wheel if and only if W is a chorded-theta.

Corollary 5 T is a minimal forbidden induced subgraph for the class of theta-ring graphs if and only if T satisfies at least one of the following conditions:

- (a) T is a θ -partial wheel.
- (b) T is a prism, pyramid or theta.

Proof First we will prove that if T satisfies (a) or (b), then T is a chorded-theta graph without transversal triangles and T does not contain a proper chorded-theta subgraph.

If T satisfies (a) with rim $C = (z_1, z_2, \dots, z_{r_1}, z_1)$ and center z , then there exist $z_i, z_j \in N_T(z)$ where $\{z_i, z_j\} \notin E(T)$. We can suppose that $i < j$, then T is a chorded-theta graph with the paths $\mathcal{L}_1(T) = (z_i, z_{i+1}, \dots, z_j)$, $\mathcal{L}_2(T) = (z_i, z, z_j)$ and $\mathcal{L}_3(T) = (z_i, z_{i-1}, \dots, z_1, z_{r_1}, z_{r_1-1}, \dots, z_j)$. Furthermore, since C does not contain chords then T is a chorded-theta without transversal triangles. On the other hand if T' is a chorded-theta subgraph of T , then one terminal vertex is in C . Thus, one path of T' has the form (z_{i_1}, z, z_{j_1}) with $\{z_{i_1}, z_{j_1}\} \notin E(T)$ and $i_1 < j_1$. Hence, the others paths of T' are $(z_{i_1}, z_{i_1+1}, \dots, z_{j_1})$ and $(z_{i_1}, z_{i_1-1}, \dots, z_1, z_{r_1}, z_{r_1-1}, \dots, z_{j_1})$. Therefore $V(T) = V(T')$ and $T = T'$.

Now, if T is a prism, then T is a chorded-theta with $\mathcal{L}_1(T) = (x_1, x_2, L_2, y_2)$, $\mathcal{L}_2(T) = (x_1, L_1, y_1, y_2)$ and $\mathcal{L}_3(T) = (x_1, x_3, L_3, y_3, y_2)$. Furthermore T has only two chords, then T does not contain transversal triangles. On the other hand, if T' is a chorded-theta in T , then the terminal vertices are x_i and y_j for some $i, j \in \{1, 2, 3\}$. Since x_i and y_j have degree at least three in T' , we have $V(T') = V(T)$ and $T' = T$.

If T is a pyramid, then we can suppose that $|E(P_1)| \geq 2$ and T is a chorded-theta with $\mathcal{L}_1(T) = (z_1, P_1, w)$, $\mathcal{L}_2(T) = (z_1, z_2, P_2, w)$ and $\mathcal{L}_3(T) = (z_1, z_3, P_3, w)$. Additionally if T has only one chord then T does not contain transversal triangles. Furthermore if T' is a chorded-theta in T then the terminal vertices are z_i and w for some $i \in \{1, 2, 3\}$. Since z_i and w have degree at least 3 in T' , we have $V(T') = V(T)$ and $T' = T$.

Finally, if T is a theta graph it is clear that T is a minimal chorded-theta graph without transversal triangles.

Now, to obtain the result it is only necessary to prove that if T satisfies (1), (2), (3) or (4) of Theorem 3 then T satisfies (a) or (b).

Case 1: T satisfies (1). If T has less than two chords and T is neither a pyramid nor a theta, then $r_3 = 1$ and $r_2 = r_1 = 1$. Thus, T satisfies (a). Now, we can suppose that T has exactly two chords. If $r_2 \geq 2$ then T is a prism. Furthermore if $r_2 = 1$ then T satisfies (a).

Case 2: T satisfies (2). If T has less than two chords, and T is neither a pyramid nor a theta then $r_3 = 1$ and $r_2 = 1$. Hence, T satisfies (a). Now, we can suppose that T has exactly two chords. If $r_3 \geq 2$ then T is a prism. Furthermore if $r_3 = 1$ then $r_2 = r_1 = 1$ and T satisfies (a).

If T satisfies (3) or (4) then T satisfies (a). $\square$

Definition 9 Let G_1, G_2 be graphs such that $G = G_1 \cup G_2$, $K = G_1 \cap G_2$. If K is a complete graph with $|K| = k$, then G is called the k -clique-sum (or clique-sum) of G_1 and G_2 in K .

Remark 4 Let r_i be the number of connected components of G_i for $i = 1, 2$. If G is 0-clique-sum of G_1 and G_2 then the number of connected components of G is $r_1 + r_2$.

In our definition of clique-sum we do not allow to delete any of the edges of K .

Proposition 7 The clique-sum of theta-ring graphs is a theta-ring graph.

Proof Let G_1 and G_2 be theta-ring graphs. We assume that G is a k -clique-sum of G_1 and G_2 . We can suppose that $G = G_1 \cup G_2$ (i.e., $G_1, G_2 \subseteq G$) and $G_1 \cap G_2 = K$ where K is a k -complete subgraph of G . By contradiction suppose that G is not a theta-ring graph. Hence, there exists a chorded-theta H (of G) without transversal triangles, where $\mathcal{L}_1(H) = (x, x_1, \dots, x_{r_1}, y)$, $\mathcal{L}_2(H) = (x, y_1, \dots, y_{r_2}, y)$ and $\mathcal{L}_3(H) = (x, z_1, \dots, z_{r_3}, y)$. We can assume that H is a simple chorded-theta. Since $\{x, y\} \notin E(G)$ then $\{x, y\} \cap (G_1 \Delta G_2) \neq \emptyset$, where $G_1 \Delta G_2$ is the symmetric difference. Thus, without loss of generality we can assume that $x \in V(G_1) \setminus V(G_2)$. If $y \in V(G_2) \setminus V(G_1)$ then $\mathcal{L}_i^o(H) \cap K \neq \emptyset$ for $i = 1, 2, 3$. But K is a complete graph therefore H has a transversal triangle, which is not possible so $y \in V(G_1)$. On the other hand G_1 is a theta-ring graph, then $V(H)$ is not contained in $V(G_1)$. Hence, there exists $w \in V(H)$ such that $w \in V(G_2) \setminus V(G_1)$. We can suppose that $w \in \mathcal{L}_1^o(H)$, i.e., $w = x_i$ for some $1 \leq i \leq r_1$. Since H is a simple chorded-theta, then $L_1 = (x, x_1, \dots, x_i = w)$ is a path between $G_1 \setminus G_2$ and $G_2 \setminus G_1$. Furthermore K is a cutset and there exists $x_{i_1} \in V(L_1) \cap V(K)$ with $i_1 < i$. On the other hand if $y \in V(K)$ then $\{y, x_{i_1}\} \in E(G)$ but this is not possible because H is a simple chorded-theta. Thus, $y \in V(G_1) \setminus V(G_2)$ and $L_2 = (w = x_i, x_{i+1}, \dots, x_{r_1}, y)$ is a path between $G_1 \setminus G_2$ and $G_2 \setminus G_1$. Hence, there exists $x_{i_2} \in V(L_2) \cap V(K)$ with $i < i_2$. Since $x_{i_1}, x_{i_2} \in V(K)$ then $\{x_{i_1}, x_{i_2}\} \in E(G)$. This is a contradiction, because H is a simple chorded-theta. $\square$

Lemma 10 Let G be a 2-connected graph and let H be an induced connected subgraph of G with $|V(H)| \geq 2$. If $x \in V(G) \setminus V(H)$ then there exist paths L_1 and L_2

between x and H such that $V(L_1) \cap V(L_2) = \{x\}$, $V(L_i) \cap V(H) = \{a_i\}$ for $i = 1, 2$, and $a_1 \neq a_2$.

Proof Let y be a vertex in H . Since G is 2-connected graph there exist paths $P_1 = (x = x_1, x_2, \dots, x_{s_1} = y)$ and $P_2 = (x = y_1, y_2, \dots, y_{s_2} = y)$ between x and y , such that $V(P_1) \cap V(P_2) = \{x, y\}$. We define $j_1 = \min\{k \mid x_k \in V(H)\}$ and $j_2 = \min\{k \mid y_k \in V(H)\}$. If $x_{j_1} \neq y_{j_2}$ we can take $L_1 = (x = x_1, x_2, \dots, x_{j_1})$ and $L_2 = (x = y_1, y_2, \dots, y_{j_2})$. Hence, we can assume that $x_{j_1} = y_{j_2} = y$, but $V(P_1) \cap V(P_2) = \{x, y\}$ therefore $x_{j_1} = y_{j_2} = y$. Since $|V(H)| \geq 2$ there must exist $a \in V(H) \setminus \{y\}$. Furthermore G is 2-connected so $G \setminus y$ is connected, and there exists a path $P_3 = (x = z_1, \dots, z_{s_3} = a)$ between x and a in $G \setminus y$. Now, we take $k_1 = \max\{k \mid z_k \in V(P_1) \cup V(P_2)\}$, $k_2 = \min\{k \mid k \geq k_1 \text{ and } z_k \in V(H)\}$. We can suppose that $z_{k_1} \in V(P_1)$ and $z_{k_1} = x_{i_1}$. Finally we take $L_1 = (x = x_1, \dots, x_{i_1} = z_{k_1}, z_{k_1+1}, \dots, z_{k_2})$ and $L_2 = P_2$. $\square$

Definition 10 Let C be a cycle without chords. C is called a hole if C is not a triangle.

Lemma 11 Let G be a theta-ring graph with a hole C . Let x be a vertex of G such that $x \notin V(C)$. If there exist paths P_1 and P_2 between x and C such that $V(P_1) \cap V(P_2) = \{x\}$ with $V(P_i) \cap V(C) = \{y_i\}$ for $i = 1, 2$, then $\{y_1, y_2\} \in E(C)$. Furthermore if P is a path between x and C with $V(P) \cap V(C) = \{y\}$ then $y \in \{y_1, y_2\}$.

Proof Set $C = (x_1, x_2, \dots, x_k = x_1)$ where $x_1 = y_1$ and $y_2 = x_s$. By way of contradiction, suppose that $\{y_1, y_2\} \notin E(C)$. Thus, there exists a chorded-theta H of G with $\mathcal{L}_1(H) = (y_1, P_1, x, P_2, y_2)$, $\mathcal{L}_2(H) = (y_1 = x_1, x_2, \dots, x_s = y_2)$, and $\mathcal{L}_3(H) = (y_1 = x_1, x_k, x_{k-1}, \dots, x_s = y_2)$ but C does not have any chord, then H has no transversal triangles. Hence, $\{y_1, y_2\} \in E(C)$. Now, if $P = (y = z_1, z_2, \dots, z_{s_2} = x)$ is a path between x and C with $V(P) \cap V(C) = \{y\}$, we will prove that $y \in \{y_1, y_2\}$. By way of contradiction suppose that $y \notin \{y_1, y_2\}$. We take $k_3 = \min\{i \mid z_i \in V(P_1) \cup V(P_2)\}$, but $y \notin V(P_1) \cup V(P_2)$ so $k_3 > 1$. We can suppose that $z_{k_3} \in V(P_1)$. Set $P_1 = (w_1 = y_1, w_2, \dots, w_{s_1} = x)$ and $z_{k_3} = w_l$. Thus, there exist paths $L_1 = (y_1 = w_1, w_2, \dots, w_l = z_{k_3})$, $L_2 = (y = z_1, z_2, \dots, z_{k_3})$, and $L_3 = (y_2, P_2, x = w_{s_1}, w_{s_1-1}, \dots, w_l = z_{k_3})$ such that $V(L_1) \cap V(C) = \{y_1\}$, $V(L_2) \cap V(C) = \{y\}$ and $V(L_3) \cap V(C) = \{y_2\}$. Then, $\{y, y_1\}, \{y, y_2\} \in E(C)$, but $\{y_1, y_2\} \in E(C)$ and C is a hole, which is a contradiction. Therefore, $y \in \{y_1, y_2\}$. $\square$

Definition 11 A graph is chordal if each of its cycles of length at least 4 has a chord.

Proposition 8 ([14]) Chordal graphs are exactly the graphs that can be formed by clique-sums of complete graphs.

Theorem 4 ($\forall \theta \exists \Delta$) Let G be a graph. The following conditions are equivalent:

- (i) G is a theta-ring graph (i.e., $\Theta_\Delta^c(G) = \emptyset$).
- (ii) G can be constructed by 0, 1, 2-clique-sums of chordal graphs and/or cycles.
- (iii) G can be constructed by clique-sums of complete graphs and/or cycles.

- (iv) G does not contain as induced subgraph any graph from the following families:
- (a) θ -partial wheels.
 - (b) Prisms, pyramids and thetas.

Proof (i) $\Leftrightarrow$ (iv) By Corollary 5.

(iii) $\Rightarrow$ (i) Complete graphs and cycles are theta-ring graphs (since they do not contain chorded-thetas). The result follows from Proposition 7.

(i) $\Rightarrow$ (iii) Suppose that G can be constructed by 0, 1, 2-clique-sums of $H_1, \dots, H_s$, where each H_i is a chordal graph or cycle. If $s = 1$, then G is a cycle or chordal graph. Hence, G satisfies (iii) by Proposition 8. Now, we can suppose that $s > 1$ and we take G' to be the induced subgraph obtained by 0, 1, 2-clique-sum of $H_1, \dots, H_{s-1}$ such that G is the j -clique-sum of G' and H_s where $j \in \{0, 1, 2\}$. Furthermore $|V(G') \cap V(H_s)| \leq j$. By induction hypothesis we see that G' can be constructed by clique-sums of $L_1, \dots, L_{l_1}$, where each L_i is a cycle or a complete subgraph of G' . If H_s is a cycle then we obtain the result. Hence, we can assume that H_s is a chordal graph. Thus, by Proposition 8, we find that H_s can be constructed by clique-sums of the complete subgraphs $K_1, K_2, \dots, K_{l_2}$. Since $|V(G') \cap V(H_s)| \leq j$ we can suppose that $G' \cap H_s \subset K_u$ for some $u \in \{1, \dots, l_2\}$. Therefore G can be constructed by clique-sums of $L_1, \dots, L_{l_1}, K_u, K_{u+1}, \dots, K_{l_2}, K_{u-1}, \dots, K_1$.

(i) $\Rightarrow$ (ii) By induction on $|V(G)|$. We can suppose that $|V(G)| > 1$. Let $B_1, \dots, B_s$ be the blocks of G . If $s > 1$ then by the induction hypothesis $B_1, \dots, B_s$ satisfy (ii). But G can be obtained as 0 or 1-clique-sums of its blocks so G satisfies (ii). Thus, we can assume that $s = 1$ and $G = B_1$. We can suppose that G is not a chordal graph. Thus, G has a cycle C without chords such that C is not a triangle, i.e., C is a hole. Since G is a block we find that G is 2-connected. Set $C = (x_1, \dots, x_k, x_{k+1} = x_1)$ and we define

$$C_i = \left\{ x \in V(G \setminus C) \left| \begin{array}{l} \text{there exist paths } L_x, L'_x \text{ between } x \text{ and } C \text{ such that} \\ V(L_x) \cap V(L'_x) = \{x\}, V(L_x) \cap V(C) = \{x_i\} \\ \text{and } V(L'_x) \cap V(C) = \{x_{i+1}\} \end{array} \right. \right\}$$

for $i = 1, \dots, k$. We will prove that $C_1, \dots, C_k$ is a partition of $V(G \setminus C)$. By Lemma 10 if $x \in V(G \setminus C)$ then there exist paths L_1 and L_2 between x and C , such that $V(L_1) \cap V(L_2) = \{x\}$ and $V(L_i) \cap V(C) = \{y_i\}$ for $i = 1, 2$ with $y_1 \neq y_2$. Hence, by Lemma 11, $\{y_1, y_2\} \in E(C)$. Thus, there exists $i \in \{1, \dots, k\}$ such that $x_i = y_1$ and $y_2 = x_{i-1}$ or $y_2 = x_{i+1}$. So, $x \in V(C_i)$ or $x \in V(C_{i-1})$. Therefore $V(G \setminus C) = \bigcup_{i=1}^k V(C_i)$. Now, we take $z \in V(C_i) \cap V(C_j)$ with $i \neq j$. By Lemma 11, we have $\{x_i, x_{i+1}\} \subset \{x_j, x_{j+1}\}$ and $\{x_j, x_{j+1}\} \subset \{x_i, x_{i+1}\}$. Then, $i = j$ which is a contradiction. Hence, $V(C_i) \cap V(C_j) = \emptyset$ and $V(C_1), \dots, V(C_r)$ is a partition of $V(G \setminus C)$. Let G_i be the subgraph induced by $V(C_i) \cup \{x_i, x_{i+1}\}$. By the definition of C_i we have $V(L_x) \cup V(L'_x) \subset V(G_i)$ for $x \in V(C_i)$, so G_i is 2-connected, and by the induction hypothesis G_i can be obtained as 0,1,2-clique-sums of chordal graphs and cycles. We define B_1 as the 2-clique-sum of C and G_1 at the edge $\{x_1, x_2\}$. Now, we will construct successively B_{i+1} as the 2-clique-sum of B_i and G_{i+1} at the edge $\{x_{i+1}, x_{i+2}\}$ for $i = 1, \dots, k - 1$. But $B_{k-1} = G$ therefore G can be constructed as 0, 1, 2-clique-sums of chordal graphs and/or cycles. $\square$

5 CIO graphs and theta-ring graphs

In this section we prove that the CIO property is closed under 0, 1, 2-clique-sums and that the chordal graphs are CIO graphs. In particular, we prove that if $D = G_{\mathcal{O}}$ is a tournament, then P_D is a binomial complete intersection. Also we prove that prisms, pyramids, thetas and θ -partial wheels (the forbidden induced subgraphs given in Corollary 5) are not CIO graphs. The main result in this section is the equivalence between CIO graphs and theta-ring graphs.

Definition 12 An oriented complete graph is called a tournament.

Proposition 9 ([11]) *If $D_n = (K_n)_{\mathcal{O}_n}$ is a tournament, then D_n has a Hamiltonian oriented path.*

Lemma 12 *Let $\mathcal{O}$ be an orientation of G and x a vertex of G with $N_G(x) = \{x_1, \dots, x_{r_1}\}$ such that $P = (x_1, \dots, x_{r_1})$ is an oriented path in $D = G_{\mathcal{O}}$. Let $\mathcal{O}'$ be the orientation of $G' = G \setminus x$ induced by $\mathcal{O}$. If $C_i = (x, x_i, x_{i+1}, x)$, then $P_D \subseteq (P_{D'} \cup \{t_{C_1}, \dots, t_{C_{r_1-1}}\})$, where $D' = G'_{\mathcal{O}'}$.*

Proof Since P is an oriented path, we can assume that $(x_j, x_{j+1}) \in E(D)$ for $j = 1, \dots, r_1 - 1$. For $1 \leq i < j \leq r_1$, we take the cycle $C_{i,j} = (x, x_i, P_{i,j}, x_j, x)$ where $P_{i,j}$ is the oriented subpath of P that joins x_i and x_j . First we will prove that $t_{C_{i,j}} \in (t_{C_1}, \dots, t_{C_{r_1-1}})$ by induction on $j - i$. If $j - i = 1$ it is clear. Using the induction hypothesis we have $t_{C_{i,j-1}} \in (t_{C_1}, \dots, t_{C_{r_1-1}})$. Thus, by Lemma 2 we have $t_{C_{i,j}} \in (t_{C_{i,j-1}}, t_{C_{j-1}})$. Therefore $t_{C_{i,j}} \in (t_{C_1}, \dots, t_{C_{r_1-1}})$.

Let C be a cycle without chords of D . If $x \notin V(C)$ then $C \subset D'$ and $t_C \in P_{D'}$. Now, if $x \in V(C)$ then $C = (x, x_{j_1}, L, x_{j_2}, x)$ where $1 \leq j_1 < j_2 \leq r_1$ and L is a path between x_{j_1} and x_{j_2} in D . Since C is a cycle without chords, $V(L) \cap N_G(x) = \{x_{j_1}, x_{j_2}\}$. On the other hand we take the cycle $C' = (x_{j_1}, L, x_{j_2}, P_{j_1, j_2}, x_{j_1})$, then $C' \subseteq D'$ and $t_{C'} \in P_{D'}$. By Lemma 2 we have $t_C \in (t_{C'}, t_{C_{i,j}})$. Furthermore $t_{C_{i,j}} \in (t_{C_1}, \dots, t_{C_{r_1-1}})$, then $t_C \in (P_{D'}, t_{C_1}, \dots, t_{C_{r_1-1}})$. Therefore $P_D \subseteq (P_{D'} \cup \{t_{C_1}, \dots, t_{C_{r_1-1}}\})$, by Proposition 4. $\square$

Definition 13 Let G be a graph. A simplicial vertex x is a vertex such that the induced subgraph obtained from $N_G(x)$ in G is a complete graph.

Proposition 10 ([4]) *If G is a chordal graph, then G has a simplicial vertex.*

Theorem 5 *If G is a chordal graph then G is a CIO graph.*

Proof By induction on $|V(G)|$. By Proposition 5 we can suppose that G is a connected graph. Since G is chordal there exists a vertex x in G such that x is a simplicial vertex. Thus, if K is the subgraph induced by $N_G(x)$, then K is a complete graph and $G' = G \setminus x$ is a connected subgraph of G . We can suppose that $N_G(x) = \{x_1, \dots, x_s\}$. Let $\mathcal{O}$ be an orientation of G and $D = G_{\mathcal{O}}$. By induction hypothesis G' is CIO, hence if $\mathcal{O}'$ is the orientation of G' induced by $\mathcal{O}$ and $D' = G'_{\mathcal{O}'}$, then there exists the binomial generating set $\mathcal{G}'$ of $P_{D'}$ with $|\mathcal{G}'| = (q - s) - (n - 1) + 1$. On the other hand,

by Proposition 9, there exists a Hamiltonian oriented path P of K . We can suppose that $P = (x_1, \dots, x_s)$. We take the triangle $C_i = (x, x_i, x_{i+1}, x)$ for $i = 1, \dots, s-1$ and $\mathcal{G} = \mathcal{G}' \cup \{t_{C_1}, \dots, t_{C_{s-1}}\}$. Hence, by Lemma 12, $\mathcal{G}$ is a generating set of P_D and $|\mathcal{G}| = q - n + 1$. Therefore G is CIO. $\square$

Corollary 6 *If D is a tournament, then P_D is a binomial complete intersection.*

Proof Since D is a tournament then $D = K_{\mathcal{O}}$ where K is a complete graph. Thus K is a chordal graph and by Theorem 5 K is a CIO graph. Therefore P_D is a binomial complete intersection. $\square$

Proposition 11 *Let G_1, G_2 be two connected graphs and let G be the 2-clique-sum of G_1 and G_2 . If G_1 and G_2 are CIO graphs then G is a CIO graph.*

Proof Let $\mathcal{O}$ be an orientation of the edges of G and let $\mathcal{O}_i$ be the orientation of the edges of G_i induced by $\mathcal{O}$, for $i = 1, 2$. If $|V(G_i)| = n_i$ and $|E(G_i)| = q_i$, for $i = 1, 2$, then there exists a set B_i with $q_i - n_i + 1$ binomials such that B_i generates P_{D_i} , where $D_i = (G_i)_{\mathcal{O}_i}$. Since G is the 2-clique-sum of G_1 and G_2 then $|E(G)| = q_1 + q_2 - 1$ and $|V(G)| = n_1 + n_2 - 2$. If C is a cycle without chords of D then $V(C) \subset V(D_1)$ or $V(C) \subset V(D_2)$. Hence, $t_C \in (B_1 \cup B_2)$ and by Proposition 4, we see that $B_1 \cup B_2$ is a generating set of $P_{G_{\mathcal{O}}}$. But

$$|B_1 \cup B_2| = \sum_{i=1}^2 (q_i - n_i + 1) = (q_1 + q_2 - 1) - (n_1 + n_2 - 2) + 1 = \text{ht}(P_{G_{\mathcal{O}}}).$$

Therefore G is CIO. $\square$

If $L = (t_{i_1}, t_{i_2}, \dots, t_{i_s})$ is an oriented path in $D = G_{\mathcal{O}}$, we denote by t_L the product of the edges of L , i.e., $t_L = t_{i_1} t_{i_2} \cdots t_{i_s}$.

Proposition 12 *Prisms, pyramids, thetas and θ -partial wheels are not CIO graphs.*

Proof Case (a). T is a partial wheel where $V(T) = \{x_1, \dots, x_m, x\}$ with center x , rim $C = (x_1, x_2, \dots, x_m, x_1)$, and $N_T(x) = \{x_{j_1}, \dots, x_{j_k}\}$.

First subcase, if $k \geq 4$ we take $\mathcal{T} = T_{\mathcal{O}}$ with the acyclic orientation $\mathcal{O}$ given into Fig. 1.

We have $\text{ht}(P_{\mathcal{T}}) = k$. Let $\mathcal{G} = \{f_1, \dots, f_s\}$ be a minimal binomial set of generators for $P_{\mathcal{T}}$. Since there are no oriented cycles in $\mathcal{T}$ and using Corollary 1, then there is no binomial having as a monomial 1. Since $g_i \in P_{\mathcal{T}}$, then there exists a binomial $f_{l_i} \in \mathcal{G}$ such that f_{l_i} has as a monomial t_i for $i = 2, \dots, k-2$ and $i = k$. Without loss of generality we can suppose that $l_i = i$ for $i = 2, \dots, k-2$ and $l_k = k-1$. Observe that if C is a cycle of T and $E(C) \cap E(L_i) \neq \emptyset$ then $E(L_i) \subset E(C)$. We will use this observation in the rest of the proof.

By the form of g_{k+1} there exists a binomial $f_j \in \mathcal{G}$ with a monomial that divides $t_{L_{k-2}} t_{L_k}$. We can suppose that $f_j = t^{\alpha_j} - t^{\beta_j}$ and $t^{\alpha_j} | t_{L_{k-2}} t_{L_k}$. By Theorem 1 there exists a cycle C_1 such that the monomials of $t_{C_1} = t^{\alpha'_j} - t^{\beta'_j}$ divide the monomials

$$\begin{aligned}
 g_1 &= t_2 - t_1 t_{L_1} \\
 g_2 &= t_3 - t_2 t_{L_2} \\
 &\vdots \\
 g_i &= t_{i+1} - t_i t_{L_i} \\
 &\vdots \\
 g_{k-3} &= t_{k-2} - t_{k-3} t_{L_{k-3}} \\
 g_{k-2} &= t_{k-2} - t_{k-1} t_{L_{k-2}} \\
 g_{k-1} &= t_k - t_{k-1} t_{L_{k-1}} \\
 g_k &= t_k - t_1 t_{L_k} \\
 g_{k+1} &= t_{L_{k-2}} t_{L_k} - t_{L_1} \cdots t_{L_{k-3}} t_{L_{k-1}}
 \end{aligned}$$

Fig. 1 Acyclic orientation of the partial wheel $\mathcal{T}$ with $k \geq 4$, and the equations of its chordless cycles

of f_j . First we prove $E(L_k) \subset E(C_1)$. Suppose that it is not true, then $E(L_k) \cap E(C_1) = \emptyset$. Hence, $E(L_{k-2}) \cap E(C_1) \neq \emptyset$ and $E(L_{k-2}) \subset E(C_1)$. Furthermore since $t^{\alpha'_j} | t_{L_{k-2}} t_{L_k}$, then $t_{k-1} \notin E(C_1)$ and $E(L_{k-1}) \subset E(C_1)$. But $E(L_k) \cap E(C_1) = \emptyset$, then $t_k \in E(C_1)$ and $t_k | t^{\alpha'_j}$. This is a contradiction because t_k does not divide $t_{L_{k-2}} t_{L_k}$. Therefore $E(L_k) \subset E(C_1)$. Now, we prove that $x \notin E(C_1)$. Suppose $x \in E(C_1)$, we take the minimal $i_1 = \min\{i \in \{1, \dots, k\} \mid t_i \in E(C_1)\}$. Thus, $\bigcup_{i=1}^{i_1} E(L_i) \subset E(C_1)$ and $t_{i_1} | t^{\alpha'_j}$. But this is not possible because t_{i_1} does not divide $t_{L_{k-2}} t_{L_k}$. Hence, $x \notin E(C_1)$ and $C_1 = \mathcal{T} \setminus x$. Therefore $t_{C_1} = g_{k+1}$. Using Proposition 1 we have $f_j = g_{k+1}$. Furthermore $f_j \neq f_i$ for $i = 2, \dots, k-1$, then we can suppose that $j = 1$.

On the other hand by the form of g_{k-1} there exists a binomial $f_{j_1} \in \mathcal{G}$ with a monomial that divides $t_{k-1} t_{L_{k-1}}$. By Theorem 1 there exists a cycle C_2 such that the monomials of $t_{C_2} = t^{\alpha'_{j_1}} - t^{\beta'_{j_1}}$ divide the monomials of f_j . We can suppose that $t^{\alpha'_{j_1}} | t_{k-1} t_{L_{k-1}}$. Suppose $E(L_{k-1}) \cap E(C_2) = \emptyset$ then $t_{k-1} \in E(C_2)$ and $E(L_{k-2}) \subset E(C_2)$. Thus, $t_{L_{k-2}} | t^{\alpha'_{j_1}}$ which is not possible because $t_{L_{k-2}}$ does not divide $t_{k-1} t_{L_{k-1}}$. So, $E(L_{k-1}) \cap E(C_2) \neq \emptyset$ and $E(L_{k-1}) \subset E(C_2)$. Now, we prove that $t_{k-1} \in E(C_2)$. If $t_{k-1} \notin E(C_2)$ then $E(L_{k-2}) \subset E(C_2)$. Furthermore, $E(L_{k-3}) \subset E(C_2)$ or $t_{k-2} \in E(C_2)$. Hence, $t_{L_{k-3}} | t^{\alpha'_{j_1}}$ or $t_{k-2} | t^{\alpha'_{j_1}}$. But this is not possible, then $t_{k-1} \in E(C_2)$. Now, $t_k \in E(C_2)$ or $E(L_k) \subset E(C_2)$. If $t_k \in E(C_2)$ then $C_2 = (t_{k-1}, L_{k-1}, t_k)$ and $t_{C_2} = -g_{k-1}$. For the other case, if $E(L_k) \subset E(C_2)$ then $C_2 = (t_{k-1}, L_{k-1}, L_k, t_1)$ and $t_{C_2} = t_{k-1} t_{L_{k-1}} - t_1 t_{L_k} = g_k - g_{k-1}$. In both cases, by Proposition 1 we have $f_{j_1} = t_{C_2}$. Furthermore, by the form of g_k there exists a binomial $f_{j_2} \in \mathcal{G}$ with a monomial that divides $t_1 t_{L_k}$. In a similar way as in the last argument, we find that there exists a cycle C_3 such that $f_{j_2} = t_{C_3}$ where $C_3 = (t_k, L_k, t_1)$ or $C_3 = (t_{k-1}, L_{k-1}, L_k, t_1)$. So, $f_{j_1}, f_{j_2} \in \{t_k - t_{k-1} t_{L_{k-1}}, t_k - t_1 t_{L_k}, t_{k-1} t_{L_{k-1}} - t_1 t_{L_k}\}$. If $\{f_1, \dots, f_{k-1}\} \cap \{f_{j_1}, f_{j_2}\} \neq \emptyset$ then $f_{j_1} = f_{k-1}$ or $f_{j_2} = f_{k-1}$ and in this case $f_{j_1} \neq f_{j_2}$. Therefore $k \leq |\{f_1, \dots, f_{k-1}, f_{j_1}, f_{j_2}\}|$.

By the form of g_{k-2} there exists a binomial f_{j_3} with a monomial that divides $t_{k-1} t_{L_{k-2}}$. Thus, there exists a cycle C_4 such that the monomials of $t_{C_4} = t^{\alpha'_{j_3}} - t^{\beta'_{j_3}}$ divide the monomials of f_{j_3} . We can suppose that $t^{\alpha'_{j_3}} | t_{k-1} t_{L_{k-2}}$. Since $t^{\alpha'_{j_3}} \neq 1$

$$\begin{aligned} g_1 &= t_3 - t_1 t_{L_1} \\ g_2 &= t_3 - t_2 t_{L_2} \\ g_3 &= t_1 t_{L_4} - t_2 t_{L_3} \\ g_4 &= t_{L_1} t_{L_3} - t_{L_2} t_{L_4} \end{aligned}$$

Fig. 2 Acyclic orientation of the partial wheel $\mathcal{T}$ with $k = 3$, and the equations of its chordless cycles

and $t_{L_{k-1}}$ does not divide $t^{\alpha'_{j_3}}$ then $V(L_{k-2}) \subset V(C_4)$. Furthermore, since t_{L_k} and t_k do not divide $t^{\alpha'_{j_3}}$ then $t_{k-1} \in E(C_4)$. Hence, if $t_{k-2} \in E(C_4)$ then $C_4 = (t_{k-2}, L_{k-2}, t_{k-1})$ and $t_{C_4} = t_{k-1} t_{L_{k-2}} - t_{k-2}$. In the other case, if $t_{k-2} \notin E(C_4)$ then $t_{C_4} = t_{k-1} t_{L_{k-2}} - t^{\beta'_{j_3}}$. Thus, $t^{\alpha'_{j_3}} = t_{k-1} t_{L_{k-2}}$ and by Proposition 1 we have $f_{j_3} = t_{C_4}$. Furthermore, if $t_{k-2} \notin E(C_4)$ then $k + 1 \leq |\{f_1, \dots, f_{k-1}, f_{j_1}, f_{j_2}, f_{j_3}\}|$ and $P_{\mathcal{T}}$ is not a binomial complete intersection. So, we can assume that $t_{k-2} \in E(C_4)$ and $f_{j_3} = t_{k-1} t_{L_{k-2}} - t_{k-2}$. Moreover, if $f_{j_3} \neq -f_{k-2}$ then $k + 1 \leq |\{f_1, \dots, f_{k-1}, f_{j_1}, f_{j_2}, f_{j_3}\}|$ and $P_{\mathcal{T}}$ is not a binomial complete intersection. Hence, we can assume $-f_{k-2} = f_{j_3} = t_{k-1} t_{L_{k-2}} - t_{k-2}$ and it is the unique binomial in $\mathcal{G}$ which has a monomial that divides either t_{k-2} or $t_{k-1} t_{L_{k-2}}$. Then,

$$\begin{aligned} t_{k-2} - t_{k-3} t_{L_{k-3}} &= \sum_{i=1}^s g_i f_i = g_{k-2} (t_{k-1} t_{L_{k-2}} - t_{k-2}) + \sum_{i \neq k-2} g_i f_i \\ &= (-1 + g'_{k-2}) (t_{k-1} t_{L_{k-2}} - t_{k-2}) + \sum_{i \neq k-2} g_i f_i. \end{aligned}$$

Here $\deg(g'_{k-2}) > 0$ or $g'_{k-2} = 0$. Using the last equation we obtain

$$t_{k-1} t_{L_{k-2}} - t_{k-3} t_{L_{k-3}} = g'_{k-2} (t_{k-1} t_{L_{k-2}} - t_{k-2}) + \sum_{i \neq k-2} g_i f_i.$$

So, there exists a f_i with $i \neq k - 2$ where one of its monomials divides $t_{k-1} t_{L_{k-2}}$. But this is not possible.

Second subcase, if $k = 3$ we take $\mathcal{T} = T_{\mathcal{O}}$ with the acyclic orientation $\mathcal{O}$ given in Fig. 2.

We have $\text{ht}(P_{\mathcal{T}}) = 3$. Let $\mathcal{G} = \{f_1, \dots, f_s\}$ be a minimum binomial set of generators of $P_{\mathcal{T}}$. For every $m \in A = \{t_3, t_1 t_{L_1}, t_2 t_{L_2}, t_1 t_{L_4}, t_2 t_{L_3}, t_{L_1} t_{L_3}, t_{L_2} t_{L_4}\}$ there exists a binomial $f \in \mathcal{G}$ such that f has a monomial t^{α} and $t^{\alpha} | m$. By Theorem 1 there exists a cycle C such that $t_C = t^{\alpha'} - t^{\beta'}$ and $t^{\alpha'} | t^{\alpha}$. But there are no cycles whose binomials have either t_{L_i} or t_j as a monomial, for $i = 1, \dots, 4$ and $j = 1, 2$. Hence, $t^{\alpha'} = t^{\alpha} = m$ and by Proposition 1 $t_C = f$. Then, for every $m \in A$ there exists $f \in \mathcal{G}$ such that m is a monomial of f . But $|A|/2 \leq |\mathcal{G}|$ and $|A| = 7$ then $4 \leq |\mathcal{G}|$. Therefore $P_{\mathcal{T}}$ is not a binomial complete intersection.

$$\begin{aligned} g_1 &= t_4 - t_5 t_6 \\ g_2 &= t_3 - t_1 t_2 \\ g_3 &= t_3 - t_{L_3} t_5 t_{L_2} \\ g_4 &= t_1 t_6 - t_{L_1} t_{L_3} \\ g_5 &= t_4 t_{L_2} - t_2 t_{L_1} \end{aligned}$$

Fig. 3 Acyclic orientation of the prism T , and the equations of its cycles without chords

$$\begin{aligned} g_1 &= t_3 - t_{L_1} t_{L_2} \\ g_2 &= t_3 - t_1 t_2 \\ g_3 &= t_{L_2} t_{L_4} - t_2 t_{L_3} \\ g_4 &= t_{L_1} t_{L_3} - t_1 t_{L_4} \end{aligned}$$

Fig. 4 Acyclic orientation of the pyramid T , and the equations of its cycles without chords

Case (b-1). If T is a prism, we take $T = T_{\mathcal{O}}$ with the acyclic orientation $\mathcal{O}$ given in Fig. 3.

We have $\text{ht}(P_T) = 4$. Let $\mathcal{G} = \{f_1, \dots, f_s\}$ be a minimal binomial set of generators of P_T . We take $A = \{t_5 t_6, t_1 t_2, t_2 t_{L_1}, t_{L_3} t_5 t_{L_2}, t_1 t_6, t_{L_1} t_{L_3}, t_3, t_4\}$. By the same argument as for partial wheels with $k = 3$, we see that for every $m \in A$ there exist $f \in \mathcal{G}$ and a cycle C of T such that m is a monomial of f and $f = t_C$. Let $f_{j_1}, f_{j_2}, f_{j_3}$ be binomials of $\mathcal{G}$ that have as a monomial $t_3, t_1 t_2$ and $t_{L_3} t_5 t_{L_2}$, respectively. Hence, if $f_{j_1} = t_{C_1}$, $f_{j_2} = t_{C_2}$ and $f_{j_3} = t_{C_3}$, then

$$C_1, C_2, C_3 \in \{(t_3, t_1, t_2), (t_1, t_2, L_2, t_5, L_3), (t_3, L_3, t_5, L_2)\}.$$

Thus, if $B = \{f_{j_1}, f_{j_2}, f_{j_3}\}$ then $|B| \geq 2$. Since if $f_{j_1} = f_{j_3}$ then $C_2 = C_3 = (t_1, t_2, L_2, t_5, L_3)$ and $f_{j_1} \neq f_{j_2}$. Furthermore for every monomial in $A' = A \setminus \{t_3, t_1 t_2, t_{L_3} t_5 t_{L_2}\}$ there exists a binomial in $\mathcal{G} \setminus B$. But $|\mathcal{G} \setminus B| \geq |A'|/2$ and $|A'| = 5$ then $|\mathcal{G} \setminus B| \geq 3$ and $|\mathcal{G}| \geq 5$. Therefore P_T is not a binomial complete intersection.

Case (b-2). If T is a pyramid, we take $T = T_{\mathcal{O}}$ with the acyclic orientation $\mathcal{O}$ given in Fig. 4.

This case is similar to the case when T is a partial wheel with $k = 3$, only change t_1 by the path L_2 . Therefore P_T is not a binomial complete intersection.

Case (b-3) If T is a theta, we take $T = T_{\mathcal{O}}$ with the acyclic orientation $\mathcal{O}$ given in Fig. 5.

We have $\text{ht}(P_T) = 2$. Let $\mathcal{G} = \{f_1, \dots, f_s\}$ be a minimum binomial set of generators of P_T . We take $A = \{t_{L_1} t_{L_4}, t_{L_2} t_{L_3}, t_{L_3} t_{L_6}, t_{L_4} t_{L_5}, t_{L_1} t_{L_6}, t_{L_2} t_{L_5}\}$. By the same

Fig. 5 Acyclic orientation of the theta $\mathcal{T}$, and the equations of its cycles without chords

argument as for partial wheels with $k = 3$, we find that for every $m \in A$ there exist $f \in \mathcal{G}$ and a cycle C of $\mathcal{T}$ such that m is a monomial of f and $f = t_C$. Hence, $|\mathcal{G}| \geq |A|/2$ but $|A| = 6$ then $|\mathcal{G}| \geq 3$. Therefore $P_{\mathcal{T}}$ is not a binomial complete intersection. $\square$

Theorem 6 *Let G be a graph. G is a CIO graph if and only if G is a theta-ring graph.*

Proof ($\Rightarrow$) If G is not a theta-ring graph, then by Theorem 4, G has an induced subgraph H that is either a prism, pyramid, theta or θ -partial wheel. Thus, by Proposition 12, H is not a CIO graph. Hence, by Corollary 4, G is not a CIO graph which is a contradiction. Therefore G is a theta-ring graph.

($\Leftarrow$) By Theorem 4 G can be constructed by 0, 1, 2-clique-sums of chordal graphs and/or cycles. By Proposition 5 chordal graphs and cycles are CIO graphs. Therefore by Proposition 5, Proposition 6 and Proposition 11, G is a CIO graph. $\square$

Definition 14 ([5]) A graph G is a *ring graph* if each block of G which is not a bridge or a vertex can be constructed successively by 2-clique-sums of cycles.

Corollary 7 ([6]) *If G is a ring graph, then G is a CIO graph. The converse holds if G is bipartite.*

Proof First we observe that ring graphs can be obtained by 0, 1, 2-clique-sums of vertices, edges and cycles. Hence, by Theorem 4, G is a theta-ring graph and by Theorem 6, G is a CIO graph.

Now, if G is bipartite and G is a CIO graph, then by Theorem 4 and Theorem 6, G can be constructed by 0, 1, 2-clique-sums of chordal graphs and/or cycles. Since G is bipartite, if H is a chordal induced subgraph of G , then H is a forest. Hence, G can be constructed by 0, 1, 2-clique-sums of cycles and/or edges. Therefore, G is a ring graph. $\square$

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