# The Ehrhart Polynomial of the Birkhoff Polytope 

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All means (even continuous) sanctify the discrete end. Doron Zeilberger ${ }^{1}$


#### Abstract

The $n$th Birkhoff polytope is the set of all doubly stochastic $n \times n$ matrices, that is, those matrices with nonnegative real coefficients in which every row and column sums to one. A wide open problem concerns the volumes of these polytopes, which have been known for $n \leq 8$. We present a new, complex-analytic way to compute the Ehrhart polynomial of the Birkhoff polytope, that is, the function counting the integer points in the dilated polytope. One reason to be interested in this counting function is that the leading term of the Ehrhart polynomial is-up to a trivial factor-the volume of the polytope. We implemented our methods in the form of a computer program, which yielded the Ehrhart polynomial (and hence the volume) of the ninth Birkhoff polytope, as well as the volume of the tenth Birkhoff polytope.


## 1. Introduction

One of the most intriguing objects of combinatorial geometry is the $n$th Birkhoff polytope

$$
\mathcal{B}_{n}=\left\{\left(\begin{array}{ccc}
x_{11} & \cdots & x_{1 n} \\
\vdots & & \vdots \\
x_{n 1} & \ldots & x_{n n}
\end{array}\right) \in \mathbb{R}^{n^{2}}: x_{j k} \geq 0, \begin{array}{l}
\sum_{j} x_{j k}=1 \text { for all } 1 \leq k \leq n \\
\sum_{k} x_{j k}=1 \text { for all } 1 \leq j \leq n
\end{array}\right\}
$$

often described as the set of all $n \times n$ doubly stochastic matrices. $\mathcal{B}_{n}$ is a convex polytope with integer vertices. It possesses fascinating combinatorial properties [4]-[6], [8], [23] and relates to many mathematical areas [10], [14]. A long-standing open problem is the determination of the relative volume of $\mathcal{B}_{n}$, which had been known only up

[^0]to $n=8$ [7], [17]. In this paper we propose a new method of calculating this volume and use it to compute $\operatorname{vol} \mathcal{B}_{9}$ and $\operatorname{vol} \mathcal{B}_{10}$.

One of the recent attempts to compute vol $\mathcal{B}_{n}$ relies on the theory of counting functions for the integer points in polytopes. Ehrhart proved [11] that for a polytope $\mathcal{P} \subset \mathbb{R}^{d}$ with integral vertices, the number

$$
L_{\mathcal{P}}(t):=\#\left(t \mathcal{P} \cap \mathbb{Z}^{d}\right)
$$

is a polynomial in the positive integer variable $t$. He showed various other properties for this counting function (in fact, in the more general setting of $\mathcal{P}$ having rational vertices), of which we mention three here:

- The degree of $L_{\mathcal{P}}$ is the dimension of $\mathcal{P}$.
- The leading term of $L_{\mathcal{P}}$ is the relative volume of $\mathcal{P}$, normalized with respect to the sublattice of $\mathbb{Z}^{d}$ on the affine subspace spanned by $\mathcal{P}$.
- Since $L_{\mathcal{P}}$ is a polynomial, we can evaluate it at nonpositive integers. These evaluations yield

$$
\begin{align*}
L_{\mathcal{P}}(0) & =\chi(\mathcal{P})  \tag{1}\\
L_{\mathcal{P}}(-t) & =(-1)^{\operatorname{dim} \mathcal{P}} L_{\mathcal{P}^{\circ}}(t) \tag{2}
\end{align*}
$$

Here $\chi(\mathcal{P})$ denotes the Euler characteristic and $\mathcal{P}^{\circ}$ the relative interior of $\mathcal{P}$. The reciprocity law (2) was in its full generality proved by Macdonald [15].

Application of this theory to the Birkhoff polytope $\mathcal{B}_{n}$ incorporates the nice interpretation of the number of integral points in $t \mathcal{B}_{n}$ as the number of semimagic squares, namely, square matrices whose nonnegative integral coefficients sum up to the same integer $t$ along each row and column.

We denote the Ehrhart polynomial of $\mathcal{B}_{n}$ by

$$
H_{n}(t):=L_{\mathcal{B}_{n}}(t)
$$

It is not hard to see that $\operatorname{dim} \mathcal{B}_{n}=(n-1)^{2}$, hence $H_{n}$ is a polynomial in $t$ of degree $(n-1)^{2}$. The first two of these polynomials are trivial:

$$
H_{1}(t)=1, \quad H_{2}(t)=t+1
$$

the first nontrivial case was computed by MacMahon [16] as

$$
\begin{equation*}
H_{3}(t)=3\binom{t+3}{4}+\binom{t+2}{2} \tag{3}
\end{equation*}
$$

The structural properties of $H_{n}$ were first studied in [12], [18], and [20]. It is a nice exercise to deduce from (2) that

$$
\begin{equation*}
H_{n}(-n-t)=(-1)^{n-1} H_{n}(t) \tag{4}
\end{equation*}
$$

and

$$
H_{n}(-1)=H_{n}(-2)=\cdots=H_{n}(-n+1)=0
$$

This allows the following strategy of computing $H_{n}$, and, therefore, the volume of $\mathcal{B}_{n}$ : compute the first $\binom{n-1}{2}$ values of $H_{n}$, use the above symmetry and trivial values of $H_{n}$, and calculate the polynomial $H_{n}$ by interpolation. In fact, as far as we are aware, the volume of $\mathcal{B}_{8}$ was computed using essentially this method, combined with some nice computational tricks [7], [17].

We propose a new, completely different approach for computing $H_{n}$ (and hence $\operatorname{vol} \mathcal{B}_{n}$ ). It is based on an analytic method by the first author of computing the Ehrhart polynomial of a polytope [2]. We introduce the application of this method to the Birkhoff polytope in the following section.

Some recent refreshing approaches-of a more algebraic-geometric/topological flavor-to the problem of computing vol $\mathcal{B}_{n}$ and $H_{n}$ can be found in [1], [9], and [22].

## 2. An Integral Counting Integers

A convex polytope $\mathcal{P} \subset \mathbb{R}^{d}$ is an intersection of halfspaces. This allows the compact description

$$
\mathcal{P}=\left\{\mathbf{x} \in \mathbb{R}^{d}: \mathbf{A x} \leq \mathbf{b}\right\}
$$

for some $m \times d$ matrix $\mathbf{A}$ and $m$-dimensional vector $\mathbf{b}$. Here the inequality is understood componentwise. In fact, we may convert these inequalities into equalities by introducing "slack variables." If $\mathcal{P}$ has rational vertices (those polytopes are called rational), we can choose $\mathbf{A}$ and $\mathbf{b}$ in such a way that all their entries are integers. In summary, we may assume that a convex rational polytope $\mathcal{P}$ is given by

$$
\begin{equation*}
\mathcal{P}=\left\{\mathbf{x} \in \mathbb{R}_{\geq 0}^{d}: \mathbf{A x}=\mathbf{b}\right\} \tag{5}
\end{equation*}
$$

where $\mathbf{A} \in M_{m \times d}(\mathbb{Z})$ and $\mathbf{b} \in \mathbb{Z}^{m}$. (If we are interested in counting the integer points in $\mathcal{P}$, we may assume that $\mathcal{P}$ is in the nonnegative orthant, i.e., the points in $\mathcal{P}$ have nonnegative coordinates, as translation by an integer vector does not change the latticepoint count.) The following straightforward theorem can be found in [2]. We use the standard multivariate notation $\mathbf{v}^{\mathbf{w}}:=v_{1}^{w_{1}} \cdots v_{n}^{w_{n}}$.

Theorem 1 [2, Theorem 8]. Suppose the convex rational polytope $\mathcal{P}$ is given by (5), and denote the columns of $\mathbf{A}$ by $\mathbf{c}_{1}, \ldots, \mathbf{c}_{d}$. Then

$$
L_{\mathcal{P}}(t)=\frac{1}{(2 \pi i)^{m}} \int_{\left|z_{1}\right|=\varepsilon_{1}} \cdots \int_{\left|z_{m}\right|=\varepsilon_{m}} \frac{z_{1}^{-t b_{1}-1} \cdots z_{m}^{-t b_{m}-1}}{\left(1-\mathbf{z}^{\mathbf{c}_{1}}\right) \cdots\left(1-\mathbf{z}^{\mathbf{c}_{d}}\right)} d \mathbf{z}
$$

Here $0<\varepsilon_{1}, \ldots, \varepsilon_{m}<1$ are distinct real numbers such that each factor in the denominator can be expanded into a geometric series.

It should be mentioned that $L_{\mathcal{P}}$ is in general not a polynomial if the vertices of $\mathcal{P}$ are not integral, but a quasipolynomial, that is, an expression of the form

$$
c_{d}(t) t^{d}+\cdots+c_{1}(t) t+c_{0}(t)
$$

where $c_{0}, \ldots, c_{d}$ are periodic functions in $t$. (See, for example, Section 4.4 of [19] for more information about quasipolynomials.) Theorem 1 applies to these slightly more
general counting functions; however, in this article we only deal with polytopes with integer vertices, for which $L_{\mathcal{P}}$ is a polynomial.

We also note here that Theorem 1 can be used to compute quickly by hand formulas for certain classes of polytopes (see, for example, [2]). In this project we take a slightly different approach and use this theorem to derive efficiently formulas with the help of a computer.

We can view the Birkhoff polytope $\mathcal{B}_{n}$ as given in the form of (5), where

$$
\mathbf{A}=\left(\begin{array}{cccccccccc}
1 & \cdots & 1 & & & & & & & \\
& & & 1 & \cdots & 1 & & & & \\
& & & & & & \ddots & & & \\
1 & & & 1 & & & & 1 & \cdots & 1 \\
& \ddots & & & \ddots & & \cdots & & \ddots & \\
& & 1 & & & 1 & & & & 1
\end{array}\right)
$$

is a $2 n \times n^{2}$ matrix and $\mathbf{b}=(1, \ldots, 1) \in \mathbb{Z}^{2 n}$. Hence Theorem 1 gives for this special case

$$
H_{n}(t)=\frac{1}{(2 \pi i)^{2 n}} \int \cdots \int \frac{\left(z_{1} \cdots z_{2 n}\right)^{-t-1}}{\left(1-z_{1} z_{n+1}\right)\left(1-z_{1} z_{n+2}\right) \cdots\left(1-z_{n} z_{2 n}\right)} d \mathbf{z}
$$

Here it is understood that each integral is over a circle with radius $<1$ centered at 0 ; all appearing radii should be different. We can separate, say, the last $n$ variables and obtain

$$
\begin{aligned}
H_{n}(t)= & \frac{1}{(2 \pi i)^{n}} \\
& \times \int \cdots \int\left(z_{1} \cdots z_{n}\right)^{-t-1}\left(\frac{1}{2 \pi i} \int \frac{z^{-t-1}}{\left(1-z_{1} z\right) \cdots\left(1-z_{n} z\right)} d z\right)^{n} d z_{n} \cdots d z_{1} .
\end{aligned}
$$

We may choose the radius of the integration circle of the innermost integral to be smaller than the radii of the other integration paths. Then this innermost integral is easy to compute: it is equal to the residue at 0 of

$$
\frac{1}{z^{t+1}\left(1-z_{1} z\right) \cdots\left(1-z_{n} z\right)}
$$

and, by the residue theorem, equal to the negative of the sum of the residues at $z_{1}^{-1}, \ldots, z_{n}^{-1}$. (Note that here we use the fact that $t>0$.) The residues at these simple poles are easily computed: the one, say, at $1 / z_{1}$ can be calculated as

$$
\begin{aligned}
& \lim _{z \rightarrow 1 / z_{1}}\left(z-\frac{1}{z_{1}}\right) \frac{1}{z^{t+1}\left(1-z_{1} z\right) \cdots\left(1-z_{n} z\right)} \\
& \quad=\lim _{z \rightarrow 1 / z_{1}} \frac{z-1 / z_{1}}{1-z_{1} z} \frac{1}{z^{t+1}\left(1-z_{2} z\right) \cdots\left(1-z_{n} z\right)} \\
& \quad=-\frac{1}{z_{1}} \frac{z_{1}^{t+1}}{\left(1-z_{2} / z_{1}\right) \cdots\left(1-z_{n} / z_{1}\right)}=-\frac{z_{1}^{t+n-1}}{\left(z_{1}-z_{2}\right) \cdots\left(z_{1}-z_{n}\right)}
\end{aligned}
$$

This yields the starting point for our computations.

Theorem 2. For any distinct $0<\varepsilon_{1}, \ldots, \varepsilon_{n}<1$,

$$
H_{n}(t)=\frac{1}{(2 \pi i)^{n}} \int_{\left|z_{1}\right|=\varepsilon_{1}} \cdots \int_{\left|z_{n}\right|=\varepsilon_{n}}\left(z_{1} \cdots z_{n}\right)^{-t-1}\left(\sum_{k=1}^{n} \frac{z_{k}^{t+n-1}}{\prod_{j \neq k}\left(z_{k}-z_{j}\right)}\right)^{n} d z_{n} \cdots d z_{1}
$$

Remark. It can be proved from the form of the integrand that $H_{n}$ is indeed a polynomial in $t$ : to compute the integral, one has to execute a (huge) number of limit computations, which yield "at worst" powers of $t$ (as a consequence of l'Hôpital's Rule). In fact, one can make this property more apparent by noticing that the expression in parenthesis is actually a polynomial; namely,

$$
H_{n}(t)=\frac{1}{(2 \pi i)^{n}} \int_{\left|z_{1}\right|=\varepsilon_{1}} \cdots \int_{\left|z_{n}\right|=\varepsilon_{n}}\left(z_{1} \cdots z_{n}\right)^{-t-1}\left(\sum_{m_{1}+\ldots+m_{n}=t} z_{1}^{m_{1}} \cdots z_{n}^{m_{n}}\right)^{n} d z_{n} \cdots d z_{1}
$$

where the sum is over all ordered partitions of $t$. This formula can also be proved "more directly" combinatorially. ${ }^{2}$

## 3. Small $\boldsymbol{n}$ Do Not Require a Computer

We now illustrate the computation of $H_{n}$ (and hence vol $\mathcal{B}_{n}$ ) by means of Theorem 2 for $n=3$ and 4 . These calculations "by hand" give an idea what computational tricks one might use in tackling larger $n$ with the aid of a computer.

By the theorem,

$$
\begin{aligned}
& H_{3}(t)=\frac{1}{(2 \pi i)^{3}} \int\left(z_{1} z_{2} z_{3}\right)^{-t-1}\left(\frac{z_{1}^{t+2}}{\left(z_{1}-z_{2}\right)\left(z_{1}-z_{3}\right)}+\frac{z_{2}^{t+2}}{\left(z_{2}-z_{1}\right)\left(z_{2}-z_{3}\right)}\right. \\
&\left.+\frac{z_{3}^{t+2}}{\left(z_{3}-z_{1}\right)\left(z_{3}-z_{2}\right)}\right)^{3} d \mathbf{z}
\end{aligned}
$$

We have to order the radii of the integration paths for each variable; we choose $0<\varepsilon_{3}<$ $\varepsilon_{2}<\varepsilon_{1}<1$. We heavily use this fact after multiplying out the cubic: integrating, for example, the term

$$
\frac{z_{1}^{-t-1} z_{2}^{-t-1} z_{3}^{2 t+5}}{\left(z_{3}-z_{2}\right)^{3}\left(z_{3}-z_{1}\right)^{3}}
$$

with respect to $z_{3}$ gives 0 , as this function is analytic at the $z_{3}$-origin and $\left|z_{1}\right|,\left|z_{2}\right|>\varepsilon_{3}$. After exploiting this observation for all the terms stemming from the cubic, the only integrals surviving are

$$
\frac{1}{(2 \pi i)^{3}} \int \frac{z_{1}^{2 t+5} z_{2}^{-t-1} z_{3}^{-t-1}}{\left(z_{1}-z_{2}\right)^{3}\left(z_{1}-z_{3}\right)^{3}} d \mathbf{z}
$$

[^1]and
$$
-\frac{3}{(2 \pi i)^{3}} \int \frac{z_{1}^{t+3} z_{2} z_{3}^{-t-1}}{\left(z_{1}-z_{2}\right)^{3}\left(z_{1}-z_{3}\right)^{2}\left(z_{2}-z_{3}\right)} d \mathbf{z} .
$$

The first integral factors and yields, again by residue calculus,

$$
\begin{aligned}
\frac{1}{(2 \pi i)^{3}} & \int \frac{z_{1}^{2 t+5} z_{2}^{-t-1} z_{3}^{-t-1}}{\left(z_{1}-z_{2}\right)^{3}\left(z_{1}-z_{3}\right)^{3}} d \mathbf{z} \\
& =\frac{1}{(2 \pi i)^{3}} \int z_{1}^{2 t+5}\left(\int \frac{z^{-t-1}}{\left(z_{1}-z\right)^{3}} d z\right)^{2} d z_{1} \\
& =\frac{1}{2 \pi i} \int z_{1}^{2 t+5}\left(-\frac{1}{2}(-t-1)(-t-2) z_{1}^{-t-3}\right)^{2} d z_{1} \\
& =\binom{t+2}{2}^{2} .
\end{aligned}
$$

For the second integral, it is most efficient to integrate with respect to $z_{2}$ first:

$$
\begin{aligned}
-\frac{3}{(2 \pi i)^{3}} & \int \frac{z_{1}^{t+3} z_{2} z_{3}^{-t-1}}{\left(z_{1}-z_{2}\right)^{3}\left(z_{1}-z_{3}\right)^{2}\left(z_{2}-z_{3}\right)} d \mathbf{z} \\
& =-\frac{3}{(2 \pi i)^{2}} \int \frac{z_{1}^{t+3} z_{3}^{-t}}{\left(z_{1}-z_{3}\right)^{5}} d z_{3} d z_{1} \\
& =-\frac{3}{2 \pi i} \int z_{1}^{t+3} \frac{1}{4!}(-t)(-t-1)(-t-2)(-t-3) z_{1}^{-t-4} d z_{1} \\
& =-3\binom{t+3}{4}
\end{aligned}
$$

Adding the last two lines gives finally

$$
H_{3}(t)=\binom{t+2}{2}^{2}-3\binom{t+3}{4}=\frac{1}{8} t^{4}+\frac{3}{4} t^{3}+\frac{15}{8} t^{2}+\frac{9}{4} t+1
$$

which is equal to (3). To obtain the volume of $\mathcal{B}_{3}$, the leading term of $H_{3}$ has to be multiplied by the relative volume of the fundamental domain of the sublattice of $\mathbb{Z}^{9}$ in the affine space spanned by $\mathcal{B}_{3}$. This volume is 9 ; hence

$$
\operatorname{vol} \mathcal{B}_{3}=\frac{9}{8} .
$$

In general, it is not hard to prove (see, for example, the appendix of [7]) that the relative volume of the fundamental domain of the sublattice of $\mathbb{Z}^{n^{2}}$ in the affine space spanned by $\mathcal{B}_{n}$ is $n^{n-1}$.

The number of integrals we have to evaluate to compute $H_{4}$ is only slightly higher. By Theorem 2,
$H_{4}(t)=\frac{1}{(2 \pi i)^{4}} \int_{\left|z_{1}\right|=\varepsilon_{1}} \int_{\left|z_{2}\right|=\varepsilon_{2}} \int_{\left|z_{3}\right|=\varepsilon_{3}} \int_{\left|z_{4}\right|=\varepsilon_{4}}\left(z_{1} z_{2} z_{3} z_{4}\right)^{-t-1}\left(\sum_{k=1}^{4} \frac{z_{k}^{t+3}}{\prod_{j \neq k}\left(z_{k}-z_{j}\right)}\right)^{4} d \mathbf{z}$.

Again we have a choice of ordering the radii; we use $0<\varepsilon_{4}<\varepsilon_{3}<\varepsilon_{2}<\varepsilon_{1}<1$. After multiplying out the quartic, we have to calculate five integrals; their evaluationagain straightforward by means of the residue theorem-is as follows. As before, we can "save" computation effort by choosing a particular order with which we integrate and by factoring an integral if possible:

$$
\begin{gathered}
\frac{1}{(2 \pi i)^{4}} \int \frac{z_{1}^{3 t+11} z_{2}^{-t-1} z_{3}^{-t-1} z_{4}^{-t-1}}{\left(z_{1}-z_{2}\right)^{4}\left(z_{1}-z_{3}\right)^{4}\left(z_{1}-z_{4}\right)^{4}} d \mathbf{z}=\frac{1}{(2 \pi i)^{4}} \int z_{1}^{3 t+11}\left(\int \frac{z^{-t-1}}{\left(z_{1}-z\right)^{4}} d z\right)^{3} d z_{1} \\
=\binom{t+3}{3}^{3} \\
-\frac{4}{(2 \pi i)^{4}} \int \frac{z_{1}^{2 t+8} z_{2}^{2} z_{3}^{-t-1} z_{4}^{-t-1}}{\left(z_{1}-z_{2}\right)^{4}\left(z_{1}-z_{3}\right)^{3}\left(z_{1}-z_{4}\right)^{3}\left(z_{2}-z_{3}\right)\left(z_{2}-z_{4}\right)} d \mathbf{z} \\
=-\frac{4}{(2 \pi i)^{4}} \int \frac{z_{1}^{2 t+8} z_{2}^{2}}{\left(z_{1}-z_{2}\right)^{4}}\left(\int \frac{z^{-t-1}}{\left(z_{1}-\right)^{3}\left(z_{2}-z\right)} d z\right)^{2} d z_{1} d z_{2} \\
=\frac{4}{(2 \pi i)^{2}} \int \frac{z_{1}^{2 t+8} z_{2}^{-t+1}}{\left(z_{1}-z_{2}\right)^{7}}\left(2\binom{t+2}{2} \frac{z_{1}^{-t-3}}{z_{1}-z_{2}}+2(t+1) \frac{z_{1}^{-t-2}}{\left(z_{1}-z_{2}\right)^{2}}\right. \\
\left.\quad+2 \frac{z_{1}^{-t-1}}{\left(z_{1}-z_{2}\right)^{3}}-\frac{z_{2}^{-t-1}}{\left(z_{1}-z_{2}\right)^{3}}\right) d \mathbf{z} \\
=8\binom{t+2}{2}\binom{t+5}{7}+8(t+1)\binom{t+6}{8}+8\binom{t+7}{9}-4\binom{2 t+8}{9}
\end{gathered}
$$

$$
\frac{4}{(2 \pi i)^{4}} \int \frac{z_{1}^{2 t+8} z_{2}^{-t-1} z_{3}^{2} z_{4}^{-t-1}}{\left(z_{1}-z_{2}\right)^{3}\left(z_{1}-z_{3}\right)^{4}\left(z_{1}-z_{4}\right)^{3}\left(z_{2}-z_{3}\right)\left(z_{3}-z_{4}\right)} d \mathbf{z}
$$

$$
=\frac{4}{(2 \pi i)^{3}} \int \frac{z_{1}^{2 t+8} z_{2}^{-t+1} z_{4}^{-t-1}}{\left(z_{1}-z_{2}\right)^{3}\left(z_{1}-z_{4}\right)^{7}\left(z_{2}-z_{4}\right)} d \mathbf{z}
$$

$$
=\frac{4}{(2 \pi i)^{2}} \int \frac{z_{1}^{2 t+8} z_{4}^{-t-1}}{\left(z_{1}-z_{4}\right)^{7}}\left(\binom{t+2}{2} \frac{z_{1}^{-t-3}}{z_{1}-z_{4}}+(t+1) \frac{z_{1}^{-t-2}}{\left(z_{1}-z_{4}\right)^{2}}+\frac{z_{1}^{-t-1}}{\left(z_{1}-z_{4}\right)^{3}}\right) d \mathbf{z}
$$

$$
=4\left(\binom{t+2}{2}\binom{t+5}{7}+(t+1)\binom{t+6}{8}+\binom{t+7}{9}\right)
$$

$$
\frac{6}{(2 \pi i)^{4}} \int \frac{z_{1}^{t+5} z_{2}^{t+5} z_{3}^{-t-1} z_{4}^{-t-1}}{\left(z_{1}-z_{2}\right)^{4}\left(z_{1}-z_{3}\right)^{2}\left(z_{1}-z_{4}\right)^{2}\left(z_{2}-z_{3}\right)^{2}\left(z_{2}-z_{4}\right)^{2}} d \mathbf{z}
$$

$$
=\frac{6}{(2 \pi i)^{4}} \int \frac{z_{1}^{t+5} z_{2}^{t+5}}{\left(z_{1}-z_{2}\right)^{4}}\left(\int \frac{z^{-t-1}}{\left(z_{1}-z\right)^{2}\left(z_{2}-z\right)^{2}} d z\right)^{2} d z_{1} d z_{2}
$$

$$
\begin{aligned}
& =\frac{6}{(2 \pi i)^{2}} \int \frac{z_{1}^{t+5} z_{2}^{t+5}}{\left(z_{1}-z_{2}\right)^{4}}\left((t+1)^{2} \frac{z_{2}^{-2 t-4}}{\left(z_{1}-z_{2}\right)^{4}}\right. \\
& \left.-4(t+1) \frac{z_{2}^{-2 t-3}}{\left(z_{1}-z_{2}\right)^{5}}+4 \frac{z_{2}^{-2 t-2}}{\left(z_{1}-z_{2}\right)^{6}}\right) d \mathbf{z} \\
& =6(t+1)^{2}\binom{t+5}{7}-24(t+1)\binom{t+5}{8}+24\binom{t+5}{9}, \\
& -\frac{12}{(2 \pi i)^{4}} \int \frac{z_{1}^{t+5} z_{2}^{2} z_{3}^{2} z_{4}^{-t-1}}{\left(z_{1}-z_{2}\right)^{3}\left(z_{1}-z_{3}\right)^{3}\left(z_{1}-z_{4}\right)^{2}\left(z_{2}-z_{3}\right)^{2}\left(z_{2}-z_{4}\right)\left(z_{3}-z_{4}\right)} d \mathbf{z} \\
& = \\
& - \\
& =-\frac{12}{(2 \pi i)^{3}} \int \frac{z_{1}^{t+5} z_{2}^{2} z_{4}^{t+1}}{(2 \pi i)^{2}} \int \frac{z_{1}^{t+5} z_{4}^{-t+1}}{\left(z_{1}-z_{4}\right)^{5}}\left(\frac{1}{\left(z_{1}-z_{4}\right)^{3}}+6 \frac{z_{4}}{\left(z_{1}-z_{4}\right)^{4}}+6 \frac{z_{4}^{2}}{\left(z_{1}-z_{4}\right)^{5}}\right) d \mathbf{z} \\
& = \\
& =-12\binom{t+5}{7}-72\binom{t+5}{8}-72\binom{t+5}{9} .
\end{aligned}
$$

Adding them all up gives

$$
\begin{aligned}
H_{4}(t)= & \frac{11}{11340} t^{9}+\frac{11}{630} t^{8}+\frac{19}{135} t^{7}+\frac{2}{3} t^{6}+\frac{1109}{540} t^{5} \\
& +\frac{43}{10} t^{4}+\frac{35117}{5670} t^{3}+\frac{379}{63} t^{2}+\frac{65}{18} t+1
\end{aligned}
$$

and hence

$$
\operatorname{vol} \mathcal{B}_{4}=4^{3} \cdot \frac{11}{11340}=\frac{176}{2835} .
$$

## 4. Larger $\boldsymbol{n}$ Do

As we have seen in the examples, after multiplying out the integrand of Theorem 2, many of the terms do not contribute to the integral. The following proposition will provide us with a general statement to that effect.

For a rational function $f$ in $n$ variables $z_{j}$ we use the notation $d_{r}(f)$ for the degree of $f$ in the variables $z_{1}, \ldots, z_{r}$.

Proposition 3. Suppose $p_{1}, \ldots, p_{n}$ are integers, $q_{j k}$ are nonnegative integers $(1 \leq$ $j<k \leq n), 1>\varepsilon_{1}>\cdots>\varepsilon_{n}>0$, and

$$
f\left(z_{1}, \ldots, z_{n}\right)=\frac{\prod_{1 \leq j \leq n} z_{j}^{p_{j}}}{\prod_{1 \leq j<k \leq n}\left(z_{j}-z_{k}\right)^{q_{j k}}} .
$$

If $1 \leq r \leq n$ and $d_{r}(f)<-r$, then

$$
\int_{\left|z_{1}\right|=\varepsilon_{1}} \cdots \int_{\left|z_{n}\right|=\varepsilon_{n}} f(\mathbf{z}) d \mathbf{z}=0
$$

Proof. We need only show that

$$
\begin{equation*}
\int_{\left|z_{1}\right|=\varepsilon_{1}} \cdots \int_{\left|z_{r}\right|=\varepsilon_{r}} f(\mathbf{z}) d \mathbf{z}=0 \tag{*}
\end{equation*}
$$

for then we can integrate over all $n$ variables by first integrating over $z_{1} \cdots z_{r}$.
If $r=1$, then $f$, considered as a function of $z_{1}$, is analytic outside the circle $\left|z_{1}\right|=\varepsilon_{1}$ and has zero residue at infinity since its degree is less than -1 , so $(*)$ follows.

We continue by induction, so suppose $r>1$ and $(*)$ is true for smaller $r$. Suppose that $d_{r}(f)<-r$; we may assume $d_{r-1}(f) \geq-(r-1)$. Note that $d_{r}(f) \geq d_{r-1}(f)+d$ where $d$ is the degree of $f$ in the single variable $z_{r}$ (the discrepancy is the sum of the exponents $q_{j r}$ for $j<r$ ). Hence $d \leq d_{r}(f)-d_{r-1}(f)<-1$. We consider $f$ as a function of $z_{r}$ and apply the residue theorem to the region outside the circle $\left|z_{r}\right|=\varepsilon_{r}$. As above $f$ has zero residue at infinity, so we only need to consider the residues at the poles $z_{j}$ for $j<r$. Evaluating these residues converts the integral of $f$ into a (possibly huge) linear combination of integrals of functions of the same form as $f$, but in the $n-1$ variables $z_{1}, \ldots, z_{r-1}, z_{r+1}, \ldots, z_{n}$. If $g$ is any one of these functions, then we easily calculate $d_{r-1}(g)=d_{r}(f)+1<-(r-1)$, and, by the induction hypothesis, the integral of $g$ is zero.

From this proposition we obtain the starting point for our algorithm.
Corollary 4. For $1>\varepsilon_{1}>\cdots>\varepsilon_{n}>0$ and $t \geq 0$,

$$
\begin{aligned}
H_{n}(t)= & \frac{1}{(2 \pi i)^{n}} \int_{\left|z_{1}\right|=\varepsilon_{1}} \cdots \int_{\left|z_{n}\right|=\varepsilon_{n}}\left(z_{1} \cdots z_{n}\right)^{-t-1} \\
& \times \sum_{m_{1}+\cdots+m_{n}=n}^{*}\binom{n}{m_{1}, \ldots, m_{n}} \prod_{k=1}^{n}\left(\frac{z_{k}^{t+n-1}}{\prod_{j \neq k}\left(z_{k}-z_{j}\right)}\right)^{m_{k}} d \mathbf{z}
\end{aligned}
$$

where $\sum^{*}$ denotes that we only sum over those n-tuples of nonnegative integers satisfying $m_{1}+\cdots+m_{n}=n$ and $m_{1}+\cdots+m_{r}>r$ if $1 \leq r<n$.

Remark. The condition on $m_{1}, \ldots, m_{n}$ can be visualized through lattice paths from $(0,0)$ to $(n, n)$ using the steps $\left(1, m_{1}\right),\left(1, m_{2}\right), \ldots,\left(1, m_{n}\right)$. The condition means that these paths stay strictly above the diagonal (except at the start and end). ${ }^{3}$

Proof. By Theorem 2,

$$
H_{n}(t)=\frac{1}{(2 \pi i)^{n}} \int \cdots \int\left(z_{1} \cdots z_{n}\right)^{-t-1}\left(\sum_{k=1}^{n} \frac{z_{k}^{t+n-1}}{\prod_{j \neq k}\left(z_{k}-z_{j}\right)}\right)^{n} d \mathbf{z}
$$

[^2]\[

$$
\begin{aligned}
= & \frac{1}{(2 \pi i)^{n}} \int \cdots \int\left(z_{1} \cdots z_{n}\right)^{-t-1} \\
& \times \sum_{m_{1}+\ldots+m_{n}=n}\binom{n}{m_{1}, \ldots, m_{n}} \prod_{k=1}^{n}\left(\frac{z_{k}^{t+n-1}}{\prod_{j \neq k}\left(z_{k}-z_{j}\right)}\right)^{m_{k}} d \mathbf{z} .
\end{aligned}
$$
\]

We select a partition $m_{1}+\cdots+m_{n}=n$ and rewrite the corresponding integrand in the language of Proposition 3:

$$
p_{j}=(n-1) m_{j}+t\left(m_{j}-1\right)-1, \quad q_{j k}=m_{j}+m_{k}
$$

Now suppose $1 \leq r<n$. The degree of the denominator in $z_{1}, \ldots, z_{r}$ is

$$
\begin{aligned}
\sum_{j=1}^{r} \sum_{k=j+1}^{n}\left(m_{j}+m_{k}\right) & =(n-1) \sum_{j=1}^{r} m_{j}+r \sum_{j=r+1}^{n} m_{j}=(n-1) \sum_{j=1}^{r} m_{j}+n r-r \sum_{j=1}^{r} m_{j} \\
& =(n-1) \sum_{j=1}^{r} m_{j}+n r-r \sum_{j=1}^{r}\left(m_{j}-1\right)-r^{2}
\end{aligned}
$$

We subtract this from $\sum_{j=1}^{r} p_{j}$ to get

$$
d_{r}(f)=(t+r) \sum_{j=1}^{r}\left(m_{j}-1\right)-r(n-r)-r .
$$

Here $t$ is nonnegative and $r(n-r)$ is positive, so Proposition 3 implies that the integral is zero unless $\sum_{j=1}^{r}\left(m_{j}-1\right)>0$.

Theoretically, Corollary 4 tells us what we have to do to compute $H_{n}$. Thus the integrals we computed in our calculations of $H_{3}$ and $H_{4}$ are exactly the nonzero integrals according to this result. For practical purposes, however, the statement is almost worthless for larger values of $n$. The first problem is that the number of terms in the sum equals the $(n-1)$ th Catalan number

$$
\frac{1}{n}\binom{2(n-1)}{n-1}=\frac{(2 n-2)!}{n!(n-1)!}
$$

which grows exponentially with $n$. Another slippery point is the evaluation of each integral. As we have seen in the examples, and as can be easily seen for the general case, we can compute each integral step by step one variable at a time. However, this means at each step we convert a rational function into a sum of rational functions (of one variable less) by means of the residue theorem. Again, this means that the number of (single-variable) integrals we have to compute grows immensely as $n$ increases. In fact, if we just "feed" the statement of Corollary 4 into a computer and tell it to integrate each summand, say, starting with $z_{1}$, then $z_{2}$, and so on, the computation time explodes once one tries $n=7$ or 8 . We feel that computationally this is as involved as calculating a sufficient number of values of $H_{n}$ and then interpolating this polynomial. However, complex analysis allows us some shortcuts which turn out to speed up the computation by a huge factor and which make Corollary 4 valuable, even from a computational
perspective. These "tricks" all showed up already in the examples and include

1. realizing when a function is analytic at the $z_{k}$-origin,
2. trying to choose the most efficient order of variables to integrate (based on estimating how many terms will be generated by the residue calculations, for each available variable), and
3. factoring the integral if some of the variables appear in a symmetric fashion.

Finally, if we are only interested in the volume of $\mathcal{B}_{n}$, we may also be
4. suppressing a particular integral if it does not contribute to the leading term of $H_{n}$.

It is worth noting that each of these computational "speed-ups" decreases the total computation time substantially. By applying them to Corollary 4, we implemented a C++ program for the specific functions we have to integrate to compute $H_{n}$. We were able to verify all previously known polynomials $(n \leq 8)$ and to compute vol $\mathcal{B}_{9}$ and $H_{9}$. The results of our calculations, including the polynomials $H_{n}$ for $n \leq 9$, and the source code for our program are available at www. math.binghamton. edu/dennis/Birkhoff. The following table gives the volumes together with the respective computing time (on a 1 GHz PC running under Linux). Note that computing the full polynomial $H_{n}$ takes longer, as we cannot make use of shortcut number 4. In fact, the computation of $H_{9}$ took about 325 days of computer time, although the elapsed time was only about 2 weeks since we distributed the calculations among a number of machines (8-40, depending on availability).

| $n$ | $\operatorname{vol} \mathcal{B}_{n}$ | Time |
| :--- | :---: | :---: |
| 1 | 1 | $<0.01 \mathrm{sec}$ |
| 2 | 2 | $<0.01 \mathrm{sec}$ |
| 3 | $\frac{9}{8}$ | $<0.01 \mathrm{sec}$ |
| 4 | $\frac{176}{2835}$ | $<0.01 \mathrm{sec}$ |
| 5 | $\frac{23590375}{167382319104}$ |  |
| 6 | $\frac{9700106723}{1319281996032 \cdot 10^{6}}$ |  |
| 7 | $\frac{77436678274508929033}{137302963682235238399868928 \cdot 10^{8}}$ |  |
| 8 | $\frac{55625533838576105333259507434329}{1258903626009547950081480942693339803308928 \cdot 10^{10}}$ | $<0.01 \mathrm{sec}$ |
| 9 | $\frac{559498129702796022246895686372766052475496691}{2153302766311808897781211017508325066061406891577233480945238016 \cdot 10^{14}}$ | 317 hr |

After this article was submitted for publication, we used the idle time on our departmental machines to compute the volume of $\mathcal{B}_{10}$. After a computation time of 6160 days, or almost

17 years (again scaled to a 1 GHz processor), we obtained

$$
\begin{aligned}
& \operatorname{vol} \mathcal{B}_{10} \\
& \quad=\frac{727291284016786420977508457990121862548823260052557333386607889}{828160860106766855125676318796872729344622463533089422677980721388055739956270293750883504892820848640000000}
\end{aligned}
$$

We hope that the reader will take this immense computing time as a challenge to improve our algorithms, and work towards $n=11$.

## 5. An Outlook Towards Transportation Polytopes

The Birkhoff polytopes are special cases of transportation polytopes, which are defined below. The study of this class of polytopes, which are naturally at least as fascinating as the Birkhoff polytopes, was motivated by problems in linear programming; for combinatorial properties see, for example, [10] and [13]. The goal of this section is to show how our methods can be applied in this more general setting.

Fix positive real numbers $a_{1}, \ldots, a_{m}, b_{1}, \ldots, b_{n}$ such that $a_{1}+\cdots+a_{m}=b_{1}+\cdots+b_{n}$. Let $\mathbf{a}=\left(a_{1}, \ldots, a_{m}\right), \mathbf{b}=\left(b_{1}, \ldots, b_{n}\right)$ :

$$
\mathcal{T}_{\mathbf{a}, \mathbf{b}}=\left\{\left(\begin{array}{ccc}
x_{11} & \cdots & x_{1 n} \\
\vdots & & \vdots \\
x_{m 1} & \ldots & x_{m n}
\end{array}\right) \in \mathbb{R}^{m n}: x_{j k} \geq 0, \sum_{k} x_{j k} x_{j k}=a_{j} \text { for all } 1 \leq j \leq m ~ b_{k} \text { for all } 1 \leq k \leq n\right\}
$$

is the set of solutions to the transportation problem with parameters $\mathbf{a}, \mathbf{b}$. It is a convex polytope of dimension $(m-1)(n-1)$ in $\mathbb{R}^{m n}$-hence we refer to $\mathcal{T}_{\mathbf{a}, \mathbf{b}}$ as a transportation polytope. Note that $\mathcal{B}_{n}=\mathcal{T}_{1,1}$ where $\mathbf{1}=(1, \ldots, 1) \in \mathbb{Z}^{n}$. We are exclusively interested in the case when $a_{1}, \ldots, a_{m}, b_{1}, \ldots, b_{n}$ are integers. One reason for this is that the vertices of $\mathcal{T}_{\mathbf{a}, \mathbf{b}}$ are then integral. Let

$$
T(\mathbf{a}, \mathbf{b})=T\left(a_{1}, \ldots, a_{m}, b_{1}, \ldots, b_{n}\right)=\#\left(\mathcal{T}_{\mathbf{a}, \mathbf{b}} \cap \mathbb{Z}^{m n}\right)
$$

denote the number of integer points in the transportation polytope with integral parameters $\mathbf{a}, \mathbf{b}$. Geometrically, each of these parameters determines the position of a hyperplane bounding the polytope $\mathcal{T}_{\mathbf{a}, \mathbf{b}}$. It is well known that $T(\mathbf{a}, \mathbf{b})$ is a piecewise-defined polynomial in $a_{1}, \ldots, a_{m}, b_{1}, \ldots, b_{n}$. The regions in which $T(\mathbf{a}, \mathbf{b})$ is a polynomial depends on the normal cones of the polytopes involved; for details, we refer to [21].

Moreover, one can derive results for the counting functions $T(\mathbf{a}, \mathbf{b})$ which are "higherdimensional" extensions of (4), by application of the main theorem in [3], which in turn is a generalization of the Ehrhart-Macdonald reciprocity theorem (2). To state this theorem, we need to define the following integer-point counting function. Suppose $\mathcal{P}$ is a rational polytope given in the form (5); denote by $\mathcal{P}(\mathbf{t})=\left\{\mathbf{x} \in \mathbb{R}_{\geq 0}^{d}: \mathbf{A x}=\mathbf{t}\right\}$ a polytope which we obtain from $\mathcal{P}=\mathcal{P}(\mathbf{b})$ by translating (some of) its bounding hyperplanes. (Classical Ehrhart dilation is the special case $\mathbf{t}=t \mathbf{b}$.) Let

$$
\mathcal{L}_{\mathcal{P}}(\mathbf{t})=\#\left(\mathcal{P}(\mathbf{t}) \cap \mathbb{Z}^{d}\right) .
$$

Theorem 4 of [3] states that $\mathcal{L}_{\mathcal{P}}$ and $\mathcal{L}_{\mathcal{P}}$ are piecewise-defined (multivariable) quasipolynomials satisfying

$$
\mathcal{L}_{\mathcal{P}}(-\mathbf{t})=(-1)^{\operatorname{dim} \mathcal{P}} \mathcal{L}_{\mathcal{P}^{\circ}}(\mathbf{t}) .
$$

As easily as (4) follows from (2), this reciprocity theorem yields a symmetry result for the transportation counting function. Let $\mathbf{1}_{d}$ denote the $d$-dimensional vector all of whose entries are one. Then

$$
T\left(-\mathbf{a}-n \mathbf{1}_{m},-\mathbf{b}-m \mathbf{1}_{n}\right)=(-1)^{m n-1} T(\mathbf{a}, \mathbf{b}) .
$$

We finally turn to the problem of writing $T(\mathbf{a}, \mathbf{b})$ in form of an integral. As with $\mathcal{B}_{n}$, we can view $\mathcal{T}_{\mathbf{a}, \mathbf{b}}$ as given in the form of (5) and apply the philosophy of Theorem 1 to obtain

$$
T(\mathbf{a}, \mathbf{b})=\frac{1}{(2 \pi i)^{m+n}} \int \cdots \int \frac{z_{1}^{-a_{1}-1} \cdots z_{m}^{-a_{m}-1} w_{1}^{-b_{1}-1} \cdots w_{n}^{-b_{n}-1}}{\prod_{\substack{1 \leq j \leq m \\ 1 \leq k \leq n}}\left(1-z_{j} w_{k}\right)} d \mathbf{z} d \mathbf{w}
$$

Again the integral with respect to each (one-dimensional) variable is over a circle with radius $<1$ centered at 0 , and all appearing radii are different. As in the Birkhoff case, we can separate, say, the w-variables and obtain

$$
T(\mathbf{a}, \mathbf{b})=\frac{1}{(2 \pi i)^{m+n}} \int \cdots \int z_{1}^{-a_{1}-1} \cdots z_{m}^{-a_{m}-1} \prod_{k=1}^{n} \int \frac{w_{k}^{-b_{k}-1}}{\prod_{j=1}^{m}\left(1-z_{j} w_{k}\right)} d w_{k} d \mathbf{z}
$$

As with the Birkhoff polytope, the innermost integrals are easy to compute by means of the residue theorem. This yields the following statement which generalizes Theorem 2.

Theorem 5. For any distinct $0<\varepsilon_{1}, \ldots, \varepsilon_{m}<1$,

$$
T(\mathbf{a}, \mathbf{b})=\frac{1}{(2 \pi i)^{m}} \int_{\left|z_{1}\right|=\varepsilon_{1}} \cdots \int_{\left|z_{m}\right|=\varepsilon_{m}} z_{1}^{-a_{1}-1} \cdots z_{m}^{-a_{m}-1} \prod_{k=1}^{n} \sum_{j=1}^{m} \frac{z_{j}^{b_{k}+m-1}}{\prod_{l \neq j}\left(z_{j}-z_{l}\right)} d \mathbf{z}
$$

Remark. As with Theorem 2, it can be proved from the form of the integrand that $T(\mathbf{a}, \mathbf{b})$ is indeed a piecewise-defined polynomial in $a_{1}, \ldots, a_{m}, b_{1}, \ldots, b_{n}$.

Theoretically we could now use this theorem to produce formulas for $T(\mathbf{a}, \mathbf{b})$ just as we did for $H_{n}(t)$. There is one major difference: $T(\mathbf{a}, \mathbf{b})$ is only a piecewise-defined polynomial. In fact, we can see this from the form of the integral in Theorem 5: whether we will get a nonzero contribution at a certain step in the computation depends heavily on the relationship between $a_{1}, \ldots, a_{m}, b_{1}, \ldots, b_{n}$.

The fact that the counting function $T(\mathbf{a}, \mathbf{b})$ is of a somewhat more delicate nature naturally has computational consequences. We feel that providing any general results on this function would go beyond the scope of this article and will hopefully be the subject of a future project. On the other hand, we adjusted our algorithm to compute values of $T(\mathbf{a}, \mathbf{b})$ for three (fixed) pairs ( $\mathbf{a}, \mathbf{b}$ ), which have been previously computed by Mount
[17] and DeLoera and Sturmfels [9]. The first example is

$$
\begin{aligned}
& T((3046,5173,6116,10928),(182,778,3635,9558,11110)) \\
& \quad=23196436596128897574829611531938753 .
\end{aligned}
$$

The authors reported their computation took 20 minutes (Mount) or 10 minutes (DeLoeraSturmfels) [9], [17]. We computed this number in 0.2 seconds, based on Theorem 5. A similar phenomenon happens with
$T((338106,574203,678876,1213008),(20202,142746,410755,1007773,1222717))$
$=316052820930116909459822049052149787748004963058022997262397$
and

$$
\begin{aligned}
& T((30201,59791,70017,41731,58270),(81016,68993,47000,43001,20000)) \\
& =2464053826815198108639701803342226405075725113340175811250949 \\
& 5633028,
\end{aligned}
$$

which reportedly took 35 minutes/ 10 days with the DeLoera-Sturmfels algorithm [9], $0.3 / 2.9$ seconds with ours. These timing comparisons ignore differences in machine speeds and implementation of the algorithms, but suggest that our methods are considerably more efficient.

## Acknowledgments

We thank the referees for carefully reading through our paper and software, and for their many helpful comments. We also thank Dean Hickerson for finding a missing factor in our original calculation of the volume of $\mathcal{B}_{9}$.

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Received February 25, 2002, and in revised form August 16, 2002. Online publication August 6, 2003.


[^0]:    ${ }^{1}$ E-mail to the first author on Wed., 14 August 2002, 16:29:35 (EDT).

[^1]:    ${ }^{2}$ The authors thank Sinai Robins and Frank Sottile for their help in the proof of this equivalence and its combinatorial interpretation.

[^2]:    ${ }^{3}$ The authors thank Lou Billera for pointing out this lattice-path interpretation.

