Nearness in Digital Images and Proximity Spaces[★]

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Abstract. The concept of “nearness”, which has been dealt with as soon as one started studying digital images, finds one of its rigorous forms in the notion of proximity space. It is this notion, together with “nearness preserving mappings”, that we investigate in this paper. We first review basic examples as they naturally occur in digital topologies, making also brief comparison studies with other concepts in digital geometry. After this we characterize proximally continuous mappings in metric spaces. Finally, we show by example that the “proximate complexity” of a finite covering in a digital picture may be too high to be adequately depicted in a finite topological space. This combinatorial result may indicate another conceptual advantage of proximities over topologies.

1 Introduction

Recently there have been quite an intense investigation of topological structures in image processing, mostly in connection with the analysis of connectivity and the operation of thinning (see e.g. [2,4,9,13], etc.). An interesting attempt to introduce richer structures than those of topology, and replacing thus “local” continuity properties by a global notion of nearness, has been done in [12] where the authors contemplated the so called semi-proximity spaces as a theoretical tool in the image processing studies. In this paper we want to go on in a similar vain by shedding light on certain questions which implicitly announced themselves in the paper [12], and by complementing the results of [12] with some new findings. (Our investigation here is relatively technical. The reader is supposed to be acquainted with the motivation for investigating digital topology and thus for the selection of the problems we pursue. Reading the papers [8,12,18] would

[★] This work was supported by the Austrian Science Foundation under grant S 7002-MAT, by the project of Czech Ministry of Education No. VS96049, and by the Grant Agency of the Czech Republic under the grant GACR 102/00/1679.

certainly be instrumental as it was for the authors. We hope, however, that our exposition is reasonably self-contained.)

The paper is organized as follows: In the section 2 the proximity spaces are introduced and related to the topological spaces, in particular to the topology introduced by Marcus and Wyse. Section 3 investigates the proximities in metric spaces. In the section 4 we discuss the problem whether nearness of finite partitions allows for a suitable finite topology on the index set. The conclusion summarizes the paper.

2 Proximity Spaces – In Relation to Topologies and Metrics

We shall deal with proximity spaces as implicitly defined by Riesz [16] and intensely pursued by several authors ([3,5,6,14]). Our definition is a modified version of the definition by E. Čech [3].

Definition 1 (Proximity Space). *A pair (X, π) is called a proximity space if X is a set and π is a binary relation on the power set of X , $\exp X$, which is subject to the following four requirements ($A, B, C, D, E \in \exp X$, the symbol $A \text{ non} \pi B$ means that $A \pi B$ is false):*

- (i) $(A \cup B) \pi C \Leftrightarrow A \pi C \text{ or } B \pi C$,
- (ii) $A \pi B \Rightarrow A \neq \emptyset \text{ and } B \neq \emptyset$,
- (iii) $A \cap B \neq \emptyset \Rightarrow A \pi B$,
- (iv) $A \pi B \Rightarrow B \pi A$.

If (X, π) fulfils the following additional requirement (v), it is called a *T-proximity*:

- (v) if $\{a\} \text{ non} \pi B$, then there is a set A such that $\{a\} \text{ non} \pi (X - A)$ and $\{d\} \text{ non} \pi B$ for each $d \in A$.

The axioms of a proximity space reflect the properties we observe when we consider the common-sense nearness. It should be noted that the less plausible axiom of T-proximity – the axiom (v) – guarantees that the proximity in question induces a topology (the proximities induce only the closure spaces – see [3] and [12]). There is an important link of proximity spaces and topological spaces.

Theorem 1 (Proximity Space, Topological Space).

- (i) Let (X, π) be a proximity (resp. T-proximity) space. Let $\bar{A} = \{x \in X \mid \{x\} \pi A\}$. Then $(X, \bar{})$ is a closure (resp. topological) space (the closure $\bar{}$ is said to be induced by the proximity π).
- (ii) Let $(X, \bar{})$ be a closure (resp. topological) space and let $(X, \bar{})$ fulfil the following condition (the Čech R_0 condition): If $a, b \in X$, then $a \in \{\bar{b}\}$ if and only if $b \in \{\bar{a}\}$. Let $A \pi B \Leftrightarrow \bar{A} \cap \bar{B} \neq \emptyset$. Then (X, π) is a proximity (resp. T-proximity) space and, moreover, (X, π) induces $(X, \bar{})$. $\square$

The proof of the above result is slightly technical (though essentially easy – one would use Th.23.B3 and Th.25.B9 of Čech [3]) and we omit it.

By the previous result, we can associate closures (resp. topologies) to proximities (resp. T-proximities) although not all topologies are proximizable. On the other hand, observe that many proximities on X may induce the same topology on X (and, vice versa, a topology may give rise to several proximities which induce it). For the reader's intuition, let us note that e. g. the following two proximities p_1, p_2 on an infinite set X define the same (discrete) topology: $A p_1 B \Leftrightarrow A \cap B \neq \emptyset$, $A p_2 B \Leftrightarrow$ either $A \cap B \neq \emptyset$ or both the sets A and B are infinite.

It is possible to adopt the notion of proximity as primary and view the notion of topology as secondary. One of the reasons for doing it is that the topology describes only *the local character of points* in a picture (or, if we want, the proximity of points and sets). But if we are to treat geometrical qualities of a picture – a situation which typically arises in image processing – it is *the proximity of sets* which matters most. E.g. in the case of scanned text documents the characters are the connected components of the black pixels of the thresholded image. The pixel sets of adjacent characters of words are closer to each other than between the words.

Thus, besides the concrete reason obvious from the study of thinning and shape deformation [12], the general reason for investigating the proximity relation lies in its fundamental role in all kinds of geometrically oriented considerations.

In view of Theorem 1 and the investigation of connectedness in digital images [20], the following proximity in Z^2 is worth recording. Recall that if $[r, s] \in Z^2$, then by the 4-neighbourhood of $[r, s]$ we mean the set $\{[r, s], [r \pm 1, s], [r, s \pm 1]\}$.

Definition 2 (Marcus–Wyse proximity). *Let Z^2 denote the subset of R^2 consisting of all points with the integer coordinates. Let $A \subset Z^2, B \subset Z^2$. Let us write $A \pi B$ if either $A \cap B \neq \emptyset$ or there exists a point $p = [r, s] \in A \cup B$ such that the following assertion holds true:*

- (i) $r + s$ is an odd number,
- (ii) if $p \in A$, then the 4-neighbourhood of p intersects B ,
- (iii) if $p \in B$, then the 4-neighbourhood of p intersects A .

The result of [20] (see also [11]) can be now reformulated proximity-wise. Following [12], let us agree to say that a set S in a proximity space (P, π) is connected if S cannot be written in the way $S = A \cup B$, where A, B are two non-proximal sets (i.e. $A \text{ non } \pi B$).

Theorem 2. *The graph-theoretic connectedness in Z^2 induced by the 4-neighbourhood adjacency relation coincides with the proximate connectedness induced by the Marcus–Wyse proximity.*

Let us take up proximities on metric spaces (here all proximities will automatically be T-proximities). Let us consider a metric space, (M, ρ) . Then (M, ρ)

can be viewed as a topological space with the closure $\bar{A} = \{x \in M \mid \rho(x, A) = 0\}$. In this way the metric space (M, ρ) induces a “topological” proximity, π_t , defined as follows: $A \pi_t B \Leftrightarrow \bar{A} \cap \bar{B} \neq \emptyset$. But (M, ρ) also induces a metric proximity, π_m . Write, for two subsets A and B , $\rho(A, B) = \inf_{a \in A, b \in B} \rho(a, b)$ and set $A \pi_m B \Leftrightarrow \rho(A, B) = 0$. It is easily seen that if $A \pi_t B$, then $A \pi_m B$ but not necessarily the other way round. If, for instance, A is the graph in R^2 of the function $f(x) = \frac{1}{x}$ and B is the x -axis in R^2 , then $A \pi_m B$ with respect to the Euclidean metric, but not $A \pi_t B$.

For a reader not trained in topology, let us explicitly clarify why (and when) the latter phenomenon (of the difference of π_t and π_m) may occur. Recall that a metric space M is called *compact* ([3, 5], etc.) if each sequence in M allows for a convergent subsequence in M . As known, a subset P of R^n endowed with the Euclidean metric (taken from R^n) is compact if and only if it is closed and bounded.

Theorem 3. *Let (M, ρ) be a compact metric space. Then $\pi_t = \pi_m$, i.e. in compact metric spaces the topological proximity agrees with the metric proximity.*

Proof. Suppose that $A, B \subset M$. It only remains to show that if $A \pi_m B$, then $A \pi_t B$, the other implication is always valid. Let $A \pi_m B$. We are to show that $\bar{A} \cap \bar{B} \neq \emptyset$. The relation $A \pi_m B$ means $\rho(A, B) = 0$. It follows that for each $n \in N$ we can find points $a_n \in A$, $b_n \in B$ such that $\rho(a_n, b_n) \leq \frac{1}{n}$. By compactness, there is a subsequence, a_{n_k} , of $\{a_n\}$ which converges to some element $a \in M$, and there is a subsequence, b_{n_k} , of $\{b_{n_k}\}$ which converges to some $b \in M$. Obviously, $a = b$ and therefore the sequences $\{a_{n_k}\}$ and $\{b_{n_k}\}$ converge to a common element $a (= b)$. Since $a \in \bar{A}$ and $b \in \bar{B}$ and since $a = b$, we see that $\bar{A} \cap \bar{B} \neq \emptyset$. This completes the proof.

The morphisms in the category of proximity spaces are the proximally continuous mappings.

Definition 3 (proximally continuous mapping). *Let (X_1, π_1) , (X_2, π_2) be proximity spaces. A mapping $f: X_1 \rightarrow X_2$ is called proximally continuous if the following implication is true:*

If $A, B \subset X_1$ and $A \pi_1 B$, then $f(A) \pi_2 f(B)$.

Thus, proximally continuous mappings are those mappings which preserve proximity. It is easily seen that a proximally continuous mapping is automatically continuous when understood as a mapping between the respective topological spaces induced by proximities. A continuous mapping does not have to be proximally continuous even if we consider it with respect to the metric proximity.

Example: Consider the function $f(x) = x^2: (0, +\infty) \rightarrow (0, +\infty)$. This mapping is obviously continuous but not proximally continuous. Indeed, let p_n be such a sequence that $(p_n + \frac{1}{2})^2 - p_n^2 \geq 1$. Let $A = \{p_n \mid n \in N\}$ and $B = \{p_n + \frac{1}{n} \mid n \in N\}$. Then $A \pi_m B$ but $f(A) \not\pi_m f(B)$.

3 Metric Proximities

In this section we are going to prove that the metric proximity of a space determines, up to a metric equivalence, the metric of the space. This relatively deep result has found applications in a number of geometric problems (see e. g. [5] for relevant comments). This result can be expressed in terms of “small” (=countable) subsets of the metric space in question and therefore it may have bearing on digital geometry. Also, we see that metric considerations of digital pictures ([17,18,19]) have a proximity character (i.e. can be expressed in proximity terms).

Let us take up the proof of the result. We provide a simplified transparent proof based on elementary reasonings only. We also point out other features of metric spaces relevant to proximity. Recall first two standard definitions.

Definition 4. Let (M, ρ) be a metric space and let $\{a_n\}$ be a sequence in M . We say that $\{a_n\}$ is a **Cauchy subsequence** if for each $\varepsilon > 0$ there is $n_0 \in \mathbb{N}$ such that $\rho(x_m, x_n) \leq \varepsilon$ provided $n \geq n_0, m \geq n_0$. We say that a sequence $(x_n)_{n \in \mathbb{N}}$ in M is **metrically discrete** (of order ε) if for any $n, m \in \mathbb{N}$ we have $\rho(x_m, x_n) \geq \varepsilon$.

The following proposition is essential in our argument. It may be interesting in its own right. Before we formulate it, let us agree to call the set $B_r(a) = \{b \in M \mid \rho(a, b) \leq r\}$ the r -ball around a .

Proposition 1 (Sequence principle in metric spaces). *Each sequence in a metric space contains either a Cauchy sequence or a metrically discrete subsequence.*

Proof. Take a sequence in a metric space and form the collection of all 1-balls centered at each of its points. If each of these balls contains only finitely many points of the sequence, we can easily construct a subsequence of the given sequence which is metrically discrete of order 1. If not, there is a 1-ball around a point of the sequence which contains infinitely many points of the sequence. Take these points and form the collection of all $\frac{1}{2}$ -balls centered at these points. If each of these $\frac{1}{2}$ -balls contains only finitely many points, we can easily construct a metrically discrete subsequence of order $\frac{1}{2}$. If not, there is a point such that the $\frac{1}{2}$ -ball around it contains infinitely many points. Going on this way inductively, either the procedure stops at the n -th step and we have an $\frac{1}{n}$ -discrete subsequence, or we obviously obtain a Cauchy subsequence of the given sequence.

We shall need one more metric notion. Let us recall it together with proximal continuity in metric space.

Definition 5. Let $(M_1, \rho_1), (M_2, \rho_2)$ be metric spaces. Let $f: M_1 \rightarrow M_2$ be a mapping. In accord with our general definition, we say that f is **proximally continuous** if the following property is fulfilled:

If P and Q are subsets of M_1 such that $\rho_1(P, Q) = 0$, then $\rho_2(f(P), f(Q)) = 0$.

We say that f is **metrically continuous** if for any $\varepsilon > 0$ there exists $\delta > 0$ such that whenever $\rho_1(x, y) < \delta$, then $\rho_2(f(x), f(y)) < \varepsilon$ (such a mapping f is sometimes called *uniformly continuous*).

Let us now formulate and prove the main result of this section. In the effort to make the proof accessible for nonspecialists in topology, we use only elementary reasonings. The novelty seems to be the utilization of the sequence principle as established in Prop. 1.

Theorem 4. *A mapping between metric spaces is proximally continuous if and only if it is metrically continuous.*

A consequence: If two metrics on a set induce the same proximity, they have to be (metrically) equivalent.

Proof. A metrically continuous mapping between metric spaces is obviously proximally continuous. Let us take up the nontrivial implication of the theorem.

Let $f: M_1 \rightarrow M_2$ be proximally continuous. Suppose f is not metrically continuous. It means that for some $\varepsilon > 0$ there exist sequences $(a_n)_{n \in N}$ and $(b_n)_{n \in N}$ in M_1 such that $\rho_1(a_n, b_n) \rightarrow 0$ whereas $\rho_2(f(a_n), f(b_n)) \geq \varepsilon$. Let us look at the sequence $f(a_n)_{n \in N}$. If it contains a Cauchy subsequence, then there is an infinite subset of this sequence which is all contained in an $\frac{\varepsilon}{4}$ -ball. Let us denote this subsequence by $f(a_{n_k})_{k \in N}$. Then the sets $\{f(a_{n_k}): k \in N\}$ and $\{f(b_{n_k}): k \in N\}$ are obviously $\frac{\varepsilon}{2}$ -apart (the triangular inequality), which is absurd. If there is a Cauchy subsequence of $f(b_n)_{n \in N}$, we can apply an analogous reasoning. Suppose that neither of the former two cases applies. By the sequence principle, we can easily construct metrically discrete subsequences $f(a_{n_k})_{k \in N}$ and $f(b_{n_k})_{k \in N}$ of $f(a_n)_{n \in N}$ and $f(b_n)_{n \in N}$, respectively. For simplicity, let us denote them again by $f(a_n)_{n \in N}$ and $f(b_n)_{n \in N}$.

Let us assume that these sequences are metrically discrete of order α . Define $r = \frac{1}{2} \min\{\varepsilon, \alpha\}$ and form the collection of all the r -balls centered at $f(a_n)_{n \in N}$. Note that each of them contains at most one of the elements $f(b_n)_{n \in N}$. Let us now proceed inductively. Put $n_1 = 1$. Take $n_2 \in N$, $n_2 > 1$ such that

$n_2 > h$ if $f(b_h)$ belongs to the r -ball centered at $f(a_1)$, and

$n_2 > k$ if the r -ball centered at $f(a_k)$ contains $f(b_1)$,

(if neither h or k exists, we simply take $n_2 = n_1 + 1$). By induction, given $n_1 < n_2 < \dots < n_i$, take $n_{i+1} \in N$, $n_{i+1} > n_i$ such that

$n_{i+1} > h$ if $f(b_h)$ belongs to the r -ball centered at $f(a_{n_i})$, and

$n_{i+1} > k$ if the r -ball centered at $f(a_k)$ contains $f(b_{n_i})$.

Note that, by our construction, if $j \in \{n_k | k \in N\}$, then $f(b_j)$ does not belong to any of the r -balls centered in $f(a_{n_k})$. The sets $f(a_{n_k})_{k \in N}$ and $f(b_{n_k})_{k \in N}$ are therefore apart of the order r . Write $A = \{a_{n_k}\}_{k \in N}$, $B = \{b_{n_k}\}_{k \in N}$. The $\rho_1(A, B) = 0$ but $\rho_1(f(A), f(B)) \geq r$. This means that f is not proximally continuous - a contradiction. This completes the proof.

The main consequence of the latter theorem (and its proof) as far as the potential application in image processing goes is contained in the next theorem. It does not seem to be explicitly formulated in the literature but it is this formulation which may have relevant bearing on the digital images studies - it reduces proximities of general sets to proximities of small (=countable, digitally

accessible) sets. (Recall that a set is said to be countable if it has the smallest possible infinite cardinality, i.e. if it has the cardinality of natural numbers.)

Theorem 5. *Let ρ_1, ρ_2 be two metrics on a set M . If $\rho_1(A, B) = 0 \Leftrightarrow \rho_2(A, B) = 0$ for all **countable** subsets of M , then the metrics ρ_1, ρ_2 are metrically equivalent. In other words, if the proximities given by ρ_1 and ρ_2 agree when restricted to countable subsets of M , then ρ_1 and ρ_2 are metrically indistinguishable.*

4 Near and Far Sets in a Finite Partition – Could “nearness” Be Controlled by a Finite Topology?

Each picture can be viewed as a partition of the underlying set into a finite family of sets. The main point in understanding the partition is specifying which sets are “near” (proximal) and which are “far” (non-proximal). In this section we exhibit an example which shows that the proximal relation of sets in a finite partition may be too complex to be described with the help of the notion of finite topology. This relates in a natural way our investigation here with topological studies in image processing (see [13, 8, 11, 18, 10, 20], etc.). (Recall that by a finite partition of a set S we mean a mutually disjoint collection $\{S_k \ (k = 1, 2, \dots, n)\}$

of non-empty subsets of S such that $\bigcup_{k=1}^n S_k = S$.)

Let us now introduce an auxiliary notion.

Definition 6. *Let (X, π) be the metric proximity space of the metric space (M, ρ) . Let $P = \{S_k \mid k = 1, 2, \dots, n\}$ be a partition of X . Let us say that P is **controlled by topology** if there is a topology t on the set $T = \{1, 2, \dots, n\}$ such that $S_k \pi S_\ell$ if and only if the set $\{k, \ell\}$ is topologically connected as a topological subspace of T .*

If each partition could be topologically controlled, which seems conceivable at first sight, we would find ourselves in an advantages situation in view of the understanding of finite topologies (see e.g. [18] and [10]). However, it is not necessarily the case. We will illustrate it by a simple example. Before, observe that we can restrict ourselves to open sets of the partition if we allow for a small degree of overlapping (in practice, we have to allow for it anyhow in view of the imperfection of our measurement).

Proposition 2. *Let (M, ρ) be a metric space. Let $\{S_k \mid k = 1, 2, \dots, n\}$ be a partition of M . Then there is an $\varepsilon, \varepsilon > 0$ such that the following statement holds true: If $T_k = \{x \in M \mid \rho(x, S_k) < \varepsilon\}, k = 1, 2, \dots, n$, then each T_k is an open subset of M , and $T_k \cap T_\ell \neq \emptyset$ if and only if $S_k \pi_\rho S_\ell$.*

Proof. All sets $T_k, k = 1, \dots, n$ are obviously open. Further, consider all couples S_k, S_ℓ such that $S_k \text{ non } \pi_\rho S_\ell$. Since there are only finitely many couples with this property, there is a sufficiently small $\varepsilon, \varepsilon > 0$ such that $T_k \cap T_\ell = \emptyset$

Fig. 1.

Fig. 2.

for all couples of T_k, T_ℓ corresponding to the couples S_k, S_ℓ considered. Since we have $S_k \pi_m S_\ell$ for the remaining couples, we infer from the definition of metric proximity that the neighbourhood sets T_k, T_ℓ , corresponding to these S_k, S_ℓ fulfil $T_k \cap T_\ell \neq \emptyset$. The proof is complete.

We see from the above proposition that we can simply construct the counterexample on the disjointness – intersection basis of an open covering.

Example: Consider the Fig. 1 above. Let $\{B_i | i \leq 6\}$ be the family of balls in R^2 as drawn in Fig. 1. Then the proximity of the family $\{B_i | i \leq 6\}$ cannot be topologically controlled.

Indeed, considering the couples of balls which are disjoint (resp. which overlap), we see that the existence of the required topology would amount to having a topology on the set $\{1, 2, 3, 4, 5, 6\}$ such that precisely the following pairs would constitute connected subspaces: $\{1, 2\}, \{2, 3\}, \{2, 5\}, \{3, 4\}, \{3, 5\}$, and $\{5, 6\}$. This would mean the existence of topology on the Fig. 2 above which would produce the 8-adjacency connectedness but this is impossible as shown in [13] and [11].

5 Conclusions

We suggest that the notion of proximity space might be a useful tool for theoretical studies in image processing. We exhibit basic examples and we link them with previous topological investigations (for instance, with the Marcus–Wyse topologies). Then we analyze more thoroughly the metric proximities. As a main result we show that the metric proximity of small sets determines in a way the metric of the underlying space. The interpretation of this result reads, roughly, that if we can verify which sequences of points in a metric space are proximal and which are not, we can in a certain sense recover the metric. In the end we show that a topological result known from the investigation of 8-adjacency relation in Z^2 disproves a conjecture about “topologizing” proximities in a partition.

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