Outlets of 2D invasion percolation and multiple-armed incipient infinite clusters

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Abstract We study invasion percolation in two dimensions, focusing on properties of the outlets of the invasion and their relation to critical percolation and to incipient infinite clusters (IICs). First we compute the exact decay rate of the distribution of both the weight of the kth outlet and the volume of the kth pond. Next we prove bounds for all moments of the distribution of the number of outlets in an annulus. This result leads to almost sure bounds for the number of outlets in a box $B(2^n)$ and for the decay rate of the weight of the kth outlet to p_c . We then prove existence of multiple-armed IIC measures for any number of arms and for any color sequence which is alternating or monochromatic. We use these measures to study the invaded region near outlets and near edges in the invasion backbone far from the origin.

Keywords Invasion percolation \cdot Invasion ponds \cdot Critical percolation \cdot Near critical percolation \cdot Correlation length \cdot Scaling relations \cdot Incipient infinite cluster

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1 Introduction

1.1 The model

Invasion percolation is a stochastic growth model both introduced and numerically studied independently by [1,14]. Let G=(V,E) be an infinite connected graph in which a distinguished vertex, the origin, is chosen. Let $(\tau_e)_{e\in E}$ be independent random variables, uniformly distributed on [0, 1]. The *invasion percolation cluster* (IPC) of the origin on G is defined as the limit of an increasing sequence (G_n) of connected subgraphs of G as follows. For an arbitrary subgraph G'=(V',E') of G, we define the outer edge boundary of G' as

$$\Delta G' = \{ e = \langle x, y \rangle \in E : e \notin E', \text{ but } x \in V' \text{ or } y \in V' \}.$$

We define G_0 to be the origin. Once the graph $G_i = (V_i, E_i)$ is defined, we select the edge e_{i+1} that minimizes τ on ΔG_i . We take $E_{i+1} = E_i \cup \{e_{i+1}\}$ and let G_{i+1} be the graph induced by the edge set E_{i+1} . The graph G_i is called the *invaded region* at time i. Let $E_{\infty} = \bigcup_{i=0}^{\infty} E_i$ and $V_{\infty} = \bigcup_{i=0}^{\infty} V_i$. Finally, define the IPC $\mathcal{S} = (V_{\infty}, E_{\infty})$.

In this paper, we study invasion percolation on two-dimensional lattices; however, for simplicity we restrict ourselves hereafter to the square lattice \mathbb{Z}^2 and denote by \mathbb{E}^2 the set of nearest-neighbor edges. The results of this paper still hold for lattices which are invariant under reflection in one of the coordinate axes and under rotation around the origin by some angle. In particular, this includes the triangular and honeycomb lattices.

We define Bernoulli percolation using the random variables τ_e to make a coupling with the invasion immediate. For any $p \in [0, 1]$ we say that an edge $e \in \mathbb{E}^2$ is p-open if $\tau_e < p$ and p-closed otherwise. It is obvious that the resulting random graph of p-open edges has the same distribution as the one obtained by declaring each edge of \mathbb{E}^2 open with probability p and closed with probability p is the probability of the state of all other edges. The percolation probability p is the probability that the origin is in the infinite cluster of p-open edges. There is a critical probability $p_c = \inf\{p : \theta(p) > 0\} \in (0, 1)$. For general background on Bernoulli percolation we refer the reader to p-open edges.

In [3], it was shown that, for any $p>p_c$, the invasion on $(\mathbb{Z}^d,\mathbb{E}^d)$ intersects the infinite p-open cluster with probability one. The definition of the invasion mechanism implies that if the invasion reaches the p-open infinite cluster for some p, it will never leave this cluster. Combining these two facts yields that if e_i is the edge added at step i then $\limsup_{i\to\infty}\tau_{e_i}=p_c$. It is well-known that for Bernoulli percolation on $(\mathbb{Z}^2,\mathbb{E}^2)$, the percolation probability at p_c is 0. This implies that, for infinitely many values of i, the weight $\tau_{e_i}>p_c$. The last two results give that $\hat{\tau}_1=\max\{\tau_e:e\in E_\infty\}$ exists and is greater than p_c . The above maximum is attained at an edge which we shall call \hat{e}_1 . Suppose that \hat{e}_1 is invaded at step i_1 , i.e. $\hat{e}_1=e_{i_1}$. Following the terminology of [16], we call the graph G_{i_1-1} the *first pond* of the invasion, denoting it by the symbol \hat{V}_1 , and we call the edge \hat{e}_1 the *first outlet*. The second pond of the invasion is defined similarly. Note that a simple extension of the above argument implies that $\hat{\tau}_2=\max\{\tau_{e_i}:e_i\in E_\infty,i>i_1\}$ exists and is greater than p_c . If we assume that $\hat{\tau}_2$



is taken on the edge \hat{e}_2 at step i_2 , we call the graph $G_{i_2-1} \setminus G_{i_1-1}$ the *second pond* of the invasion, and we denote it \hat{V}_2 . The edge \hat{e}_2 is called the *second outlet*. The further ponds \hat{V}_k and outlets \hat{e}_k are defined analogously. For a hydrological interpretation of the ponds we refer the reader to [20].

In this paper we study the sequence of outlets (\hat{e}_k) and the sequence of their weights $(\hat{\tau}_k)$. In Theorem 1.1 we give the asymptotic behaviour for the distribution of $\hat{\tau}_k$ for any fixed k. For k>1, we compute the exact decay rate of the distribution of the size of the kth pond in Theorem 1.2. This result can be also seen as a statement about the sequence of steps i_k at which \hat{e}_k are invaded. In Theorem 1.3, we find uniform bounds on all moments of the number of outlets in an annulus. We use this result in Theorem 1.4 to derive almost sure bounds on the number of outlets in a box $B(2^n)$. An important consequence of Theorem 1.4 is Corollary 1.1; it states almost sure bounds on the difference $(\hat{\tau}_k - p_c)$ and on the radii of the ponds.

In Theorem 1.6 we prove the existence of an IIC with several infinite p_c -open and p_c -closed paths from a neighborhood of the origin. Last, we show in Theorems 1.8 and 1.9 that the local description of the invaded region near the backbone of the IPC far away from the origin is given by the IIC with two infinite p_c -open paths, and the local description of the invaded region near an outlet of the IPC far away from the origin is given by the IIC with two infinite p_c -open paths and two infinite p_c -closed paths so that these paths alternate.

1.2 Notation

In this section we collect most of the notation and the definitions used in the paper.

For $a \in \mathbb{R}$, we write |a| for the absolute value of a, and, for a site $x = (x_1, x_2) \in \mathbb{Z}^2$, we write |x| for $\max(|x_1|, |x_2|)$. For n > 0 and $x \in \mathbb{Z}^2$, let $B(x, n) = \{y \in \mathbb{Z}^2 : |y - x| \le n\}$ and $\partial B(x, n) = \{y \in \mathbb{Z}^2 : |y - x| = n\}$. We write B(n) for B(0, n) and $\partial B(n)$ for $\partial B(0, n)$. For m < n and $x \in \mathbb{Z}^2$, we define the annulus $Ann(x; m, n) = B(x, n) \setminus B(x, m)$. We write Ann(m, n) for Ann(0; m, n).

We consider the square lattice $(\mathbb{Z}^2, \mathbb{E}^2)$, where $\mathbb{E}^2 = \{\langle x, y \rangle \in \mathbb{Z}^2 \times \mathbb{Z}^2 : |x-y| = 1\}$. Let $(\mathbb{Z}^2)^* = (1/2, 1/2) + \mathbb{Z}^2$ and $(\mathbb{E}^2)^* = (1/2, 1/2) + \mathbb{E}^2$ be the vertices and the edges of the dual lattice. For $x \in \mathbb{Z}^2$, we write x^* for x + (1/2, 1/2). For an edge $e \in \mathbb{E}^2$ we denote its endpoints (left respectively right or bottom respectively top) by e_x , $e_y \in \mathbb{Z}^2$. The edge $e^* = \langle e_x + (1/2, 1/2), e_y - (1/2, 1/2) \rangle$ is called the *dual edge* to e. Its endpoints (bottom respectively top or left respectively right) are denoted by e_x^* and e_y^* . Note that, in general, e_x^* and e_y^* are not the same as $(e_x)^*$ and $(e_y)^*$. For a subset $\mathcal{K} \subset \mathbb{Z}^2$, let $\mathcal{K}^* = (1/2, 1/2) + \mathcal{K}$. We say that an edge $e \in \mathbb{E}^2$ is in $\mathcal{K} \subset \mathbb{Z}^2$ if both its endpoints are in \mathcal{K} . For any graph \mathcal{G} we write $|\mathcal{G}|$ for the number of vertices in \mathcal{G} .

Let $(\tau_e)_{e \in \mathbb{E}^2}$ be independent random variables, uniformly distributed on [0,1], indexed by edges. We call τ_e the *weight* of an edge e. We define the weight of an edge e^* as $\tau_{e^*} = \tau_e$. We denote the underlying probability measure by \mathbb{P} and the space of configurations by $([0,1]^{\mathbb{E}^2}, \mathcal{F})$, where \mathcal{F} is the natural σ -field on $[0,1]^{\mathbb{E}^2}$. We say that an edge e is p-open if $\tau_e < p$ and p-closed if $\tau_e \ge p$. An edge e^* is p-open if e is p-open, and it is p-closed if e is e-closed. Accordingly, for e is e-closed.



the edge configuration $\omega_p \in \{0,1\}^{\mathbb{Z}^2}$ by $\omega_p(e) = 1$ if $\tau_e < p$ and 0 if $\tau_e \ge p$. We make a similar definition for ω_p^* , the dual edge configuration. The event that two sets of sites $\mathcal{K}_1, \mathcal{K}_2 \subset \mathbb{Z}^2$ are connected by a p-open path is denoted by $\mathcal{K}_1 \stackrel{p}{\longleftrightarrow} \mathcal{K}_2$, and the event that two sets of sites $\mathcal{K}_1^*, \mathcal{K}_2^* \subset (\mathbb{Z}^2)^*$ are connected by a p-closed path in the dual lattice is denoted by $\mathcal{K}_1^* \stackrel{p^*}{\longleftrightarrow} \mathcal{K}_2^*$. For any $n \ge 1$ and $p \in [0, 1]$, we define the event

 $B_{n,p} = \{\text{There is a } p\text{-closed circuit with radius at least } n \text{ around the origin in the dual lattice} \}.$

For $p \in [0, 1]$, we consider a probability space $(\Omega_p, \mathcal{F}_p, \mathbb{P}_p)$, where $\Omega_p = \{0, 1\}^{\mathbb{E}^2}$, \mathcal{F}_p is the σ -field generated by the finite-dimensional cylinders of Ω_p , and \mathbb{P}_p is a product measure on $(\Omega_p, \mathcal{F}_p)$, defined as $\mathbb{P}_p = \prod_{e \in \mathbb{E}^2} \mu_e$, where μ_e is the probability measure on $\{0, 1\}$ with $\mathbb{P}_p(\omega_e = 1) = \mu_e(\{1\}) = 1 - \mu_e(\{0\}) = 1 - \mathbb{P}_p(\omega_e = 0) = p$. We say that an edge e is open or occupied if $\omega_e = 1$, and e is closed or vacant if $\omega_e = 0$. We say that an edge e^* is open or occupied if e is open, and it is closed or vacant if e is closed. The event that two sets of sites $\mathcal{K}_1, \mathcal{K}_2 \subset \mathbb{Z}^2$ are connected by an open path is denoted by $\mathcal{K}_1 \leftrightarrow \mathcal{K}_2$, and the event that two sets of sites $\mathcal{K}_1^*, \mathcal{K}_2^* \subset \mathbb{Z}^2$ are connected by a closed path in the dual lattice is denoted by $\mathcal{K}_1^* \leftrightarrow \mathcal{K}_2^*$. For any $n \geq 1$ and $p \in [0, 1]$, let $\pi_n = \mathbb{P}_{p_c}(0 \leftrightarrow \partial B(n))$ and $\pi(n, p) = \mathbb{P}_p(0 \leftrightarrow \partial B(n))$. Also define the event

 $B_n = \{\text{There is a closed circuit with radius at least } n \text{ around the origin in the dual lattice}\}.$

For any $k \ge 1$, let \hat{R}_k be the radius of the union of the first k ponds. In other words, $\hat{R}_k = \max\{|x| : x \in \bigcup_{j=1}^k \hat{V}_j\}$. For two functions g and h from a set \mathcal{X} to \mathbb{R} , we write $g(z) \asymp h(z)$ to indicate that g(z)/h(z) is bounded away from 0 and ∞ , uniformly in $z \in \mathcal{X}$. Throughout this paper we write log for \log_2 . We also write \mathbb{P}_{cr} for \mathbb{P}_{p_c} . All the constants (C_i) in the proofs are strictly positive and finite. Their exact values may be different from proof to proof.

1.3 Main results

1.3.1 Weight of the kth outlet

Let $\hat{\tau}_k$ be the weight of the kth outlet, as defined in Sect. 1.1.

Theorem 1.1 For any $k \geq 1$,

$$\mathbb{P}(\hat{\tau}_k < p) \simeq (\log L(p))^{k-1} \theta(p), \quad p > p_c, \tag{1.1}$$

where the correlation length L(p) is defined in Sect. 2.



Remark 1 Note that the statement is trivial in the case k = 1. Indeed, it follows from the definition of the invasion that $\mathbb{P}(\hat{\tau}_1 < p) = \theta(p)$ for all p.

1.3.2 Volumes of the ponds

Theorem 1.2 For any $k \ge 1$,

$$\mathbb{P}(|\hat{V}_k| \ge n^2 \pi_n) \times (\log n)^{k-1} \pi_n, \quad n \ge 2. \tag{1.2}$$

In particular,

$$\mathbb{P}(|\hat{V}_k| \ge n) \times (\log n)^{k-1} \mathbb{P}_{cr}(|C(0)| \ge n), \quad n \ge 2.$$
 (1.3)

Remark 2 The case k = 1 is considered in [20].

Remark 3 The second set of inequalities follows from the first one using the relations $\mathbb{P}_{cr}(|C(0)| \ge n^2 \pi_n) \times \pi_n$ (see [20, Theorem 2]) and $\log(n^2 \pi_n) \times \log n$.

Remark 4 Let i_k be the index such that $e_{i_k} = \hat{e}_k$. Then i_k is comparable to $|\hat{V}_1| + \cdots + |\hat{V}_k|$. Therefore the statements (1.2) and (1.3) hold with $|\hat{V}_k|$ replaced by i_k .

1.3.3 Almost sure bounds

For any m < n, let O(m, n) be the number of outlets in Ann(m, n), and let O(n) be the number of outlets in B(n). We first give n-independent bounds on all moments of O(n, 2n).

Theorem 1.3 There exists $c_1 > 0$ such that for all $t, n \ge 1$,

$$\mathbb{E}(O(n,2n)^t) \le (c_1t)^{3t}. \tag{1.4}$$

In particular, there exists c_2 , $\lambda > 0$ such that for all n,

$$\mathbb{E}(\exp(\lambda O(n, 2n)^{1/3})) < c_2. \tag{1.5}$$

Next we show almost sure bounds on the sequence of random variables $(O(2^n))_{n>1}$.

Theorem 1.4 There exists c_3 , $c_4 > 0$ such that with probability one, for all large n,

$$c_3 n \le O(2^n) \le c_4 n. \tag{1.6}$$

Theorem 1.4 implies related bounds on the convergence rate of the weights $\hat{\tau}_k$ to p_c and on the growth of the radii $(\hat{R}_k)_{k\geq 1}$.

Corollary 1.1 1. There exists c_5 and c_6 with $1 < c_5$, $c_6 < \infty$ such that with probability one, for all large k,

$$(c_5)^k \le \hat{R}_k \le (c_6)^k. \tag{1.7}$$



2. There exists c_7 and c_8 with $0 < c_7, c_8 < 1$ such that with probability one, for all large k,

$$(c_7)^k \le \hat{\tau}_k - p_c \le (c_8)^k. \tag{1.8}$$

Remark 5 Asymptotics of various ponds statistics as well as CLT-type and large deviations results for deviations of those quantities away from their limits are studied in [8] for invasion percolation on regular trees. Not only do the results in [8] imply exponential almost sure bounds similar to (1.7) and (1.8), they are very explicit. For instance, it is shown that $\lim_{k\to\infty}\frac{1}{k}\ln(\hat{\tau}_k-p_c)=-1$ and $\lim_{k\to\infty}\frac{1}{k}\ln(\hat{R}_k)=1$ a.s.

Our last theorem concerns ratios of successive terms of the sequence $(\hat{\tau}_k - p_c)_{k \ge 1}$.

Theorem 1.5 With probability one, the set

$$\left\{ \frac{\hat{\tau}_{k+1} - p_c}{\hat{\tau}_k - p_c} : k \ge 1 \right\}$$

is a dense subset of [0, 1].

1.3.4 Outlets and multiple-armed IICs

First we recall the definition of the incipient infinite cluster from [12]. It is shown in [12] that the limit

$$\nu(E) = \lim_{N \to \infty} \mathbb{P}_{cr}(E \mid 0 \leftrightarrow \partial B(N))$$

exists for any event E that depends on the state of finitely many edges in \mathbb{E}^2 . The unique extension of ν to a probability measure on configurations of open and closed edges exists. Under this measure, the open cluster of the origin is a.s. infinite. It is called the *incipient infinite cluster* (IIC). In Theorem 1.7 [10, Theorem 3], a relation between IPC and IIC is given.

In this section we introduce multiple-armed IIC measures (Theorem 1.6) and study their relation to invasion percolation (Theorems 1.8, 1.9). For this, let $k \geq 1$ and $\sigma \in \{\text{open, closed}\}^k$. Let r_1 be the number of 'open' entries in σ and let r_2 be the number of 'closed' entries in σ . For l < n such that $|\partial B(l)| > |\sigma|$, we say that B(l) is σ -connected to $\partial B(n)$, denoted $B(l) \leftrightarrow_{\sigma} \partial B(n)$, if there exist r_1 disjoint open paths between B(l) and $\partial B(n)$ and r_2 disjoint dual closed paths between $B(l)^*$ and $\partial B(n)^*$ such that the relative counterclockwise arrangement of these paths is given by σ . In the definition above we allow $n = \infty$, in this case we write $B(l) \leftrightarrow_{\sigma} \infty$.

Theorem 1.6 Suppose that σ is alternating and let l be the minimal number such that $|\partial B(l)| \ge |\sigma|$. For every cylinder event E, the limit

$$\nu_{\sigma}(E) = \lim_{n \to \infty} \mathbb{P}_{cr}(E \mid B(l) \leftrightarrow_{\sigma} \partial B(n))$$
 (1.9)



exists. For the unique extension of v_{σ} to a probability measure on the configurations of open and closed edges,

$$\nu_{\sigma}(B(l) \leftrightarrow_{\sigma} \infty) = 1.$$

We call the resulting measure v_{σ} the σ -incipient infinite cluster measure.

Remark 6 Note that Kesten's IIC measure corresponds to the case $\sigma = \{\text{open}\}\)$ with l = 0. One can check that Kesten's original proof [12] also works for the case $\sigma = \{\text{open}, \text{open}\}\)$. We use this second IIC measure in Theorem 1.8.

Remark 7 The proof we present of Theorem 1.6 can be easily modified to give the existence of IICs for σ 's which either do not contain neighboring open paths or do not contain neighboring closed paths (here we take the first and last elements of σ to be neighbors). In particular, it works for any 3-arm IIC and for monochromatic IICs. In the case when there are neighboring open paths (but no neighboring closed paths) one needs to change the proof by considering closed circuits with defects instead of open circuits.

Define $\mathcal{O}=\{\hat{e}_k: k\geq 1\}$, the set of *outlets* of the invasion and let \mathcal{B} be the *backbone*, i.e., those vertices which are connected in the IPC of the origin by two disjoint paths, one to the origin and one to ∞ . For any vertex v, define the shift operator θ_v on configurations ω so that for any edge e, $\theta_v(\omega)(e)=\omega(e-v)$, where $e-v=\langle e_x-v,e_y-v\rangle$. For any event E, define

$$\theta_v E = \{\theta_v(\omega) : \omega \in E\},\$$

and if K is a set of edges in \mathbb{E}^2 , define

$$E_{\mathcal{K}} = \{\mathcal{K} \subset \mathcal{S}\}, \quad \theta_v \mathcal{K} = \{e \in \mathbb{E}^2 : e - v \in \mathcal{K}\}, \quad \text{and} \quad \Theta_v E_{\mathcal{K}} = \{\theta_v \mathcal{K} \subset \mathcal{S}\}.$$

For an edge e, let ρ_e be the rotation of the lattice around the origin that maps $e-e_x$ to $\langle (0,0),(1,0)\rangle$. We define the operator θ_e on configurations ω so that for any edge $f,\theta_e(\omega)(f)=\omega(\rho_e(f-e_x))$. We define $\theta_e E,\theta_e \mathcal{K}$, and $\Theta_e E_{\mathcal{K}}$ similarly. Let $E_{\mathcal{K}}'$ be the event that \mathcal{K} is contained in the cluster of the origin. We recall [10, Theorem 3], which states that asymptotically the distribution of invaded edges near a vertex v is given by the IIC measure.

Theorem 1.7 Let E be an event which depends on finitely many values $\omega_{p_c}(\cdot)$ and let $\mathcal{K} \subset \mathbb{E}^2$ be finite.

$$\lim_{|v|\to\infty} \mathbb{P}(\theta_v E \mid v \in \mathcal{S}) = \nu(E) \quad and \quad \lim_{|v|\to\infty} \mathbb{P}(\Theta_v E_{\mathcal{K}} \mid v \in \mathcal{S}) = \nu(E_{\mathcal{K}}'),$$

where the measure on the right is the IIC measure.

We are interested in the distribution of invaded edges near the backbone (Theorem 1.8) or near an outlet (Theorem 1.9). While the analysis of the distribution of the



invaded edges near the backbone is very similar to the proof of Theorem 1.7, the study of the distribution of the invaded edges near an outlet is more involved. Define $\tilde{v}^{2,0}$ to be the measure constructed in the same way as $v^{2,0}$ except that we condition that the origin is connected to $\partial B(n)$ by two disjoint open paths and take n to ∞ . Define $\tilde{v}^{2,2}$ similarly, but by conditioning that the endpoints of the edge $e_0 = \langle (0,0), (1,0) \rangle$ are connected to $\partial B(n)$ by two disjoint open paths and that the endpoints of e_0^* are connected to $\partial B(n)^*$ by two disjoint closed dual paths and taking n to ∞ . Obvious modifications of Theorem 1.6 hold for these measures.

Theorem 1.8 Let E be an event which depends on finitely many values $\omega_{p_c}(\cdot)$ and let $\mathcal{K} \subset \mathbb{E}^2$ be finite.

$$\lim_{|v|\to\infty} \mathbb{P}(\theta_v E \mid v \in \mathcal{B}) = \tilde{v}^{2,0}(E) \quad and \quad \lim_{|v|\to\infty} \mathbb{P}(\Theta_v E_{\mathcal{K}} \mid v \in \mathcal{B}) = \tilde{v}^{2,0}(E_{\mathcal{K}}').$$

Proof Similar to the proof of [10, Theorem 3].

Theorem 1.9 Let E be an event which depends on finitely many values $\omega_{p_c}(\cdot)$ (but not on $\omega_{p_c}(e_0)$), and let K be a finite set of edges such that $e_0 \notin K$.

$$\lim_{|e|\to\infty} \mathbb{P}(\theta_e E \mid e \in \mathcal{O}) = \tilde{v}^{2,2}(E) \quad and \quad \lim_{|e|\to\infty} \mathbb{P}(\Theta_e E_{\mathcal{K}} \mid e \in \mathcal{O}) = \tilde{v}^{2,2}(E_{\mathcal{K}}').$$

1.4 Structure of the paper

We define the correlation length and state some of its properties in Sect. 2. We prove Theorems 1.1 and 1.2 in Sects. 3 and 4, respectively. The proofs of Theorems 1.3–1.5 are in Sect. 5: the proof of Theorem 1.3 is in Sect. 5.1; the proofs of Theorem 1.4 and Corollary 1.1 are in Sect. 5.2; and the proof of Theorem 1.5 is in Sect. 5.3. We prove Theorem 1.6 in Sect. 6 and Theorem 1.9 in Sect. 7. For the notation in Sects. 3–7 we refer the reader to Sect. 1.2.

2 Correlation length and preliminary results

In this section we define the correlation length that will play a crucial role in our proofs. The correlation length was introduced in [2] and further studied in [13].

2.1 Correlation length

For m, n positive integers and $p \in (p_c, 1]$ let

 $\sigma(n, m, p) = \mathbb{P}_p$ (there is an open horizontal crossing of $[0, n] \times [0, m]$).

Given $\varepsilon > 0$ and $p > p_c$, we define

$$L(p,\varepsilon) = \min\{n : \sigma(n,n,p) > 1 - \varepsilon\}. \tag{2.1}$$



 $L(p, \varepsilon)$ is called the finite-size scaling correlation length and it is known that $L(p, \varepsilon)$ scales like the usual correlation length (see [13]). It was also shown in [13] that the scaling of $L(p, \varepsilon)$ is independent of ε given that it is small enough, i.e. there exists $\varepsilon_0 > 0$ such that for all $0 < \varepsilon_1, \varepsilon_2 \le \varepsilon_0$ we have $L(p, \varepsilon_1) \times L(p, \varepsilon_2)$. For simplicity we will write $L(p) = L(p, \varepsilon_0)$ for the entire paper. We also define

$$p_n = \sup\{p : L(p) > n\}.$$

It is easy to see that $L(p) \to \infty$ as $p \to p_c$ and $L(p) \to 0$ as $p \to 1$. In particular, the probability p_n is well-defined. It is clear from the definitions of L(p) and p_n and from the RSW theorem that, for positive integers k and k, there exists $\delta_{k,l} > 0$ such that, for any positive integer k and for all k and k are considered as k.

 \mathbb{P}_p (there is an open horizontal crossing of $[0, kn] \times [0, ln]$) $> \delta_{k,l}$

and

 \mathbb{P}_p (there is a closed horizontal dual crossing of $([0, kn] \times [0, ln])^*$) $> \delta_{k,l}$.

By the FKG inequality and a standard gluing argument [9, Section 11.7] we get that, for positive integers n and $k \ge 2$ and for all $p \in [p_c, p_n]$,

 $\mathbb{P}_p(Ann(n,kn))$ contains an open circuit around the origin) $> (\delta_{k,k-2})^4$

and

 $\mathbb{P}_p(Ann(n,kn)^*$ contains a closed dual circuit around the origin) $> (\delta_{2k,k-1})^4$.

2.2 Preliminary results

For any positive l we define $\log^{(0)} l = l$ and $\log^{(j)} l = \log(\log^{(j-1)} l)$ for all $j \ge 1$, as long as the right-hand side is well defined. For l > 10, let

$$\log^* l = \min\{j > 0 : \log^{(j)} l \text{ is well-defined and } \log^{(j)} l \le 10\}.$$
 (2.2)

Our choice of the constant 10 is quite arbitrary, we could take any other large enough positive number instead of 10. For l > 10, let

$$p_{l}(j) = \begin{cases} \inf \left\{ p > p_{c} : L(p) \leq \frac{l}{C_{*} \log^{(j)} l} \right\} & \text{if } j \in (0, \log^{*} l), \\ p_{c} & \text{if } j \geq \log^{*} l, \\ 1 & \text{if } j = 0. \end{cases}$$
 (2.3)

The value of C_* will be chosen differently in each proof. For any C_* , notice that there exists a universal constant $L_0(C_*) > 10$ such that $p_l(j)$ are well-defined if $l > L_0(C_*)$ and non-increasing in l. The last observation follows from monotonicity of L(p) and



the fact that the functions $l/\log^{(j)}l$ are non-decreasing in l for $j\in(0,\log^*l)$ and $l\geq 3$.

We give the following results without proofs.

1. ([10, (2.10)]) There exists a universal constant D_1 such that, for every $l > L_0(C_*)$ and $j \in (0, \log^* l)$,

$$C_* \log^{(j)} l \le \frac{l}{L(p_l(j))} \le D_1 C_* \log^{(j)} l.$$
 (2.4)

2. ([13, Theorem 2]) There is a constant D_2 such that, for all $p > p_c$,

$$\theta(p) \le \mathbb{P}_p \left[0 \leftrightarrow \partial B(L(p)) \right] \le D_2 \mathbb{P}_{cr} \left[0 \leftrightarrow \partial B(L(p)) \right],$$
 (2.5)

where $\theta(p) = \mathbb{P}_p(0 \to \infty)$ is the percolation function for Bernoulli percolation.

3. ([17, Section 4]) There is a constant D_3 such that, for all $n \ge 1$,

$$\mathbb{P}_{p_n}(B(n) \leftrightarrow \infty) \ge D_3. \tag{2.6}$$

4. ([13, (3.61)]) There is a constant D_4 such that, for all positive integers $r \leq s$,

$$\frac{\mathbb{P}_{cr}(0 \leftrightarrow \partial B(s))}{\mathbb{P}_{cr}(0 \leftrightarrow \partial B(r))} \ge D_4 \sqrt{\frac{r}{s}}.$$
(2.7)

5. There exist positive constants D_5 and D_6 such that, for all $p > p_c$,

$$\mathbb{P}_p(B_n) \le D_5 \exp\left\{-D_6 \frac{n}{L(p)}\right\}. \tag{2.8}$$

It follows, for example, from [10, (2.6) and (2.8)] (see also [18, Lemma 37 and Remark 38]).

6. ([18, Proposition 34]) Fix $e = \langle (0,0), (1,0) \rangle$, and let $A_n^{2,2}$ be the event that e_x and e_y are connected to $\partial B(n)$ by open paths, and e_x^* and e_y^* are connected to $\partial B(n)^*$ by closed dual paths. Note that these four paths are disjoint and alternate. Then

$$(p_n - p_c)n^2 \mathbb{P}_{cr}(A_n^{2,2}) \approx 1, \quad n \ge 1.$$
 (2.9)

3 Proof of Theorem 1.1

We give the proof for the case k = 2. The proof for $k \ge 3$ is similar to the proof for k = 2, and we omit the details. Note that [13, Theorem 2] it is sufficient to prove that

$$\mathbb{P}(\hat{\tau}_2 < p) \simeq (\log L(p)) \mathbb{P}_{cr} (0 \leftrightarrow \partial B(L(p))). \tag{3.1}$$



We first prove the upper bound. We partition the box B(L(p)) into $\lfloor \log L(p) \rfloor$ disjoint annuli:

$$\mathbb{P}(\hat{\tau}_2 < p) \leq \mathbb{P}(\hat{R}_1 \geq L(p)) + \sum_{k=0}^{\lfloor \log L(p) \rfloor} \mathbb{P}\left(\hat{\tau}_2 < p; \hat{R}_1 \in \left[\frac{L(p)}{2^{k+1}}, \frac{L(p)}{2^k}\right)\right).$$

We show that there is a universal constant C_1 such that for any $p > p_c$ and $m \le L(p)/2$,

$$\mathbb{P}\left(\hat{\tau}_2 < p; \hat{R}_1 \in [m, 2m]\right) \le C_1 \mathbb{P}_{cr}\left(0 \leftrightarrow \partial B(L(p))\right). \tag{3.2}$$

From [20], $\mathbb{P}(\hat{R}_1 \geq L(p)) \leq C_2 \mathbb{P}_{cr} \ (0 \leftrightarrow \partial B(L(p)))$. Therefore the upper bound in (3.1) will immediately follow from (3.2). We partition the event $\{\hat{\tau}_2 < p; \hat{R}_1 \in [m, 2m]\}$ according to the value of $\hat{\tau}_1$:

$$\sum_{j=1}^{\log^* m} \mathbb{P}\left(\hat{\tau}_2 < p; \hat{R}_1 \in [m, 2m]; \hat{\tau}_1 \in [p_m(j), p_m(j-1))\right). \tag{3.3}$$

Note that if the event $\{\hat{R}_1 \geq m, \hat{\tau}_1 \in [p_m(j), p_m(j-1))\}$ occurs then (a) there is a $p_m(j-1)$ -open path from the origin to $\partial B(m)$, and (b) the origin is surrounded by a $p_m(j)$ -closed circuit of diameter at least m in the dual lattice. Also note that if the event $\{\hat{\tau}_2 < p, \hat{R}_1 \leq 2m\}$ occurs then there is a p-open path from B(2m) to $\partial B(L(p))$.

From the above observations, it follows that the sum (3.3) is bounded from above by

$$\sum_{i=1}^{\log^* m} \mathbb{P}\left(0 \stackrel{p_m(j-1)}{\longleftrightarrow} \partial B(m); B(2m) \stackrel{p}{\longleftrightarrow} \partial B(L(p)); B_{m,p_m(j)}\right).$$

The FKG inequality and independence give an upper bound of

$$\sum_{j=1}^{\log^* m} \mathbb{P}_{p_m(j-1)}(0 \leftrightarrow \partial B(m)) \mathbb{P}_p(B(2m) \leftrightarrow \partial B(L(p))) \mathbb{P}_{p_m(j)}(B_m).$$

It follows from (2.4) and (2.8) that $\mathbb{P}_{p_m(j)}(B_m) \leq C_3(\log^{(j-1)} m)^{-C_4}$ for some C_3 and C_4 , where C_4 can be made arbitrarily large given that C_* is made large enough. Inequalities (2.7) and (2.5) give

$$\mathbb{P}_{D_m(j-1)}(0 \leftrightarrow \partial B(m)) \le C_5(\log^{(j-1)} m)^{\frac{1}{2}} \mathbb{P}_{Cr}(0 \leftrightarrow \partial B(m)),$$



and (2.5) and the RSW Theorem give

$$\mathbb{P}_p(B(2m) \leftrightarrow \partial B(L(p))) \le C_6 \mathbb{P}_{cr}(B(2m) \leftrightarrow \partial B(L(p))).$$

Also, the RSW Theorem and the FKG inequality imply that

$$\mathbb{P}_{cr}(0 \leftrightarrow \partial B(m))\mathbb{P}_{cr}(B(2m) \leftrightarrow \partial B(L(p))) \leq C_7 \mathbb{P}_{cr}(0 \leftrightarrow \partial B(L(p))).$$

Therefore, we obtain that the probability $\mathbb{P}\left(\hat{\tau}_2 < p; \hat{R}_1 \in [m, 2m]\right)$ is bounded from above by

$$C_8 \mathbb{P}_{cr}(0 \leftrightarrow \partial B(L(p))) \sum_{j=1}^{\log^* m} (\log^{(j-1)} m)^{-C_4+1/2}.$$

As in [10, (2.26)], one can easily show that, for $C_4 > 1$,

$$\sum_{j=1}^{\log^* m} (\log^{(j-1)} m)^{-C_4+1/2} < C_9.$$

The upper bound in (3.1) follows.

We now prove the lower bound in (1.1). For $p > p_c$ and a positive integer m < L(p)/2, we consider the event $C_{m,p}$ that there exists an edge $e \in Ann(m, 2m)$ such that

- $\tau_e \in (p_c, p_m);$
- there exist two p_c -open paths in $B(2L(p))\setminus \{e\}$, one connecting the origin to one of the endpoints of e, and another connecting the other endpoint of e to the boundary of B(2L(p));
- there exists a p_m -closed dual path P in $Ann(m, 2m)^* \setminus \{e^*\}$ connecting the endpoints of e^* such that $P \cup \{e^*\}$ is a circuit around the origin;
- there exists a p_c -open circuit around the origin in Ann(L(p), 2L(p));
- there exists a p-open path connecting B(L(p)) to infinity.

See Fig. 1 for an illustration of the event $C_{m,p}$. It can be shown similarly to [4, Corollary 6.2] that

$$\mathbb{P}(C_{m,p}) \ge C_{10} \mathbb{P}_{cr}(0 \leftrightarrow \partial B(L(p))),$$

where we also use the fact that $\mathbb{P}_p(B(L(p)) \leftrightarrow \infty) > C_{11}$ [see (2.6)]. It remains to notice that for fixed p, the events $C_{\lfloor L(p)/2^k \rfloor, p}$ are disjoint and each of them implies the event $\{\hat{\tau}_2 < p\}$. Therefore,

$$\mathbb{P}(\hat{\tau}_2 < p) \ge \sum_{k=0}^{\lfloor \log L(p) \rfloor - 1} \mathbb{P}(C_{\lfloor L(p)/2^k \rfloor, p}) \ge C_{10} \lfloor \log L(p) \rfloor \mathbb{P}_{cr}(0 \leftrightarrow \partial B(L(p))).$$



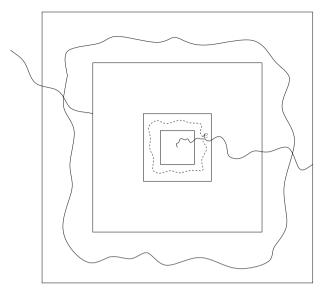


Fig. 1 The event $C_{m,p}$. The *boxes*, in order from smallest to largest, are B(m), B(2m), B(L(p)), and B(2L(p)). The *dotted path* is p_m -closed, the path to infinity is p_m -open, and all other paths are p_c -open

4 Proof of Theorem 1.2

The case k = 1 is considered in [20, Theorem 2]. We give the proof for k = 2. The proof for $k \ge 3$ is similar to the proof for k = 2, and we omit the details.

We first prove the upper bound. By the RSW Theorem, it is sufficient to bound the probability $\mathbb{P}(|\hat{V}_2| \geq 2n^2\pi_n)$. We partition this probability according to the value of the radii \hat{R}_1 and \hat{R}_2 , defined in Sect. 1.2. Without loss of generality we can assume that $n = 2^N$.

$$\mathbb{P}(|\hat{V}_2| \ge 2n^2 \pi_n) \le \mathbb{P}(\hat{R}_2 \ge n) + \sum_{m=1}^{N} \sum_{k=1}^{m} \mathbb{P}\left(|\hat{V}_2| \ge 2n^2 \pi_n; \, \hat{R}_1 \in [2^{k-1}, 2^k); \, \hat{R}_2 \in [2^{m-1}, 2^m)\right).$$

It follows from [4] that $\mathbb{P}(\hat{R}_2 \ge n) \times (\log n)\pi_n$. We now consider the second term. We decompose the probability of the event

$$E_{n,k,m} = \left\{ |\hat{V}_2| \ge 2n^2 \pi_n; \, \hat{R}_1 \in [2^{k-1}, 2^k); \, \hat{R}_2 \in [2^{m-1}, 2^m) \right\}$$

according to the values of $\hat{\tau}_1$ and $\hat{\tau}_2$:

$$\sum_{i=1}^{\log^* 2^k} \sum_{j=1}^{2^m} \mathbb{P}\left(E_{n,m,k}; \hat{\tau}_1 \in [p_{2^k}(i), p_{2^k}(i-1)); \hat{\tau}_2 \in [p_{2^m}(j), p_{2^m}(j-1))\right). \tag{4.1}$$



We consider the event $D_{n,k,m}$ that the number of vetices in the annulus $Ann(2^k, 2^m)$ connected to $B(2^k)$ inside $Ann(2^k, 2^m)$ is at least $n^2\pi_n$. If the vertices in the definition of $D_{n,k,m}$ are connected to $B(2^k)$ by p-open paths, we denote the corresponding event by $D_{n,k,m}(p)$. We also consider the event $D_{n,k}$ that the number of vertices in the box $B(2^k)$ connected to the boundary $\partial B(2^k)$ is at least $n^2\pi_n$. If the vertices in the definition of $D_{n,k}$ are connected to $\partial B(2^k)$ by p-open paths, we denote the corresponding event by $D_{n,k}(p)$. The probability of a typical summand in (4.1) can be bounded from above by

$$\mathbb{P}\left(\begin{array}{c} B_{2^{k-1},p_{2^{k}}(i)}; \, B_{2^{m-1},p_{2^{m}}(j)}; \, 0 \overset{p_{2^{k}}(i-1)}{\longleftrightarrow} \, \partial B(2^{k}); \, B(2^{k}) \overset{p_{2^{m}}(j-1)}{\longleftrightarrow} \, \infty; \\ D_{n,k,m}(p_{2^{m}}(j-1)) \cup D_{n,k}(p_{2^{k}}(i-1)) \end{array}\right),$$

where we use the fact that $\hat{\tau}_1 > \hat{\tau}_2$ a.s.

We use the FKG inequality and independence to estimate the above probability. It is no greater than

$$\mathbb{P}\left(B_{2^{k-1},p_{2^{k}}(i)}; B_{2^{m-1},p_{2^{m}}(j)}\right) \mathbb{P}_{p_{2^{k}}(i-1)}\left(0 \leftrightarrow \partial B(2^{k})\right) \mathbb{P}_{p_{2^{m}}(j-1)} \\
\times \left(B(2^{k}) \leftrightarrow \infty; D_{n,k,m}\right) \qquad (4.2) \\
+ \mathbb{P}\left(B_{2^{k-1},p_{2^{k}}(i)}; B_{2^{m-1},p_{2^{m}}(j)}\right) \mathbb{P}_{p_{2^{k}}(i-1)}\left(0 \leftrightarrow \partial B(2^{k}); D_{n,k}\right) \mathbb{P}_{p_{2^{m}}(j-1)} \\
\times \left(B(2^{k}) \leftrightarrow \infty\right). \qquad (4.3)$$

The probability $\mathbb{P}(B_{2^{k-1},p_{>k}(i)}; B_{2^{m-1},p_{2^m}(j)})$ is bounded from above by [4, (6.6)]

$$C_1(\log^{(i-1)} 2^k)^{-C_2}(\log^{(j-1)} 2^m)^{-C_2}$$

where the constant C_2 can be made arbitrarily large given C_* is made large enough. We first estimate (4.2). It follows from (2.7) that

$$\begin{split} \mathbb{P}_{p_{2^k}(i-1)} \left(0 \leftrightarrow \partial B(2^k) \right) &\leq C_3 (\log^{(i-1)} 2^k)^{1/2} \mathbb{P}_{cr}(0 \leftrightarrow \partial B(2^k)) \\ &\leq C_4 (\log^{(i-1)} 2^k)^{1/2} (\log^{(j-1)} 2^m)^{1/2} \\ &\times \frac{\mathbb{P}_{cr}(0 \leftrightarrow \partial B(2^m))}{\mathbb{P}_{p_{2^m}(j-1)}(B(2^k) \leftrightarrow \infty)}. \end{split}$$

Substitution gives the following upper bound for (4.2):

$$\mathbb{P}_{cr}(0 \leftrightarrow \partial B(2^m)) C_1 C_4(\log^{(i-1)} 2^k \log^{(j-1)} 2^m)^{-C_2 + 1/2} \mathbb{P}_{p_{2^m}(j-1)} \times (D_{n,k,m} \mid B(2^k) \leftrightarrow \infty).$$

We now estimate (4.3). It follows from the FKG inequality and (2.7) that

$$\mathbb{P}_{p_{2^k}(i-1)}\left(0 \leftrightarrow \partial B(2^k)\right) \leq C_5(\log^{(i-1)} 2^k)^{1/2} \mathbb{P}_{cr}\left(0 \leftrightarrow \partial B(2^k)\right).$$



Substitution gives the following upper bound for (4.3):

$$\mathbb{P}_{cr}(0 \leftrightarrow \partial B(2^k)) C_1 C_5(\log^{(i-1)} 2^k \log^{(j-1)} 2^m)^{-C_2 + 1/2} \mathbb{P}_{p_{2^k}(i-1)} \times (D_{n,k} \mid 0 \leftrightarrow \partial B(2^k)).$$

Therefore, the sum (4.1) is bounded from above by

$$C_{6}\mathbb{P}_{cr}(0 \leftrightarrow \partial B(2^{m})) \sum_{i=1}^{\log^{*} 2^{k}} \sum_{j=1}^{\log^{*} 2^{m}} (\log^{(i-1)} 2^{k} \log^{(j-1)} 2^{m})^{-C_{2}+1/2} \mathbb{P}_{p_{2^{m}}(j-1)}$$

$$\times (D_{n,k,m} \mid B(2^{k}) \leftrightarrow \infty)$$

$$+ C_{6}\mathbb{P}_{cr}(0 \leftrightarrow \partial B(2^{k})) \sum_{i=1}^{\log^{*} 2^{k}} \sum_{j=1}^{\log^{*} 2^{m}} (\log^{(i-1)} 2^{k} \log^{(j-1)} 2^{m})^{-C_{2}+1/2} \mathbb{P}_{p_{2^{k}}(i-1)}$$

$$\times (D_{n,k} \mid 0 \leftrightarrow \partial B(2^{k})).$$

Note that [10, (2.26)] if $C_2 > 1/2$, then there exists $C_7 > 0$ such that for all k,

$$\sum_{i=1}^{\log^* 2^k} (\log^{(i-1)} 2^k)^{-C_2 + 1/2} \le C_7 < \infty.$$

Also note that analogously to [20, Lemma 4] one can show that there exist $C_8 - C_{11}$ such that, for all $p > p_c$,

$$\mathbb{P}_p(D_{n,k,m} \mid B(2^k) \leftrightarrow \infty) \le C_8 \exp\left\{-C_9 \frac{n^2 \pi_n}{2^{2m} \pi(2^m, p)}\right\}$$

and

$$\mathbb{P}_{p}(D_{n,k} \mid 0 \leftrightarrow \partial B(2^{k})) \leq C_{10} \exp\left\{-C_{11} \frac{n^{2} \pi_{n}}{2^{2k} \pi(2^{k}, p)}\right\},\,$$

where π_n and $\pi(n, p)$ are defined in Sect. 1.2. In particular,

$$\mathbb{P}_{p_{2^m}(j-1)}(D_{n,k,m} \mid B(2^k) \leftrightarrow \infty) \le C_8 \exp\left\{-C_9 \frac{n^2 \pi_n}{2^{2m} \pi (2^m, p_{2^m}(j-1))}\right\}$$

$$\le C_8 \exp\left\{-C_{12} \frac{n^2 \pi_n}{2^{2m} \pi_{2^m}} (\log^{(j-1)} 2^m)^{-1/2}\right\},$$



and, similarly,

$$\mathbb{P}_{p_{2^k}(i-1)}(D_{n,k} \mid 0 \leftrightarrow \partial B(2^k)) \le C_{10} \exp\left\{-C_{11} \frac{n^2 \pi_n}{2^{2k} \pi (2^k, p_{2^k}(i-1))}\right\}$$

$$\le C_{10} \exp\left\{-C_{13} \frac{n^2 \pi_n}{2^{2k} \pi_{2^k}} (\log^{(i-1)} 2^k)^{-1/2}\right\}.$$

Therefore, the sum $\sum_{m=1}^{N} \sum_{k=1}^{m} \mathbb{P}(E_{n,k,m})$ is not bigger than

$$C_{14}(\log n)\pi_{n} \sum_{m=1}^{N} \frac{\pi_{2^{m}}}{\pi_{n}} \sum_{j=1}^{\log^{*} 2^{m}} (\log^{(j-1)} 2^{m})^{-C_{2}+1/2}$$

$$\times \exp\left\{-C_{12} \frac{n^{2} \pi_{n}}{2^{2m} \pi_{2^{m}}} (\log^{(j-1)} 2^{m})^{-1/2}\right\}$$

$$+C_{14}(\log n)\pi_{n} \sum_{k=1}^{N} \frac{\pi_{2^{k}}}{\pi_{n}} \sum_{i=1}^{\log^{*} 2^{k}} (\log^{(i-1)} 2^{k})^{-C_{2}+1/2}$$

$$\times \exp\left\{-C_{13} \frac{n^{2} \pi_{n}}{2^{2k} \pi_{2^{k}}} (\log^{(i-1)} 2^{k})^{-1/2}\right\}, \tag{4.4}$$

where $\log n$ comes from the fact that $\sum_{k=1}^{m} 1 = m \le N = \log n$. Finally, it follows from [20, p. 419] that

$$\sum_{m=1}^{N} \frac{\pi_{2^m}}{\pi_n} \sum_{j=1}^{\log^* 2^m} (\log^{(j-1)} 2^m)^{-C_2 + 1/2} \exp\left\{ -C_{12} \frac{n^2 \pi_n}{2^{2m} \pi_{2^m}} (\log^{(j-1)} 2^m)^{-1/2} \right\}$$

$$\leq C_{15} < \infty.$$

A similar bound holds for the summand (4.4). The proof for the second inequality in (1.2) is completed.

We now prove the first inequality in (1.2). For $m \le N$, let $C_{n,m}$ be the event that there exists an edge in $Ann(2^{m-1}, 2^m)$ such that

- its weight $\tau_e \in (p_c, p_{2^m})$;
- there exist two disjoint p_c -open paths, one connecting an end of e to the origin, and one connecting the other end of e to $\partial B(2n)$;
- there exist a p_{2^m} -closed dual path connecting the endpoints of e^* in $Ann(2^{m-1}, 2^m)^*$;
- there exists a p_c -open circuit in Ann(n, 2n).

It can be shown similarly to [4, Corollary 6.2] that $\mathbb{P}(C_{n,m}) \simeq \pi_n$. We also note that the events $C_{n,m}$ are disjoint and each of them implies the event $\{\hat{R}_2 \geq n\}$. Using the arguments from the proof of [4, Corollary 6.2], it follows that, for any $x \in Ann(2^{N-1}, n) =: A_n$ and $1 \leq m \leq N-2$,

$$\mathbb{P}(x \stackrel{p_c}{\longleftrightarrow} \partial B(2n) \mid C_{n,m}) \ge C_{16}\pi_n,$$



from which we conclude that

$$\mathbb{E}(|(\hat{V}_1 \cup \hat{V}_2) \cap A_n| | C_{n,m}) \ge C_{17} n^2 \pi_n. \tag{4.5}$$

We will show later that, for $1 \le m \le N - 2$,

$$\mathbb{E}(|(\hat{V}_1 \cup \hat{V}_2) \cap A_n|^2 \mid C_{n,m}) \le C_{18} \left(\mathbb{E}(|(\hat{V}_1 \cup \hat{V}_2) \cap A_n| \mid C_{n,m}) \right)^2. \tag{4.6}$$

If (4.6) holds, the second moment estimate gives that, for some $C_{19} > 0$,

$$\mathbb{P}(|\hat{V}_1 \cup \hat{V}_2| \ge C_{19}n^2\pi_n; C_{n,m}) \ge C_{19}\mathbb{P}(C_{n,m}) \ge C_{20}\pi_n.$$

Therefore

$$\mathbb{P}(|\hat{V}_1 \cup \hat{V}_2| \ge C_{19}n^2\pi_n; \, \hat{R}_2 \ge n) \ge \sum_{m=1}^{N-2} \mathbb{P}\left(|\hat{V}_1 \cup \hat{V}_2| \ge C_{19}n^2\pi_n; \, C_{n,m}\right)$$

$$\ge C_{20}(N-2)\pi_n.$$

In particular, using (2.7), we obtain $\mathbb{P}(|\hat{V}_1 \cup \hat{V}_2| \ge n^2 \pi_n) \ge C_{21}(\log n)\pi_n$. Recall that $\mathbb{P}(|\hat{V}_1| \ge n^2 \pi_n) \times \pi_n$. It immediately gives the inequality $\mathbb{P}(|\hat{V}_2| \ge n^2 \pi_n) \ge C_{22}(\log n)\pi_n$.

It remains to prove (4.6). Note that

$$\mathbb{E}(|(\hat{V}_1 \cup \hat{V}_2) \cap A_n|^2 | C_{n,m}) = \sum_{x,y \in A_n} \mathbb{P}(x, y \in \hat{V}_1 | C_{n,m}) + \sum_{x,y \in A_n} \mathbb{P}(x, y \in \hat{V}_2 | C_{n,m}),$$
(4.7)

where we use the fact that, by construction, \hat{V}_1 and \hat{V}_2 cannot both intersect A_n . We estimate the two sums on the r.h.s. separately. We only consider the first sum. The other sum is treated similarly. We decompose the probability $\mathbb{P}(x, y \in \hat{V}_1; C_{n,m})$ according to the value of $\hat{\tau}_1$:

$$\sum_{j=1}^{\log^* n} \mathbb{P}(x, y \in \hat{V}_1; C_{n,m}; \hat{\tau}_1 \in [p_n(j), p_n(j-1))).$$

Using arguments as in the first part of the proof of this theorem, the above sum is bounded from above by



$$\begin{split} & \sum_{j=1}^{\log^* n} \mathbb{P} \left(\begin{array}{c} 0 & \stackrel{p_c}{\longleftrightarrow} \partial B(2^{m-1}); \ B(2^m) & \stackrel{p_c}{\longleftrightarrow} \partial B(2^{N-2}); \\ x & \stackrel{p_n(j-1)}{\longleftrightarrow} \partial B(x, 2^{N-2}); \ y & \stackrel{p_n(j-1)}{\longleftrightarrow} \partial B(y, 2^{N-2}); \\ B_{n, p_n(j)} & \end{array} \right) \\ & \leq \mathbb{P}_{cr}(0 & \leftrightarrow \partial B(2^{m-1})) \mathbb{P}_{cr}(B(2^m) & \leftrightarrow \partial B(2^{N-2})) \\ & \times \sum_{j=1}^{\log^* n} \mathbb{P}_{p_n(j)}(B_n) \mathbb{P}_{p_n(j-1)}(x & \leftrightarrow \partial B(x, 2^{N-2}); \ y & \leftrightarrow \partial B(y, 2^{N-2})). \end{split}$$

Again, using tools from the first part of the proof of this theorem (see also the proof of Theorem 1.5 in [4]), the above sum is no greater than

$$C_{23}\mathbb{P}_{cr}(0 \leftrightarrow \partial B(n))\mathbb{P}_{cr}(x \leftrightarrow \partial B(x, 2^{N-2}); y \leftrightarrow \partial B(y, 2^{N-2})).$$

Similar arguments apply to the second sum in (4.7). Since $\mathbb{P}(C_{m,n}) \times \pi_n$, we get

$$\mathbb{E}(|(\hat{V}_1 \cup \hat{V}_2) \cap A_n|^2 \mid C_{n,m})$$

$$\leq C_{24} \sum_{x,y \in A_n} \mathbb{P}_{cr} \left(x \leftrightarrow \partial B(x, 2^{N-2}), y \leftrightarrow \partial B(y, 2^{N-2}) \right).$$

The last sum is bounded from above by $C_{25}n^4\pi_n^2$ (see, e.g., the proof of Theorem 8 in [12]), which along with (4.5) gives (4.6).

5 Proof of Theorems 1.3–1.5

5.1 Proof of Theorem 1.3

We will use the following lemma. For $m, n \ge 1$, and $p \in [0, 1]$, let N(m, n, p) be the number of edges e in the annulus Ann(n, 2n) such that (a) e is connected to $\partial B(e_x, m)$ by two disjoint p-open paths, (b) e^* is connected to $\partial B(e_x, m)^*$ by two disjoint p_c -closed paths, (c) the open and closed paths are disjoint and alternate, and (d) $\tau_e \in [p_c, p]$.

Lemma 5.1 Let m be such that $m \le L(p)$ and $m \le n$. There exists C_1 such that for all t, n,

$$\mathbb{E}(N(m, n, p)^t) \le t! \left(C_1 \frac{n}{m}\right)^{2t}. \tag{5.1}$$

Proof The proof is very similar to the proof of the upper bound in [12, Theorem 8], where we need to use [4, Lemma 6.3] to deal with *p*-open paths. We omit the details.

To continue the proof of Theorem 1.3, define for $n \ge 1$ and k with $0 \le k \le \log^* n$, the event

$$H_{n,k} = \left\{ \begin{array}{l} \text{There exists a } p_n(k) - \text{ open circuit in } Ann(n/4, n/2) \\ \text{which is connected to infinity by a } p_n(k) - \text{ open path.} \end{array} \right\}, \quad (5.2)$$



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where $p_n(k)$ is defined in (2.3). Let us decompose the tth moment of O(n, 2n) according to the events $H_{n,k}$. By (2.4) and (2.8), there exists C_2 , C_3 such that for all n, k,

$$\mathbb{P}(H_{n,k}^c) \le C_2(\log^{(k-1)} n)^{-C_*C_3}. \tag{5.3}$$

Writing $n_k = \frac{n}{C_* \log^{(k)} n}$ and using the Cauchy-Schwarz inequality for $1 < k < \log^* n$,

$$\mathbb{E}(O(n,2n)^{t}; H_{n,k}, H_{n,k+1}^{c}) \leq (\mathbb{E}(N(n_{k}, n, p_{n}(k)))^{2t})^{1/2} (\mathbb{P}(H_{n,k+1}^{c}))^{1/2}$$

$$\leq ((2t)!)^{1/2} (C_{1}C_{*}\log^{(k)} n)^{2t} C_{2}^{1/2} (\log^{(k)} n)^{-\frac{C_{*}C_{3}}{2}}$$

$$= (C_{2}(2t)!)^{1/2} (C_{*}C_{1})^{2t} (\log^{(k)} n)^{\frac{4t-C_{*}C_{3}}{2}}.$$

Choosing $C_* = \frac{4t+2}{C_3}$, this becomes

$$(C_2(2t)!)^{1/2} \left(\frac{(4t+1)C_1}{C_3} \right)^{2t} (\log^{(k)} n)^{-1} \le (C_4t)^{3t} (\log^{(k)} n)^{-1}$$

for some C_4 . For the case k = 0, we have

$$\mathbb{E}(O(n,2n)^t; H_{n,1}^c) \le n^{2t} \mathbb{P}(H_{n,1}^c) \le \frac{C_2}{n} \le C_2.$$

If we sum over k and bound $\sum_{k=1}^{(\log^* n)-1} (\log^{(k)} n)^{-1}$ independent of n as in [10, (2.26)], we get

$$\mathbb{E}(O(n,2n)^t) \le (Ct)^{3t}.$$

5.2 Proof of Theorem 1.4

Proof of upper bound Consider the event A that, for all large n, for all $1 \le i \le n$, the annulus $Ann(2^i, 2^{i+c\log n})$ contains a p_c -open circuit around the origin. Note that $\mathbb{P}(A) = 1$ for large enough c. We assume that c is an integer. Then $2^{c\log n} = n^c$ is an integer too.

In the annulus $Ann(2^i, 2^{i+2c\log n+1})$, we define the graph \mathcal{G}_i^n as follows. Let \mathcal{U} be the union of p_c -open clusters in $Ann(2^i, 2^{i+2c\log n+1})$ attached to $\partial B(2^{i+2c\log n+1})$. In particular, we assume that all the sites in $\partial B(2^{i+2c\log n+1})$ are in \mathcal{U} . If \mathcal{U} contains a path from $B(2^i)$ to $\partial B(2^{i+2c\log n+1})$, we define \mathcal{G}_i^n as \mathcal{U} . Otherwise, we consider the invasion percolation cluster \mathcal{I} in $Ann(2^i, 2^{i+2c\log n+1})$ of the invasion percolation process with $G_0 = B(2^i)$ (that is $B(2^i)$ is assumed to be invaded at step 0) terminated at the first time a site from \mathcal{U} is invaded, and define \mathcal{G}_i^n as $\mathcal{I} \cup \mathcal{U}$. We say that an edge e is e is e is e in e



Let X_i^n be the number of disconnecting edges for \mathcal{G}_i^n in $Ann(2^{i+c\log n}, 2^{i+c\log n+1})$. Note that if the event A occurs then, for all large n, X_i^n dominates $O(2^{i+c\log n}, 2^{i+c\log n+1})$, the number of outlets of the IPC \mathcal{S} of the origin in $Ann(2^{i+c\log n}, 2^{i+c\log n+1})$. Moreover, for any $i < \lfloor 3c\log n \rfloor$, $(X_{i+k\lfloor 3c\log n \rfloor}^n)_{k=0}^{\lfloor n/3c\log n \rfloor - 1}$ are independent and the reader can verify that the proof of Theorem 1.3 is valid when the number of outlets is replaced with X_i^n . Therefore, there exist constants $\lambda > 0$ and $C_5 < \infty$ so that, for all n and i,

$$\mathbb{E}\exp\left(\lambda\left(X_{i}^{n}\right)^{1/3}\right) < C_{5}.$$

Let Y_i be a sequence of independent integer-valued random variables with $\mathbb{P}(Y_i > n) = \min\{1, C_5 e^{-\lambda n^{1/3}}\}$. Then, for any $i < \lfloor 3c \log n \rfloor$, $(X_{i+k \lfloor 3c \log n \rfloor}^n)_{k=0}^{\lfloor n/3c \log n \rfloor - 1}$ is stochastically dominated by $(Y_k)_{k=0}^{\lfloor n/3c \log n \rfloor - 1}$. In particular,

$$\mathbb{P}\left(\sum_{i=1}^{n} X_{i}^{n} > C_{6}n\right) \leq 3c \log n \mathbb{P}\left(\sum_{i=1}^{\lfloor n/3c \log n \rfloor - 1} Y_{i} > C_{6}n/3c \log n\right)$$

$$\leq C_{7} \log n \exp(-C_{8}n^{C_{9}}).$$

The last inequality follows, for example, from [15]. Therefore, a.s., for all large n, $\sum_{i=1}^{n} X_i^n \leq C_6 n$.

Note that, if the event A occurs, then, for all large n,

$$O(2^{c \log n}, 2^n) \le \sum_{i=1}^n X_i^n \le C_6 n.$$

Finally, since the event A occurs with probability one,

$$O(2^n) \le O(2^{c \log n}, 2^n) + O(2^{c \log(c \log n)}, 2^{c \log n}) + |B(2^{c \log(c \log n)})| \le C_{10}n.$$

Proof of lower bound For $i \ge 1$, let G_i be the event that there is no p_{2^i} -closed dual circuit around the origin with radius larger than $2^{i+\log i}$, and let G be the event that G_i occurs for all but finitely many i. It is easy to see [using inequality (2.8)] that $\mathbb{P}(G) = 1$.

For $i \geq 1$, let K_i be the event that (a) there exists a p_{2^i} -closed dual circuit $\mathcal C$ around the origin in $Ann(2^i, 2^{i+1})^*$, (b) there exists a p_c -open circuit $\mathcal C'$ around the origin in $Ann(2^i, 2^{i+1})$ and (c) the circuit $\mathcal C'$ is connected to $\partial B(2^{i+\log i})$ by a p_{2^i} -open path. See Fig. 2 for an illustration of the event $G_i \cap K_i$. Note that $\mathcal C'$ is in $B(2^{i+1}) \cap ext(\mathcal C)$. By RSW theorem and (2.6),

$$\mathbb{P}(K_i) > C_{11} > 0,$$



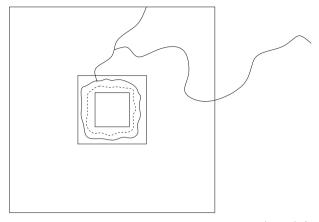


Fig. 2 The event $G_i \cap K_i$. The *boxes*, in order from smallest to largest, are $B(2^i)$, $B(2^{i+1})$, and $B(2^{i+\log i})$. Because there is no p_{2^i} -closed circuit around the origin of radius larger than $2^{i+\log i}$, the p_{2^i} -open path which connects $\partial B(2^{i+\log i})$ to the circuit in $Ann(2^i, 2^{i+1})$ must be connected to ∞ by a p_{2^i} -open path

for some C_{11} that does not depend on i. Fix an integer n, and let j be an integer between 1 and $\log n$. We consider events $K_i^j = K_{j+i\log n}$. Note that, for any fixed j, the events $(K_i^j)_{i=0}^{\lfloor \lfloor n/\log n\rfloor -1}$ are independent.

Let $X_i^j = I_{K_i^j}$. Recall that $\mathbb{P}(X_i^j = 1) > C_{11}$. We need the following lemma. Its proof is standard, so we omit it.

Lemma 5.2 Let c > 0. There exist $\alpha > 0$ and $\beta < 1$ depending on c with the following property. If X_i are independent 0/1 random variables (not necessarily identically distributed) with $\mathbb{P}(X_i = 1) > c$ for all i, then for all n,

$$\mathbb{P}\left(\sum_{i=1}^n X_i < \alpha n\right) < \beta^n.$$

It follows that there exist $\alpha > 0$ and $\beta < 1$ such that for any n and $1 \le j \le \log n$

$$\mathbb{P}\left(\sum_{i=0}^{\lfloor n/\log n\rfloor -1} X_i^j < \frac{\alpha n}{\log n}\right) < \beta^{n/\log n}.$$

Therefore,

$$\mathbb{P}\left(\sum_{j=1}^{\log n}\sum_{i=0}^{\lfloor n/\log n\rfloor-1}X_i^j < \alpha n\right)$$

$$\leq \mathbb{P}\left(\sum_{i=0}^{\lfloor n/\log n\rfloor-1}X_i^j < \alpha n/\log n \text{ for some } j \in [1,\log n]\right) \leq \log n\beta^{n/\log n}.$$



In particular, it follows from Borel–Cantelli's lemma that, with probability one, for all large n,

$$\sum_{i=1}^n I_{K_i} \geq \alpha n.$$

Finally, observe that the event G occurs with probability one, and the event $G_i \cap K_i$ implies that there exists an outlet in $Ann(2^i, 2^{i+1})$. The lower bound in (1.6) follows.

Proof of Corollary 1.1 The inequalities (1.7) follow immediately from those in Theorem 1.4. Therefore we will only prove (1.8). First we show the upper bound.

Choose c_5 from (1.7). Using (2.4) and (2.8), we can show that if C_* is made sufficiently large, then with probability one, for all large n, after the invasion has reached $\partial B(n)$, the weight of each further accepted edge is no larger than $p_n(1)$, where $p_n(1)$ is defined in (2.3). Therefore, for all large k,

$$\hat{\tau}_k - p_c \le p_{(c_5)^k}(1) - p_c.$$

Since there exists C_{12} , $C_{13} > 0$ such that for all n, $p_n(1) - p_c \le C_{12}n^{-C_{13}}$ [use (2.9)] and the fact that the 4-arm exponent is strictly smaller than 2 (see, e.g., Section 6.4 in [21]), we have

$$\hat{\tau}_k - p_c \le C_{12}(c_5)^{-C_{13}k},$$

proving the upper bound. To show the lower bound, choose c_6 from (1.7). For a < 1, we obtain

$$\begin{split} \mathbb{P}(\hat{\tau}_k < p_c + a^k, \, \hat{R}_k < (c_6)^k) & \stackrel{\leq}{\leftarrow} \mathbb{P}(B((c_6)^k) \stackrel{p_c + a^k}{\longleftrightarrow} \infty) \\ & \leq C_{14} \mathbb{P}_{cr}(B((c_6)^k) \leftrightarrow \partial B(L(p_c + a^k))) \\ & \leq C_{15} \left(\frac{(c_6)^k}{a^{-C_{16}k}}\right)^{C_{17}} \leq C_{18} e^{-C_{19}k}, \end{split}$$

for constants $C_{14} - C_{19}$, where the last inequality holds for small enough a. The second inequality follows from (2.5). The third one follows from, for example, [9, eq. 11.90] and the fact that $L(p) > (p - p_c)^{-\delta}$ for some $\delta > 0$ (see, e.g., Cor. 1 and eq. 2.3 from [13]). Borel–Cantelli's lemma gives the lower bound of (1.8).

5.3 Proof of Theorem 1.5

Given any nonempty subinterval of (0, 1], we will show that with probability one, $(\frac{\hat{\tau}_{k+1}-p_c}{\hat{\tau}_k-p_c})$ is in this subinterval for infinitely many k. We will use the following fact. From [13, (4.35)] it follows that, for any a > 0,

$$L(p_c + a\delta) \simeq L(p_c + \delta).$$



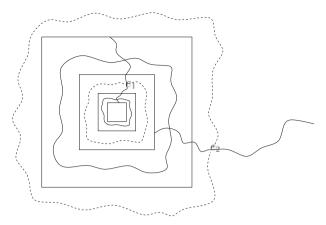


Fig. 3 The event D_n . The boxes, from smallest to largest, are B(n), B(2n), B(4n), and B(8n). The edge e_1 is connected to both B(n) and $\partial B(4n)$. The edge e_2 is connected to both B(4n) and infinity. The solid curves represent occupied paths and the dotted curves represent vacant dual paths. In the figure, both e_1 and e_2 are outlets of the invasion

The constants above depend on a but do not depend on δ so long as δ is sufficiently small.

Pick a nonempty interval $[x, y] \subset (0, 1]$ and choose a, b > 0 such that

$$a < b$$
 and $1 < bx < ay$.

We consider the event D_n that there exist p_c -open circuits around the origin in the annulus Ann(n, 2n); and in the annulus A(4n, 8n), and there exist two edges, $e_1 \in Ann(2n, 3n)$ and $e_2 \in Ann(8n, 9n)$ such that

- there is a p_c -open path connecting one of the ends of e_1 to B(n), and there is a p_c -open path connecting the other end of e_1 to $\partial B(8n)$;
- there is a p_c -open path connecting one of the ends of e_2 to B(4n), and there is a p_n -open path connecting the other end of e_2 to infinity;
- there is a $(p_c + b(p_n p_c))$ -closed path P_1 in the dual lattice inside $Ann(2n, 3n)^*$ connecting the ends of e_1^* such that $P_1 \cup \{e_1^*\}$ is a circuit around the origin;
- there is a $(p_c + ay(p_n p_c))$ -closed path P_2 in the dual lattice inside $Ann(8n, 9n)^*$ connecting the ends of e_2^* such that $P_2 \cup \{e_2^*\}$ is a circuit around the origin;
- the weight $\tau_{e_1} \in (p_c + a(p_n p_c), p_c + b(p_n p_c))$, and the weight $\tau_{e_2} \in (p_c + bx(p_n p_c), p_c + ay(p_n p_c))$.

See Fig. 3 for an illustration of the event D_n . By RSW arguments and [4, Lemma 6.3] (similar to the proof of [4, Corollary 6.2]), there exists a constant $C_{20} > 0$ which depends on a, b, x, and y but not on n such that

$$\mathbb{P}(D_n) \ge C_{20}. \tag{5.4}$$

Since $\limsup_n D_n$ does not depend on the states of finitely many edges, $\mathbb{P}(\limsup_n D_n) \in \{0, 1\}$. Assume that this probability is 0. Then there exists N (deterministic)



such that

$$\mathbb{P}(D_n \text{ occurs for some } n \geq N) < C_{20}/2.$$

But this probability is, in fact, at least $\mathbb{P}(D_N)$. This contradicts (5.4). Therefore

$$\mathbb{P}(\limsup_{n} D_n) = 1. \tag{5.5}$$

Note that the event D_n implies that there exists k=k(n) such that e_1 and e_2 are respectively the kth and (k+1)st outlets of the invasion. In particular, using the above bounds for τ_{e_1} and τ_{e_2} , $\frac{\hat{\tau}_{k+1}-p_c}{\hat{\tau}_k-p_c} \in [x,y]$. Combining this with (5.5), we get $\mathbb{P}(\frac{\hat{\tau}_{k+1}-p_c}{\hat{\tau}_k-p_c} \in [x,y]$ for infinitely many k)=1. This completes the proof.

6 Proof of Theorem 1.6

Since the proof is very similar to the proof of Theorem 3 in [12], we only sketch the main ideas. From now on we fix $\sigma \in \{\text{open, closed}\}^{2m}$, and assume that σ consists of m 'open' and m 'closed'.

The RSW theorem implies that there exists $\delta > 0$ such that for all N,

 \mathbb{P}_{cr} (there exists an occupied circuit in $Ann(N, 2N) \geq \delta$.

Since events depending on the state of edges in disjoint annuli are independent, we can find an increasing sequence N_i such that

$$\alpha_i = \mathbb{P}_{cr}$$
 (there exists an occupied circuit in $Ann(N_i, N_{i+1}) \rightarrow 1$,

as $i \to \infty$. We fix the sequence N_i and write A_i for $Ann(N_i, N_{i+1})$.

$$\mathbb{P}_{cr}(E \cap \{B(l) \leftrightarrow_{\sigma} \partial B(n)\}) = \mathbb{P}_{cr}(E \cap \{B(l) \leftrightarrow_{\sigma} \partial B(n)\} \cap (F_i^{(m)})^c) + \sum_{\mathcal{C} \subset A_i} \sum_{e_1, \dots, e_m \in \mathcal{O} \cap \mathbb{E}^2} \mathbb{P}_{cr}(E \cap \{B(l) \leftrightarrow_{\sigma} \partial B(n)\} \cap F_i(\mathcal{C}_{e_1, \dots, e_m})).$$

Let $\{B(l) \leftrightarrow_{\sigma} \mathcal{C}_{e_1,\dots,e_m}\}$ denote the event that B(l) is σ -connected to \mathcal{C} so that the m disjoint closed dual paths connect $B(l)^*$ to the edges e_1^*,\dots,e_m^* in the interior of \mathcal{C} . Similarly, let $\{\mathcal{C}_{e_1,\dots,e_m} \leftrightarrow_{\sigma} \partial B(n)\}$ denote the event that \mathcal{C} is σ -connected to $\partial B(n)$



so that the m disjoint closed dual paths connect $\partial B(n)^*$ to the edges e_1^*, \ldots, e_m^* in the exterior of C.

We now estimate the probability $\mathbb{P}_{cr}(E \cap \{B(l) \leftrightarrow_{\sigma} \partial B(n)\} \cap (F_i^{(m)})^c)$. By Menger's theorem [6, Theorem 3.3.1], the event $\{B(l) \leftrightarrow_{\sigma} \partial B(n)\} \cap (F_i^{(m)})^c$ implies that there exist (m+1) disjoint closed crossings of the annulus A_i . We use Reimer's inequality [19] to conclude that the probability $\mathbb{P}_{cr}(E \cap \{B(l) \leftrightarrow_{\sigma} \partial B(n)\} \cap (F_i^{(m)})^c)$ is bounded from above by

$$\mathbb{P}_{cr}(B(l) \leftrightarrow_{\sigma} \partial B(n)) \mathbb{P}_{cr}(\text{ there exists a closed crossing of } A_i)$$

 $\leq (1 - \alpha_i) \mathbb{P}_{cr}(B(l) \leftrightarrow_{\sigma} \partial B(n)).$

We have just shown how a statement similar to (17) in [12] is obtained. An analogous statement to (18) in [12] is also valid. The remainder of the proof is similar to the proof of Kesten [12], where in the proof of the statement analogous to Lemma 23 in [12] we use extensions of arm separation techniques from [18, Section 4]. We use the following analogue of Kesten's Lemma 23.

Lemma 6.1 Consider circuits C in annulus A_i , D in annulus A_{i+3} , sets of edges e_1, \ldots, e_m on C and f_1, \ldots, f_m on D respectively. Let P(C, D) be the probability, conditional on the event that all edges in $C \setminus \{e_1, \ldots, e_m\}$ are open and e_1, \ldots, e_m are closed, that (1) there are disjoint closed dual paths from e_i^* to f_i^* , (2) there are m disjoint open paths that connect C to D such that, for any two of them, there is a closed dual path (one of the paths from (1)) between them, (3) D is the innermost open circuit with defects f_1, \ldots, f_m in annulus A_{i+3} , (4) there is an open circuit with m defects in annulus A_{i+2} . (Dependence on the edges e_i and f_i is suppressed in the notation.) We similarly define C', D', etc. There exists a finite constant C_1 that may depend only on m (it does not depend on particular choice of circuits or defects) such that

$$\frac{P(\mathcal{C},\mathcal{D})P(\mathcal{C}',\mathcal{D}')}{P(\mathcal{C},\mathcal{D}')P(\mathcal{C}',\mathcal{D})} < C_1.$$

To prove Lemma 6.1, we need the following extension of Kesten's arm separation [13, Lemmas 4 and 5]. Let \mathcal{I} be a fixed partition of $\partial B(1)$ (in \mathbb{R}^2) into 2m disjoint connected subsets \mathcal{I}_i , each of diameter at least 1/(2m) (ordered clockwise). Let $\mathcal{I}(s)$ be the corresponding partition of $\partial B(s)$ into 2m disjoint connected subsets $\mathcal{I}_i(s) = s\mathcal{I}_i = \{sx : x \in \mathcal{I}_i\}$. Let $\mathcal{I}(n, n')$ be the partition of $\overline{Ann(n, n')}$ into 2m disjoint connected subsets $\mathcal{I}_i(n, n') = \bigcup_{n \leq s \leq n'} \mathcal{I}_i(s)$.

Lemma 6.2 (external arm separation) Let n_0 and n be positive integers with $n_0 \le n-3$. We consider a circuit C in $B(2^{n_0})$ and a set of edges e_1, \ldots, e_m on C. Let $E(C_{e_1,\ldots,e_m})$ be the event that (1) the edges in $C \setminus \{e_1,\ldots,e_m\}$ are open and e_1,\ldots,e_m are closed, (2) there are m disjoint closed dual paths from e_j^* to $\partial B(2^n)^*$, (3) there are m disjoint open paths from C to the boundary of $B(2^n)$ in $(B(2^n) \setminus int(C)) \setminus \{e_1,\ldots,e_m\}$ such that these paths alternate with the closed paths defined in (2). Let $\widetilde{E}(C_{e_1,\ldots,e_m})$ be the event that $E(C_{e_1,\ldots,e_m})$ occurs with 2m paths P_1,\ldots,P_{2m} (ordered clockwise, all



paths with odd indices are closed, and the ones with even indices are open) satisfying the requirement that, for all $1 \le i \le 2m$, $P_i \cap \overline{Ann(2^{n-1}, 2^n)} \subset \mathcal{I}_i(2^{n-1}, 2^n)$. Then

$$\mathbb{P}(E(\mathcal{C}_{e_1,\ldots,e_m})) \leq C_2 \mathbb{P}(\widetilde{E}(\mathcal{C}_{e_1,\ldots,e_m})),$$

where the constant C_2 may depend on m but not on n, n_0 , or the choice of circuit.

Remark 8 The event \widetilde{E} is reminiscent of the event Δ in [13, p. 127 and Figure 8].

Remark 9 It is actually believed [7] and is the aim of ongoing work of Garban and Pete that a much stronger statement holds: given any configuration inside $B(2^{n_0})$, if we condition on the existence of m open paths and m closed dual paths from a neighborhood of the origin to $\partial B(2^n)$ and these paths are alternating, then they will be well-separated (refer to [13] for this definition) on $\partial B(2^n)$ with positive probability independent of n and the configuration inside $B(2^{n_0})$.

Lemma 6.3 (internal arm separation) Let n and n_1 be positive integers with $n+3 \le n_1$. Consider a circuit \mathcal{D} in $B(2^{n_1})^c$ and a set of edges f_1, \ldots, f_m on \mathcal{D} . Let $F(\mathcal{D}_{f_1,\ldots,f_m})$ be the event that (1) the edges in $\mathcal{D}\setminus\{f_1,\ldots,f_m\}$ are open and f_1,\ldots,f_m are closed, (2) there are m disjoint closed dual paths from f_j^* to $B(2^n)^*$, (3) there are m disjoint open paths from \mathcal{D} to $B(2^n)$ in $\overline{int(\mathcal{D})}$ such that these paths alternate with the closed dual paths defined in (2). Let $\widetilde{F}(\mathcal{D}_{f_1,\ldots,f_m})$ be the event that the event $F(\mathcal{D}_{f_1,\ldots,f_m})$ occurs with 2m paths P_1,\ldots,P_{2m} (ordered clockwise, all paths with odd indices are closed, and the ones with even indices are open) satisfying the requirement that, for all 1 < i < 2m, $P_i \cap \overline{Ann(2^n, 2^{n+1})} \subset T_i(2^n, 2^{n+1})$. Then

$$\mathbb{P}(F(\mathcal{D}_{f_1,\ldots,f_m})) \leq C_3 \mathbb{P}(\widetilde{F}(\mathcal{D}_{f_1,\ldots,f_m})),$$

where the constant C_3 may depend on m but not on n, n_1 , or the choice of circuit.

The proofs of Lemmas 6.2 and 6.3 are similar, and we only give the proof of Lemma 6.2 here. Moreover, parts of the proof of Lemma 6.2 are similar to the proof of Lemma 4 in [13]. We will refer the reader to [13] for the proof of those parts. Before we give the proof of Lemma 6.2, we show how to deduce Lemma 6.1 from the above two lemmas. Using Lemmas 6.2, 6.3 and "gluing" arguments (see [13,18]), we prove

Lemma 6.4 For two circuits, C_1 in annulus A_i and D_1 in annulus A_{i+2} , sets of edges e_1, \ldots, e_m on C_1 and f_1, \ldots, f_m on D_1 , if $M(C_1, D_1)$ is the probability, conditioned on the event that all edges in $C_1 \setminus \{e_1, \ldots, e_m\}$ and in $D_1 \setminus \{f_1, \ldots, f_m\}$ are open and e_1, \ldots, e_m , f_1, \ldots, f_m are closed, that there are disjoint closed dual paths from e_i^* to f_i^* for all i, and there are m disjoint open paths from C_1 to D_1 in $\overline{int}(D_1) \setminus int(C_1)$, which alternate with the closed dual paths defined above (and similar definitions for C_2 and D_2), then

$$\frac{M(\mathcal{C}_1,\mathcal{D}_1)M(\mathcal{C}_2,\mathcal{D}_2)}{M(\mathcal{C}_1,\mathcal{D}_2)M(\mathcal{C}_2,\mathcal{D}_1)} < C_4,$$

for some constant C_4 that does not depend on the particular choice of circuits or defects. (Dependence on the edges e_i and f_i is suppressed in the notation.)



Proof This lemma follows from Lemmas 6.2, 6.3, the RSW theorem (Section 11.7 in [9]), and the generalized FKG inequality [13, Lemma 3]. For more details we refer the reader to the proof of (2.43) in [13].

Proof of Lemma 6.1 Consider circuits C_1 in annulus A_{i+2} , D_1 in A_{i+3} , sets of edges g_1, \ldots, g_m on C_1 and h_1, \ldots, h_m on D_1 respectively. Let $H(C_1, D_1)$ be the probability of the event that (1) C_1 is the outermost open circuit with defects g_1, \ldots, g_m in annulus A_{i+2} , (2) D_1 is the innermost open circuit with defects h_1, \ldots, h_m in annulus A_{i+3} , (3) there are disjoint closed dual paths from g_i^* to h_i^* , and (4) there are m disjoint open paths from C_1 to D_1 in $\overline{int(D_1)} \setminus int(C_1)$, which alternate with the closed dual paths defined above. (Dependence on the edges g_i and h_i is suppressed in the notation.)

We write

$$P(\mathcal{C}, \mathcal{D})P(\mathcal{C}', \mathcal{D}') = \sum_{\mathcal{C}_1} M(\mathcal{C}, \mathcal{C}_1)H(\mathcal{C}_1, \mathcal{D}) \sum_{\mathcal{C}'_1} M(\mathcal{C}', \mathcal{C}'_1)H(\mathcal{C}'_1, \mathcal{D}').$$

We then apply the previous lemma to C, C_1 , C' and C'_1 .

Proof of Lemma 6.2 We only consider the case m=2. The case m=1 is simpler, and the general case is similar to the case m=2. Fix a circuit \mathcal{C} in $B(2^{n_0})$ and edges e, f on \mathcal{C} , and assume that the event $E(\mathcal{C}_{e,f})$ occurs.

We define γ_1^l as the leftmost closed dual path from e^* to $\partial B(2^{n_0}+1/2)$ in $B(2^{n_0}+1/2)$ in the first vertex on $\partial B(2^{n_0})$ to the left of γ_1^l as a_1 , and the first vertex on $\partial B(2^{n_0})$ to the right of γ_1^r as a_2 . Let γ_2^l be the leftmost open path from the right end-vertex of e (using the clockwise ordering of vertices end edges on C) to a_2 . This path is necessarily contained in $B(2^{n_0})$ int(C). Let C0. Let C0 be the rightmost open path from the left end-vertex of E1 (using the clockwise ordering of vertices end edges on C2 to C1 in C2 (see Fig. 4).

For $i \in \{1, 2, 3, 4\}$, let T_i be the piece of $\partial B(2^{n_0})$ between (and including) a_i and a_{i+1} that does not contain a_{i+2} or a_{i+3} , where we use the convention $a_i = a_{i-4}$ for i > 4. Note that it is possible that $a_2 = a_3$ (in which case $T_2 = \{a_2\}$) or $a_4 = a_1$ (in which case $T_4 = \{a_4\}$); however, we necessarily have $a_1 \neq a_2$ and $a_3 \neq a_4$.

Let γ_i be the part of $\gamma_i^l \cup \gamma_i^r \cup \mathcal{C}$ that consists of the piece of γ_i^l from the last intersection with $\gamma_i^r \cup \mathcal{C}$, the piece of γ_i^r from the last intersection with $\gamma_i^l \cup \mathcal{C}$, and the piece of \mathcal{C} that connects the first two pieces (if the pieces are disconnected). Note that it is possible that γ_2 or γ_4 is a single point set on $\partial B(2^{n_0})$, which happens if $a_2 = a_3$ or $a_4 = a_1$, respectively. Let R_i denote the connected subset of \mathbb{R}^2 with the boundary that consists of T_i and γ_i (see Fig. 4). Note that these sets are disjoint. Moreover, if γ_2 or γ_4 is a single point set ($\{a_2\}$ or $\{a_4\}$, respectively), then R_2 or R_4 is the same single point set. Let $S := B(2^{n_0})^c \cup R_1 \cup R_2 \cup R_3 \cup R_4$. Note that once $\gamma_1, \gamma_2, \gamma_3$, and γ_4 are fixed, the percolation process in S is still an independent Bernoulli percolation.

Let $E(\gamma_1, \ldots, \gamma_4)$ be the event that (1) γ_1 and γ_3 are connected to $\partial B(2^n)^*$ by closed dual paths P_1 and P_3 in S, and (2) γ_2 and γ_4 are connected to $\partial B(2^n)$ by open paths P_2 and P_4 in S. Let $\widetilde{E}(\gamma_1, \ldots, \gamma_4)$ be the event that $E(\gamma_1, \ldots, \gamma_4)$ occurs with paths



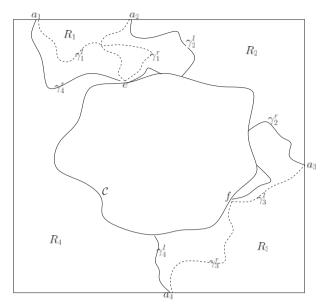


Fig. 4 The event $E(C_{e_1,...,e_m})$ occurs if and only if γ_1 and γ_3 are connected by closed paths to $\partial B(2^n)^*$, and γ_2 and γ_4 are connected by open paths to $\partial B(2^n)$ in $B(2^{n_0})^c \cup R_1 \cup R_2 \cup R_3 \cup R_4$

 P_1, \ldots, P_4 satisfying the requirement that, for all $1 \le i \le 4$, $P_i \cap \overline{Ann(2^{n-1}, 2^n)} \subset \mathcal{I}_i(2^{n-1}, 2^n)$. Lemma 6.2 follows if there exists a constant C_2 which does not depend on n, n_0 , or the choice of γ_i 's, such that

$$\mathbb{P}(E(\gamma_1, \dots, \gamma_4)) \le C_2 \mathbb{P}(\widetilde{E}(\gamma_1, \dots, \gamma_4)). \tag{6.1}$$

If T_1, \ldots, T_4 are comparable in size, the proof of (6.1) is essentially the same as the proof of Lemma 4 in [13]. If T_1, \ldots, T_4 are of different scales, the proof of (6.1) is similar in spirit to the proof of Lemma 4 in [13], but more involved. We indicate the differences below. We first construct a family of disjoint annuli in four stages. We define

$$l_i(1) = \min\{l : \exists x \in 2^l \mathbb{Z}^2 \cap \partial B(2^{n_0}) \text{ s.t. } B(x, 2^l) \supset T_i\}$$

if such l exists (the definition implies that it is no bigger than n_0), and let $B_i(1) = B(x_i(1), 2^{l_i(1)})$ be such a box. If there are several choices for the box, we pick the first one in clockwise ordering. If there are no such l, we let $l_i(1) = n_0 + 1$ and $B_i(1) = B(2^{n_0+1})$ (in this case $x_i(1) = 0$). The boxes $B_1(1), \ldots, B_4(1)$ form a covering of $\partial B(2^{n_0})$ such that $B_i(1) \supset T_i$, and either $x_i(1) \in T_i$ or $x_i(1) = 0$ (in which case $B_i(1) = B(2^{n_0+1})$). We also define

$$\tilde{l}_i(1) = \min\{l \ge l_i(1) : B(x_i(1), 2^l) \supset B_j(1) \text{ for some } j \ne i\}.$$



Let $Ann_i(1) = Ann(x_i(1); 2^{l_i(1)}, 2^{\widetilde{l}_i(1)-3}) = B(x_i(1), 2^{\widetilde{l}_i(1)-3}) \setminus B_i(1)$, if $\widetilde{l}_i(1) - 5 > l_i(1)$; otherwise, let $Ann_i(1) = \emptyset$. Note that if $l_i(1) \geq l_{i-1}(1)$ or $l_i(1) \geq l_{i+1}(1)$, then $Ann_i(1) = \emptyset$. In particular, if $l_i(1) = n_0 + 1$, then $Ann_i(1) = \emptyset$. If $Ann_i(1) \neq \emptyset$, we let $\widetilde{B}_i(1) = B(x_i(1), 2^{\widetilde{l}_i(1)-3})$; otherwise, we let $\widetilde{B}_i(1) = B_i(1)$. We write $\widetilde{B}_i(1) = B(x_i(1), 2^{l_i(1)'})$. We observe that the boxes $\widetilde{B}_i(1)$ form a covering of $\partial B(2^{n_0})$, and that there exists i such that $|l_{i+1}(1)' - l_i(1)'| \leq 5$, in other words, $\widetilde{B}_{i+1}(1)$ and $\widetilde{B}_i(1)$ are comparable in size. This completes the first stage of our construction.

We proceed further by defining

$$l_i(2) = \min\{l : \exists x \in 2^l \mathbb{Z}^2 \cap \partial B(2^{n_0}) \text{ s.t. } B(x, 2^l) \supset (\widetilde{B}_i(1) \cup \widetilde{B}_{i+1}(1))\}$$

if such l exists (it is necessarily not bigger than n_0), and let $B_i(2) = B(x_i(2), 2^{l_i(2)})$ be such a box (if there are several choices, we pick the first one in clockwise ordering). If there are no such l, we let $l_i(2) = n_0 + 1$ and $B_i(2) = B(2^{n_0+1})$ (in this case $x_i(2) = 0$). We also define

$$\tilde{l}_i(2) = \min\{l \ge l_i(2) : B(x_i(2), 2^l) \supset \tilde{B}_j(1) \text{ for some } j \ne i, i+1\}.$$

Let $Ann_i(2) = Ann(x_i(2); 2^{l_i(2)}, 2^{\widetilde{l_i}(2)-3}) = B(x_i(2), 2^{\widetilde{l_i}(2)-3}) \setminus B_i(2)$, if $\widetilde{l_i}(2)-5 > l_i(2)$; otherwise, let $Ann_i(2) = \emptyset$. If $Ann_i(2) \neq \emptyset$, we define $\widetilde{B_i}(2) = \widetilde{B_{i+1}}(2) = B(x_i(2), 2^{\widetilde{l_i}(2)-3})$. For all remaining indices i for which $\widetilde{B_i}(2)$ is not yet defined, we let $\widetilde{B_i}(2) = \widetilde{B_i}(1)$. In other words, if we have not succeeded in building a nonempty annulus around $\widetilde{B_i}(1)$, we take this box unchanged to the next stage of our construction. We write $\widetilde{B_i}(2) = B(x_i(2), 2^{l_i(2)'})$. We observe that the boxes $\widetilde{B_i}(2)$ form a covering of $\partial B(2^{n_0})$, and that there exists i such that $|l_{i+1}(2)' - l_i(2)'| \leq 5$, $|l_{i+2}(2)' - l_i(2)'| \leq 5$ and $|l_{i+2}(2)' - l_{i+1}(2)'| \leq 5$. In other words $\widetilde{B_i}(2)$, $\widetilde{B_{i+1}}(2)$ and $\widetilde{B_{i+2}}(2)$ are comparable in size.

In the third stage, for any i, we define

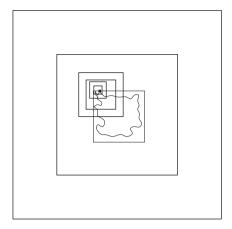
$$l_i(3) = \min\{l : \exists x \in 2^l \mathbb{Z}^2 \cap \partial B(2^{n_0}) \text{ s.t. } B(x, 2^l) \supset (\widetilde{B}_i(2) \cup \widetilde{B}_{i+1}(2) \cup \widetilde{B}_{i+2}(2))\}$$

if such l exists (it is necessarily not bigger than n_0), and let $B_i(3) = B(x_i(3), 2^{l_i(3)})$ be such a box (if there are several choices, we pick the first one in clockwise ordering). If there are no such l, we let $l_i(3) = n_0 + 1$ and $B_i(3) = B(2^{n_0+1})$ (in this case $x_i(3) = 0$). We also define

$$\tilde{l}_i(3) = n_0 + 1.$$

Let $Ann_i(3) = Ann(x_i(3); 2^{l_i(3)}, 2^{\widetilde{l}_i(3)-3}) = B(x_i(3), 2^{\widetilde{l}_i(3)-3}) \setminus B_i(3)$, if $\widetilde{l}_i(3) - 5 > l_i(3)$, otherwise let $Ann_i(3) = \emptyset$. If $Ann_i(3) \neq \emptyset$, we define $\widetilde{B}_i(3) = \widetilde{B}_{i+1}(3) = \widetilde{B}_{i+2}(3) = B(x_i(3), 2^{\widetilde{l}_i(3)-3})$. For all remaining indices i for which $\widetilde{B}_i(3)$ is not yet defined, we let $\widetilde{B}_i(3) = \widetilde{B}_i(2)$. In other words, if we have not succeeded in building a nonempty annulus around $\widetilde{B}_i(2)$, we take this box unchanged to the next stage of our construction. We write $\widetilde{B}_i(3) = B(x_i(3), 2^{l_i(3)'})$. We observe that the boxes $\widetilde{B}_i(3)$





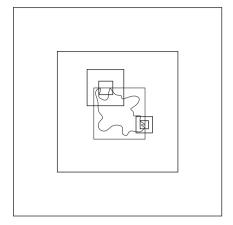


Fig. 5 The family of annuli in the *left figure* consists of one annulus of each level. The family of annuli in the *right figure* consists of two level 1 annuli and one level 4 annulus. In general, there are at most two non-empty level 1 annuli, at most one level 2, 3 or 4 annulus each

form a covering of $\partial B(2^{n_0})$ such that $|l_{i+1}(3)' - l_i(3)'| \le 5$ for all i. Moreover, all of these boxes are contained in $B(2^{n_0+1})$.

Finally, we define the annulus $Ann_i(4) = Ann(2^{n_0+1}, 2^n)$.

Note that $(Ann_i(1))_i$, $(Ann_j(2))_j$, $(Ann_k(3))_k$, and $(Ann_l(4))_l$ are disjoint among levels and between levels. In addition, the event $E(\gamma_1, \ldots, \gamma_4)$ implies the existence of crossings of annulus $Ann_i(1)$ by path P_i , annulus $Ann_i(2)$ by paths P_i and P_{i+1} , annulus $Ann_i(3)$ by paths P_i , P_{i+1} and P_{i+2} , and annulus $Ann_i(4)$ by all four paths P_1, \ldots, P_4 . Some examples of families of annuli are illustrated on Fig. 5.

To show (6.1), our strategy is to bound the probability of the event $E(\gamma_1, \ldots, \gamma_4)$ by the product of probabilities of crossing events in such annuli. We should be more careful though, since we have to take into account that we consider paths in S. For k < 4 and for each nonempty annulus $Ann_i(k)$, we define $a_i(k)$ as the first point on γ_{i-1} (seen as an oriented path from a_{i-1} to a_i) that belongs to $Ann_i(k)$, and $b_i(k)$ as the last point on γ_{i+k} (seen as an oriented path from a_{i+k} to a_{i+k+1}) that belongs to $Ann_i(k)$. Note that such points always exist if $Ann_i(k) \neq \emptyset$. We then define the set $S_i(k)$ as the subset of S with boundary that consists of four pieces (in clockwise order): the piece of $\partial B(x_i(k), 2^{\tilde{l}_i(k)-3})$ between $a_i(k)$ and $b_i(k)$, the piece of $\partial B(x_i(k), 2^{l_i(k)})$ in S, and the piece of γ_{i-1} from $a_i(k)$ to the last intersection with $\partial B(x_i(k), 2^{l_i(k)})$. Let $L_i(k)$ be the common piece of the boundary of $S_i(k)$ and $\partial B(x_i(k), 2^{\tilde{l}_i(k)-3})$ between $a_i(k)$ and $b_i(k)$ (Fig. 6).

For k < 4 and for each nonempty annulus $Ann_i(k)$, let $E_i(k)$ be the event that there exist k disjoint paths from $B(x_i(k), 2^{l_i(k)})$ to $L_i(k)$ in $S_i(k)$ such that the order and the status of paths (occupied or vacant) are induced by the order and the status of P_i, \ldots, P_{i+k-1} . For k < 4 and for each empty annulus $Ann_i(k)$, let $E_i(k)$ be the sure event. Let $E_i(4)$ be the event that the annulus $Ann_i(4)$ is crossed by two open and two closed dual paths such that the open paths are separated by the closed paths.



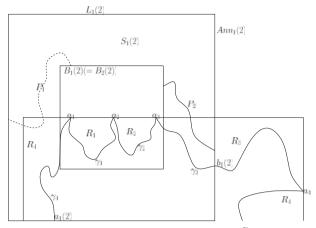


Fig. 6 The *small box* is $B(x_1(2), 2^{l_1(2)})$ and the *big box* is $\partial B(x_1(2), 2^{\tilde{l}_1(2)})$. The event $E_1(2)$ occurs if there is a closed dual path P_1 and an open path P_2 from $B(x_1(2), 2^{l_1(2)})$ to $L_1(2)$ in $S_1(2)$

Since the sets $S_i(k)$ and $Ann_i(4)$ are disjoint, we obtain

$$\mathbb{P}(E(\gamma_1,\ldots,\gamma_4)) \leq \prod_{i,k=1}^4 \mathbb{P}(E_i(k)).$$

It remains to prove an arm separation statement (analogous to Lemmas 4 and 5 in [13]) for each of the crossing events $E_i(k)$. Once this is done, we can proceed similarly to the proof of (2.43) in [13]. We can 'glue' those crossing into solid paths from the γ_i 's to $\partial B(2^n)$ in such a way that the event $\widetilde{E}(\gamma_1, \ldots, \gamma_4)$ occurs, and there exists a constant C_5 such that for any choice of γ_i 's, $\prod_{i,k=1}^4 \mathbb{P}(E_i(k)) \leq C_5 \mathbb{P}(\widetilde{E}(\gamma_1, \ldots, \gamma_4))$. Since this is a standard application of the RSW theorem (Section 11.7 in [9]) and the generalized FKG inequality [13, Lemma 3], we omit the details.

The main difficulty in the proof of the arm separation statement for $E_i(k)$ is that the boundary of $S_i(k)$ is irregular. In the proof of Lemma 4 in [13] it was enough to make sure that with high probability (uniformly in r) all the crossings of $Ann(2^r, 2^{r+1})$ have well-separated extremities. In our case, it is not enough to know that with high probability (uniformly in r) all the crossings of $Ann(x_i(k); 2^r, 2^{r+1})$ in $S_i(k)$ are well-separated. It is possible that their extremities are "trapped" in the sense that they cannot be connected or can only be connected through narrow bottlenecks to $L_i(k)$ in $S_i(k) \setminus B(x_i(k), 2^{r+1})$. An application of the RSW theorem shows that this is very unlikely. The remaining strategy of the proof is similar to the proof of Lemma 4 in [13], and we omit it. The interested reader may find the details in the second version of [5].

7 Proof of Theorem 1.9

We prove only the first statement; the proof of the second is similar (see the proof of the second statement in [10, Theorem 3]). We follow the same method used in



[10, Theorem 3] but because many difficulties arise, we present details of the entire proof. Pick an edge e and let n=2|e|/3 (this choice makes e in the middle of Ann(n,2n)). Let $\epsilon>0$.

Step 1. First we give a lower bound for the probability that $e \in \mathcal{O}$. The constant C_* will be determined later. Consider the event D_e that

- 1. there exist p_c -open circuits around the origin in the annuli Ann(n/2, n) and Ann(2n, 4n);
- 2. there exist two disjoint p_c -open paths, one connecting e_y to the circuit in Ann(2n, 4n) and one connecting e_x to the circuit in Ann(n/2, n);
- 3. there exists a $(2p_n p_c)$ -closed dual circuit with one defect around 0 in the annulus $Ann(n, 2n)^*$ which includes the edge e^* as its defect;
- 4. $\tau_e \in [p_n, 2p_n p_c)$; and
- 5. the p_c -open circuit in Ann(2n, 4n) is connected to ∞ by a p_n -open path.

By RSW arguments, [4, Lemma 6.3] and the fact that $L(2p_n - p_c)$ is comparable with n (see e.g. [13, (4.35)]),

$$\mathbb{P}(D_e) \simeq (p_n - p_c) \mathbb{P}_{cr}(A_n^{2,2}), \tag{7.1}$$

where $A_n^{2,2}$ is the event that the edge $e - e_x$ (we recall this notation means $(0, e_y - e_x)$) is connected to $\partial B(n)$ by two disjoint p_c -open paths and $(e - e_x)^*$ is connected to $\partial B(n)^*$ by two disjoint p_c -closed dual paths such that the open and closed paths alternate. Since D_e implies that $e \in \mathcal{O}$, we have for all e,

$$\mathbb{P}(e \in \mathcal{O}) \ge C_1(p_n - p_c)\mathbb{P}_{cr}(A_n^{2,2}). \tag{7.2}$$

Step 2. Let $A_{N,M}(e_x, p_c)$ be the event that there is a p_c -open circuit with 2 defects around e_x in $Ann(e_x, N, M)$. We will show that $\mathbb{P}(A_{N,M}(e_x, p_c), \theta_e E \mid e \in \mathcal{O})$ is close to $\mathbb{P}(\theta_e E \mid e \in \mathcal{O})$ for certain values of N < M. To this end, recall the definition of the event $H_{n,k}$ in (5.2) and write H for the event $H_{n,1}$. By (5.3) and (7.2), we can choose C_* independent of n such that

$$\mathbb{P}(\theta_e E, H^c \mid e \in \mathcal{O}) < \epsilon. \tag{7.3}$$

When the event H occurs, the invasion enters the $p_n(1)$ -open infinite cluster before it reaches e. Hence if e is an outlet, then e must be connected to $\partial B(e_x, n/4)$ by two disjoint $p_n(1)$ -open paths and e^* must be connected to $\partial B(e_x, n/4)^*$ by two disjoint p_c -closed dual paths such that the open and closed paths alternate and are all disjoint. Also, the weight τ_e must be in the interval $[p_c, p_n(1)]$. If, in addition, the event $A_{N,M}(e_x, p_c)$ does not occur, then there must be yet another p_c -closed dual path from $B(e_x, N)^*$ to $\partial B(e_x, M)^*$. This crossing has the property that it is disjoint from the two p_c -closed paths which are already present; however, it does not need to be disjoint from the $p_n(1)$ -open crossings. Therefore, $\mathbb{P}(\theta_e E, H, A_{N,M}(e_x, p_c)^c, e \in \mathcal{O})$ is at most

$$\leq C_2(p_n(1) - p_c) \mathbb{P}(A_n^{2,2}(p_n(1), p_c)) \mathbb{P}(A_{NM}^{2,3*}(p_n(1), p_c) \mid A_{NM}^{2,2}(p_n(1), p_c)),$$



where $A_{N,M}^{2,2}(p,q)$ denotes the event that B(N) is connected to $\partial B(M)$ by two p-open paths and that $B(N)^*$ is connected to $\partial B(M)^*$ by two q-closed paths so that the open and closed paths alternate and are all disjoint. The symbol $A_{N,M}^{2,3*}(p,q)$ signifies the event that $A_{N,M}^{2,2}(p,q)$ occurs but that there is an additional q-closed path connecting B(N) to $\partial B(M)$ which is disjoint from the two other q-closed paths but not necessarily from the two p-open paths. The above inequality, along with the estimate (7.2), gives that $\mathbb{P}(\theta_e E, H, A_{N,M}(e_x, p_c)^c \mid e \in \mathcal{O})$ is at most

$$\frac{C_2(p_n(1)-p_c)\mathbb{P}(A_n^{2,2}(p_n(1),p_c))}{C_1(p_n-p_c)\mathbb{P}_{cr}(A_n^{2,2})}\mathbb{P}(A_{N,M}^{2,3*}(p_n(1),p_c)\mid A_{N,M}^{2,2}(p_n(1),p_c)).$$

From (2.9) and Lemma 6.3 in [4], we can deduce

$$\frac{C_2(p_n(1) - p_c)\mathbb{P}(A_n^{2,2}(p_n(1), p_c))}{C_1(p_n - p_c)\mathbb{P}_{cr}(A_n^{2,2})} \le C_3(C_* \log n)^2, \tag{7.4}$$

so that

$$\mathbb{P}(\theta_{e}E, H, A_{N,M}(e_{x}, p_{c})^{c} | e \in \mathcal{O})$$

$$\leq C_{3}(C_{*} \log n)^{2} \mathbb{P}(A_{N,M}^{2,3*}(p_{n}(1), p_{c}) | A_{N,M}^{2,2}(p_{n}(1), p_{c})). \tag{7.5}$$

The above can be made less than ϵ provided that M/N grows fast enough with n. Let us assume this for the moment; we shall choose precise values for M and N at the end of the proof. Therefore, using (7.3), we have

$$|\mathbb{P}(\theta_e E \mid e \in \mathcal{O}) - \mathbb{P}(\theta_e E, H, A_{N,M}(e_x, p_c) \mid e \in \mathcal{O})| < 2\epsilon.$$
 (7.6)

Step 3. We now condition on the outermost p_c -open circuit with 2 defects in $Ann(e_x, N, M)$. For any circuit \mathcal{C} with 2 defects around the origin in the annulus Ann(N, M), let $D(\mathcal{C})$ be the event that it is the outermost p_c -open circuit with 2 defects. Notice that $D(\mathcal{C})$ depends only on the state of edges on or outside \mathcal{C} . For distinct \mathcal{C} , \mathcal{C}' (i.e. the sets of edges in \mathcal{C} and \mathcal{C}' are different or the sets of edges in \mathcal{C} and \mathcal{C}' are the same but the defects are different), the events $D(\mathcal{C})$, $D(\mathcal{C}')$ are disjoint. Therefore, the second term of (7.6) is equal to

$$\frac{1}{\mathbb{P}(e \in \mathcal{O})} \sum_{\mathcal{C} \subset Ann(N,M)} \mathbb{P}(\theta_e E, H, \theta_e D(\mathcal{C}), e \in \mathcal{O}), \tag{7.7}$$

where it is implied that in the sum, and in future sums like it, we only use circuits which enclose the origin.

Step 4. Let $Q(\theta_e C)$ be the event that there exists $f \neq e$ interior to $\theta_e C$ with $\tau_f \in [p_c, p_n(1)]$. We will now show that with high probability, the event $Q(\theta_e C)$ does not occur. In other words, we will bound the probability of the event $\{H, Q(\theta_e C), \theta_e D(C), e \in C\}$. Supposing that this event occurs, then both $\tau_e \in [p_c, p_n(1))$ and



 $A_{M,n}^{2,2}(p_n(1), p_c)$ must occur. Notice that the events $A_{M,n}^{2,2}(p_n(1), p_c)$, $\theta_e D(\mathcal{C})$, $\{\tau_e \in [p_c, p_n(1))\}$, and $Q(\theta_e \mathcal{C})$ are all independent. Hence $\mathbb{P}(H, Q(\theta_e \mathcal{C}), \theta_e D(\mathcal{C}), e \in \mathcal{O})$ is at most

$$\begin{split} & \mathbb{P}(A_{M,n}^{2,2}(p_n(1), p_c)) \mathbb{P}(\theta_e D(\mathcal{C})) \mathbb{P}(Q(\theta_e \mathcal{C})) \mathbb{P}(\tau_e \in [p_c, p_n(1))) \\ & \leq C_4 M^2 \frac{(p_n(1) - p_c)^2}{\mathbb{P}(A_M^{2,2}(p_n(1), p_c))} \mathbb{P}(A_n^{2,2}(p_n(1), p_c)) \mathbb{P}(\theta_e D(\mathcal{C})), \end{split}$$

where in the last inequality we used Corollary 6.1 from [4]. Consequently,

$$\begin{split} \mathbb{P}(H,\,Q(\theta_e\mathcal{C}),\,\theta_eD(\mathcal{C})|e\in\mathcal{O}) &\leq \left[\frac{C_4M^2(p_n(1)-p_c)}{\mathbb{P}_{p_n(1)}(A_M^{2,2})}\right] \\ &\times \left[\frac{(p_n(1)-p_c)\mathbb{P}(A_n^{2,2}(p_n(1),\,p_c))}{C_1(p_n-p_c)\mathbb{P}_{cr}(A_n^{2,2})}\right]\mathbb{P}(\theta_eD(\mathcal{C})), \end{split}$$

which, by (7.4), is at most

$$\frac{C_5(C_* \log n)^2 M^2}{\mathbb{P}_{p_n(1)}(A_M^{2,2})}(p_n(1) - p_c) \mathbb{P}(\theta_e D(\mathcal{C})).$$

As long as M is not too big, from $\mathbb{P}_{p_n(1)}(A_M^{2,2}) \simeq \mathbb{P}_{cr}(A_M^{2,2}) \geq cM^{-2}$ (see, e.g., Theorem 24 and Theorem 27 in [18]), we get an upper bound of

$$C_6(C_* \log n)^2 M^4(p_n(1) - p_c) \mathbb{P}(\theta_e D(\mathcal{C})) < \epsilon \mathbb{P}(\theta_e D(\mathcal{C})). \tag{7.8}$$

We will be able to choose such an M (in fact it will be of the order of a power of $\log n$), but we delay justification of this to the end of the proof. We henceforth assume that $\mathbb{P}(\theta_e E \mid e \in \mathcal{O})$ is within 3ϵ of

$$\frac{1}{\mathbb{P}(e \in \mathcal{O})} \sum_{\mathcal{C} \subset Ann(N,M)} \mathbb{P}(\theta_e E, H, \theta_e D(\mathcal{C}), Q(\theta_e \mathcal{C})^c, e \in \mathcal{O}). \tag{7.9}$$

Step 5. We write our configuration ω as $\eta \oplus \xi$, where η is the configuration outside or on $\theta_e \mathcal{C}$ and ξ is the configuration inside $\theta_e \mathcal{C}$. We condition on both η and τ_e : the summand of the numerator in (7.9) becomes

$$\mathbb{E}\left[\mathbb{P}(\theta_e E, H, \theta_e D(\mathcal{C}), Q(\theta_e \mathcal{C})^c, e \in \mathcal{O} \mid \tau_e, \eta)\right]. \tag{7.10}$$

Call the defected dual edges in $\theta_e C$ e_1^* and e_2^* . Given the value of τ_e , on the event $\theta_e D(C) \cap H \cap Q(\theta_e C)^c$, the event $\{e \in C\}$ occurs if and only if all of the following occur:



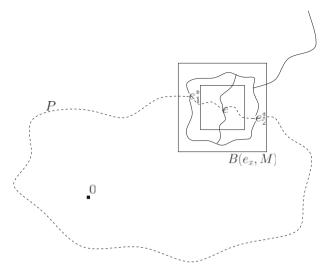


Fig. 7 The edge e is connected to the circuit $\theta_e C$ by two τ_e -open paths (the solid lines) and two p_c -closed paths (the dotted lines). The outer dotted circuit represents the τ_e -closed path P and the circuit $\theta_e C$ is connected to ∞ by a τ_e -open path. It is assumed that the invasion from the origin touches $\theta_e C$ before it touches P

- 1. e is connected to $\theta_e C \setminus \{e_1^*, e_2^*\}$ in the interior of $\theta_e C$ by two disjoint p_c -open paths;
- 2. e^* is connected to $\{e_1^*, e_2^*\}$ in the interior of $\theta_e C$ by two disjoint p_c -closed dual paths so that the p_c -closed paths and the p_c -open paths from item 1 alternate and are disjoint;
- 3. outside of $\theta_e C$, $\theta_e C$ is connected by a τ_e -open path to ∞ ;
- 4. $\tau_e \in [p_c, p_n(1))$; and
- 5. there exists a τ_e -closed dual path P outside of $\theta_e \mathcal{C}$, connecting e_1^* to e_2^* (both of which are τ_e -closed) such that (a) $P \cup B(e_x, M)^*$ contains a circuit around the origin and (b) the invasion graph contains a vertex from \mathcal{C} before it contains an edge f with f^* from P.

We will denote by $e \leftrightarrow_{2,2,p_c} \theta_e \mathcal{C}$ the event that the first two events occur, we will denote by $\theta_e \mathcal{C} \leftrightarrow_{\tau_e} \infty$ the third event, and we will use the symbol $X(\mathcal{C})$ for the fifth event. See Fig. 7 for an illustration of the intersection of these events. The term (7.10) becomes

$$\mathbb{E}\left[\mathbb{P}(\theta_{e}E, H, \theta_{e}D(C), e \leftrightarrow_{2,2,p_{c}} \theta_{e}C, \theta_{e}C \leftrightarrow_{\tau_{e}} \infty, Q(\theta_{e}C)^{c},\right.$$

$$\tau_{e} \in [p_{c}, p_{n}(1)), X(C) \mid \tau_{e}, \eta)\right]. \tag{7.11}$$

On the event $\{e \leftrightarrow_{2,2,p_c} \theta_e \mathcal{C}\} \cap \{\theta_e \mathcal{C} \leftrightarrow_{\tau_e} \infty\} \cap \tau_e \in [p_c, p_n(1)), \text{ the event } H \text{ occurs if and only if there exists a } p_n(1)\text{-open circuit } \mathcal{C}_0 \text{ enclosing the origin in } Ann(n/4, n/2)$ and either $\mathcal{C}_0 \overset{p_n(1)}{\longleftrightarrow} \theta_e \mathcal{C}$ or $\mathcal{C}_0 \overset{p_n(1)}{\longleftrightarrow} \infty$ outside of $\theta_e \mathcal{C}$. Denote by Y the event that such a circuit \mathcal{C}_0 exists and that either one of the above occur. Note that Y is measurable with respect to η . The term (7.11) becomes



$$\mathbb{E}\left[\mathbb{P}(\theta_{e}E, Y, \theta_{e}D(\mathcal{C}), e \leftrightarrow_{2,2,p_{c}}\theta_{e}\mathcal{C}, \theta_{e}\mathcal{C} \leftrightarrow_{\tau_{e}}\infty, Q(\theta_{e}\mathcal{C})^{c}, \tau_{e} \in [p_{c}, p_{n}(1)), X(\mathcal{C}) \mid \tau_{e}, \eta)\right]$$

$$= \mathbb{E}\left[1_{Y}1_{\theta_{e}D(\mathcal{C})}1_{\theta_{e}\mathcal{C} \leftrightarrow_{\tau_{e}}} \times 1_{\tau_{e} \in [p_{c}, p_{n}(1))}1_{X(\mathcal{C})}\mathbb{P}(\theta_{e}E, e \leftrightarrow_{2,2,p_{c}}\theta_{e}\mathcal{C}, Q(\theta_{e}\mathcal{C})^{c} \mid \tau_{e}, \eta)\right].$$
(7.12)

We now inspect the inner conditional probability. Clearly we have

$$\mathbb{P}(\theta_{e}E, e \leftrightarrow_{2,2,p_{c}} \theta_{e}C, Q(\theta_{e}C)^{c} | \tau_{e}, \eta) \leq \mathbb{P}(\theta_{e}E, e \leftrightarrow_{2,2,p_{c}} \theta_{e}C | \tau_{e}, \eta)
\leq \mathbb{P}(\theta_{e}E, e \leftrightarrow_{2,2,p_{c}} \theta_{e}C, Q(\theta_{e}C)^{c} | \tau_{e}, \eta)
+ \mathbb{P}(Q(\theta_{e}C)).$$
(7.13)

Using arguments similar to those that led to (7.9), one can show that the same choice of M and N that will make (7.8) hold will also make

$$\frac{1}{\mathbb{P}(e \in \mathcal{O})} \sum_{\mathcal{C} \subset Ann(N,M)} \mathbb{E}\left[1_{Y} 1_{\theta_{e}D(\mathcal{C})} 1_{\theta_{e}\mathcal{C} \leftrightarrow_{\tau_{e}} \infty} 1_{\tau_{e} \in [p_{c},p_{n}(1))} 1_{X(\mathcal{C})} \mathbb{P}(Q(\theta_{e}\mathcal{C}))\right] < \epsilon.$$

Therefore we conclude from (7.13) that $\mathbb{P}(\theta_e E \mid e \in \mathcal{O})$ is within 4ϵ of

$$\frac{1}{\mathbb{P}(e \in \mathcal{O})} \sum_{\mathcal{C} \subset Ann(N,M)} \mathbb{E} \left[1_{Y} 1_{\theta_{e}D(\mathcal{C})} 1_{\theta_{e}\mathcal{C} \leftrightarrow \tau_{e}} \otimes 1_{\tau_{e} \in [p_{c},p_{n}(1))} 1_{X(\mathcal{C})} \right. \\
\left. \times \mathbb{P}(\theta_{e}E, e \leftrightarrow_{2,2,p_{c}} \theta_{e}\mathcal{C} \mid \tau_{e}, \eta) \right]. \tag{7.14}$$

Step 6. Notice that since the events $\theta_e E$ and $e \leftrightarrow_{2,2,p_c} \theta_e C$ do not depend on τ_e or on η , we have

$$\mathbb{P}(\theta_e E, e \leftrightarrow_{2,2,p_c} \theta_e C \mid \tau_e, \eta) = \mathbb{P}_{cr}(E, 0 \leftrightarrow_{2,2} C) \text{ a.s.}, \tag{7.15}$$

where $0 \leftrightarrow_{2,2} \mathcal{C}$ denotes the event that the edge $e - e_x$ is connected to \mathcal{C} by two open paths and the dual edge $(e - e_x)^*$ is connected to $\{(e_1 - e_x)^*, (e_2 - e_x)^*\}$ by two closed paths such that all of these connections occur inside \mathcal{C} and the open and closed paths alternate. The quantity $\mathbb{P}_{cr}(E|0\leftrightarrow_{2,2} \mathcal{C})$ from the right side of (7.15) approaches $\tilde{v}^{2,2}(E)$ as long as $N \to \infty$ as $|e| \to \infty$ (this is a slight extension of Theorem 1.6) so, assuming this growth on N, we have

$$\frac{1}{(1+\epsilon)}\mathbb{P}_{cr}(E,0\leftrightarrow_{2,2}\mathcal{C})\leq \tilde{\nu}^{2,2}(E)\mathbb{P}_{cr}(0\leftrightarrow_{2,2}\mathcal{C})\leq \frac{1}{(1-\epsilon)}\mathbb{P}_{cr}(E,0\leftrightarrow_{2,2}\mathcal{C}).$$

It is straightforward now (following the end of the proof of [10, Theorem 3]) to show that

$$(1-\epsilon)(1-4\epsilon)\tilde{v}^{2,2}(E) < R < (1+\epsilon)(1+4\epsilon)\tilde{v}^{2,2}(E),$$

where *R* is the term which comprises the entire line of (7.15). Since *R* is within 4ϵ of $\mathbb{P}(\theta_e E \mid e \in \mathcal{O})$, all that remains is to choose *M* and *N* correctly.



Step 7. Choice of M and N. Recall, from (7.5), that we need the inequality

$$C_8(C_* \log n)^2 \mathbb{P}\left(A_{N,M}^{2,3*}(p_n(1), p_c) \mid A_{N,M}^{2,2}(p_n(1), p_c)\right) < \epsilon \tag{7.16}$$

to hold. In addition, we need to satisfy (7.8). Using the facts that $\mathbb{P}_{cr}(A_M^{2,2}) \geq C_9/M^2$ and

$$\mathbb{P}\left(A_{N,M}^{2,3*}(p_n(1), p_c) \mid A_{N,M}^{2,2}(p_n(1), p_c)\right) < C_{10}\left(\frac{N}{M}\right)^{\beta}$$
(7.17)

for some $\beta > 0$ (which is easily proved for what will be our choice of M and N, and which we assume for the moment), the reader may check that a choice of

$$N = \log n, \quad M = (\log n)^{2+2/\beta}$$

satisfies these two conditions for n large. The reason that this choice satisfies (7.8) is that $(\log n)^{\gamma}(p_n(1) - p_c) \to 0$ for any γ [use (2.9)] and the fact that the 4-arm exponent is strictly smaller than 2 (see, e.g., Section 6.4 in [21])).

We now prove (7.17). Let Q(M) be the event that there exists an edge in B(M) which has weight in the interval $[p_c, p_n(1))$. If Q(M) does not occur then the event $A_{N,M}^{2,3*}(p_n(1), p_c)$ implies the event $A_{N,M}^{2,3}(p_n(1), p_c)$ (i.e. the same event but with all five paths disjoint). Therefore, by Reimer's inequality, $\mathbb{P}(A_{N,M}^{2,3*}(p_n(1), p_c))$ is at most

$$\mathbb{P}\left(A_{N,M}^{2,3}(p_n(1), p_c)\right) + \mathbb{P}(Q(M)) \leq \mathbb{P}\left(A_{N,M}^{2,2}(p_n(1), p_c)\right) \mathbb{P}_{cr}\left(A_{N,M}^{0,1}\right) + |B(M)|(p_n(1) - p_c), \tag{7.18}$$

where $A_{N,M}^{0,1}$ is the event that B(N) is connected to $\partial B(M)$ by a p_c -closed path. Putting this estimate into (7.17), the term on its left is at most

$$\mathbb{P}_{cr}\left(A_{N,M}^{0,1}\right) + |B(M)| \frac{p_n(1) - p_c}{\mathbb{P}\left(A_{N,M}^{2,2}(p_n(1), p_c)\right)}.$$

Using the fact that

$$\mathbb{P}\left(A_{N,M}^{2,2}(p_n(1),\,p_c)\right) \geq \frac{\mathbb{P}\left(A_{N,M}^{2,3}(p_n(1),\,p_c)\right)}{\mathbb{P}_{cr}\left(A_{N,M}^{0,1}\right)} \geq \frac{C_{11}N^2}{M^2\mathbb{P}_{cr}\left(A_{N,M}^{0,1}\right)},$$

we see that the term (7.18) is at most

$$\mathbb{P}_{cr}\left(A_{N,M}^{0,1}\right) \left[1 + \frac{C_{12}M^4}{N^2} (p_n(1) - p_c)\right] \le 2\mathbb{P}_{cr}\left(A_{N,M}^{1,0}\right) \le C_{13} \left(\frac{N}{M}\right)^{\beta}$$



for some $\beta > 0$, as we have chosen N and M on the order of $\log n$. This shows (7.17) and completes the proof.

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